

# TFB Soundness

If  $e \Rightarrow v$  and  $\Gamma \vdash e : \tau$  then  $\Gamma \vdash v : \tau$ .

$$\text{True} \frac{}{\Gamma \vdash \text{True} : \text{Bool}} + \frac{\Gamma \vdash e_1 : \text{Int} \quad \Gamma \vdash e_2 : \text{Int}}{\Gamma \vdash e_1 + e_2 : \text{Int}}$$

$$\frac{\Gamma \vdash 4 : \text{Bool}}{\Gamma \vdash 4 : \text{Bool}} \text{ And } \frac{\Gamma \vdash e_1 : \tau_1 \rightarrow \tau_2}{\Gamma \vdash e_1 \text{ And } e_2 : \text{Int}}$$

$\Gamma \vdash e : \tau$  and  $x$  not free in  $e$  then  $\Gamma, x : \tau' \vdash e : \tau$

By induction on height of  $\Gamma \vdash e : \tau$  then by case analysis on the rule used.

Case of Int Rule:  $\frac{}{\Gamma \vdash n : \text{Int}}$  so by Int Rule  $\frac{}{\Gamma, x : \tau' \vdash n : \text{Int}}$

Case of + Rule:  $\frac{\Gamma \vdash e_1 : \text{Int} \quad \Gamma \vdash e_2 : \text{Int}}{\Gamma \vdash e_1 + e_2 : \text{Int}}$  By Ind,  $\Gamma, x : \tau' \vdash e_1 : \text{Int}$   
 $\Gamma, x : \tau' \vdash e_2 : \text{Int}$

By + Rule:

$$\Gamma, x : \tau' \vdash e_1 + e_2 : \text{Int}$$

Case of Let:  $\frac{\Gamma \vdash e_1 : \tau'' \quad \Gamma, x' : \tau'' \vdash e_2 : \tau}{\Gamma \vdash (\text{Let } x' : \tau'' = e_1 \text{ In } e_2) : \tau}$

By ind.  $\Gamma, x : \tau' \vdash e_2 : \tau''$

Two subcases:  $x = x'$  or  $x \neq x'$

Case  $x \neq x'$ : By ind.  $\Gamma, x' : \tau'', x : \tau' \vdash e_2 : \tau$

because  $x \neq x'$

$$\Gamma, x : \tau', x' : \tau'' \vdash e_2 : \tau$$

Case  $x = x'$ :

We know  $\Gamma, x' : \tau'' \vdash e_2 : \tau$ .  
 WTS  $\Gamma, x : \tau', x' : \tau'' \vdash e_2 : \tau$ .

$$\Gamma, x' : \tau'' = \Gamma, x : \tau', x' : \tau'' \quad \Gamma, c : \text{Bool} = \left\{ \begin{array}{l} a \mapsto \text{Int} \\ b \mapsto \text{Bool} \\ c \mapsto \text{Bool} \end{array} \right\}$$

$$\Gamma = \left\{ \begin{array}{l} a \mapsto \text{Int} \\ b \mapsto \text{Bool} \end{array} \right\} \quad \Gamma, c : \text{Int} = \left\{ \begin{array}{l} a \mapsto \text{Int} \\ b \mapsto \text{Bool} \\ c \mapsto \text{Int} \end{array} \right\} \quad \Gamma, c : \text{Int}; c : \text{Bool} = \left\{ \begin{array}{l} a \mapsto \text{Int} \\ b \mapsto \text{Bool} \\ c \mapsto \text{Bool} \end{array} \right\}$$

$\Gamma \vdash v : \tau$  and  $\Gamma, x : \tau \vdash e : \tau'$  then  $\Gamma \vdash e[v/x] : \tau'$

By induction on height of  $\Gamma, x : \tau \vdash e : \tau'$  then by case analysis on rule:

Integer case:  $e = n$  for  $n \in \mathbb{Z}$

$\Gamma, x : \tau \vdash e : \tau'$  is  $\frac{}{\Gamma, x : \tau \vdash n : \text{Int}}$

Want to show is  $\Gamma \vdash n[v/x] : \text{Int}$

$n[v/x] = n$   $\Gamma \vdash n : \text{Int}$  (by Int Rule)

+ case:  $e = e_1 + e_2$

$\Gamma, x : \tau \vdash e : \tau'$  is  $\frac{\frac{}{\Gamma, x : \tau \vdash e_1 : \text{Int}} \quad \frac{}{\Gamma, x : \tau \vdash e_2 : \text{Int}}}{\Gamma, x : \tau \vdash e_1 + e_2 : \text{Int}}}$

Want to show is

$\Gamma \vdash e[v/x] : \text{Int}$   
 $\Gamma \vdash (e_1 + e_2)[v/x] : \text{Int}$   
 $\Gamma \vdash (e_1[v/x]) + (e_2[v/x]) : \text{Int}$

By ind

$\Gamma \vdash e_1[v/x] : \text{Int}$

$\Gamma \vdash e_2[v/x] : \text{Int}$

By + rule

Function Case:  $\Gamma, x : \tau \vdash e : \tau'$  is

$\frac{\Gamma, x : \tau, x' : \tau_1 \vdash e' : \tau_2}{\Gamma, x : \tau \vdash (\text{Function } x' : \tau_1 \rightarrow e') : \tau_1 \rightarrow \tau_2}$

WTS

$\Gamma \vdash (\text{Function } x' : \tau_1 \rightarrow e')[v/x] : \tau_1 \rightarrow \tau_2$

Case  $x \neq x'$ :  $(\text{Function } x' : \tau_1 \rightarrow e')[v/x] = \text{Function } x' : \tau_1 \rightarrow (e'[v/x])$

$\Gamma, x' : \tau_1, x : \tau \vdash e' : \tau_2$  by Ind  $\Gamma, x' : \tau_1 \vdash e'[v/x] : \tau_2$

$\Gamma \vdash \text{Function } x' : \tau_1 \rightarrow (e'[v/x]) : \tau_1 \rightarrow \tau_2$

Case  $x = x'$ :  $(\text{Function } x' : \tau_1 \rightarrow e')[v/x] = \text{Function } x' : \tau_1 \rightarrow e'$

$\Gamma \vdash \text{Function } x' : \tau_1 \rightarrow e' : \tau_1 \rightarrow \tau_2$