

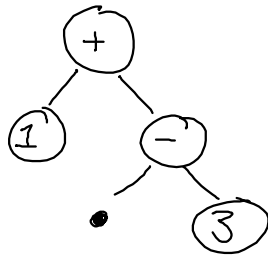
Operational Equivalence

$$4 \quad \cong \quad 1 + 3$$

A context is an expression with a hole in it.

$$C ::= \bullet \mid \text{Not } C \mid C \text{ Or } e \mid e \text{ Or } C \mid \dots$$

$$1 + (\bullet - 3)$$



$$(1 + (\bullet - 3))[4] = (1 + (4 - 3))$$

Filling a context: we write $C[e]$ to refer to the expression in which e replaces the $\bullet$ in C .

Operational equivalence: we write $e_1 \cong e_2$ to mean

$$\forall C. C[e_1] \Rightarrow v_1 \text{ iff } C[e_2] \Rightarrow v_2 \\ (\text{for some } v_1 \text{ and } v_2)$$

Values v_1 and v_2 can be proven to be the same.

$$4 \not\cong 5 \quad \text{If } \bullet = 4 \text{ Then } 0 \text{ Else } 1 + \bullet$$

Practice w/ contexts: for each of these contexts, find an expression to plug in that produces False. (Then: find an expression that doesn't contain boolean constant.)

(Function $a \rightarrow \bullet$) True	False	$1=8$
$\bullet \quad \circ$	(Function $a \rightarrow \text{False}$)	(Function $b \rightarrow b=1$)
(Function $f \rightarrow \text{Not } (f \ 4)$) $\bullet$	(Function $a \rightarrow \text{True}$)	(Function $a \rightarrow 1=1$)

we write $e_1 \cong e_2$ to mean

$\forall C. C[e_1] \Rightarrow v_1 \text{ iff } C[e_2] \Rightarrow v_2$
(for some v_1 and v_2)

$\text{True} \not\cong \text{False}$	(Function $a \rightarrow \text{If } a \text{ Then } \circ \text{ Else } 1+\text{True}$) $\bullet$
$n \not\cong n+1-1$	(Function $n \rightarrow \text{Not } \bullet$) True
(Function $a \rightarrow b$) $\not\cong$ (Function $a \rightarrow d$)	Let $b=4$ In $\bullet \ 3$

$\cong$ is an equivalence relation

reflexive

$$e \cong e$$

symmetric

$$e_1 \cong e_2 \text{ iff } e_2 \cong e_1$$

transitive

$$\text{if } e_1 \cong e_2 \text{ and } e_2 \cong e_3 \text{ then } e_1 \cong e_3$$

Op eq:

congruent

$$\text{if } e_1 \cong e_2 \text{ then } \forall C. C[e_1] \cong C[e_2]$$

All divergent expressions are operationally equivalent.

An expression is "used" during evaluation if a subproof shows that it evaluates.

Assume for contradiction that two divergent e_1 and e_2 and some C for which $C[e_1] \Rightarrow v$ and $C[e_2] \not\Rightarrow$.

$C[e_1]$ is closed. Since $C[e_1]$ evaluates, it must not use $\bullet$.

Then $C[e_2]$ does not use $\bullet$. So $C[e_2] \Rightarrow v$.

By contradiction: no such C exists.