

$$\frac{(x:\tau) \in \Gamma}{\Gamma \vdash x:\tau}$$

$$\frac{v \in \mathbb{Z}}{\Gamma \vdash v:\text{Int}}$$

$$\frac{\Gamma \{x:\tau\} \vdash e:\tau'}{\Gamma \vdash (\text{Function } x:\tau \rightarrow e):\tau \rightarrow \tau'}$$

$$\frac{\Gamma \vdash e_1:\tau_1 \quad \dots \quad \Gamma \vdash e_n:\tau_n}{\Gamma \vdash \{l_1=e_1; \dots; l_n=e_n\}:\{\tau_1; \dots; \tau_n\}}$$

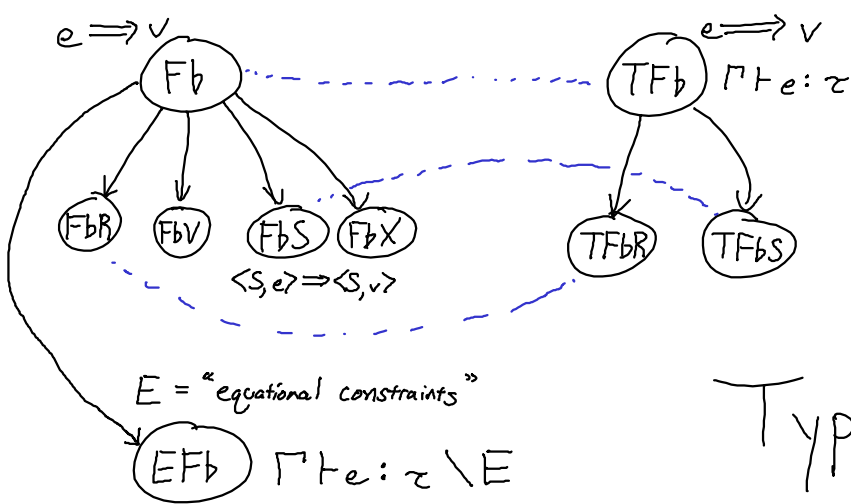
$$\frac{\Gamma \vdash e:\tau \quad \tau <:\tau'}{\Gamma \vdash e:\tau'}$$

Prove:

$$\frac{\frac{4 \in \mathbb{Z}}{\{r:\{a:\text{Int}\}\} \vdash 4:\text{Int}}}{\emptyset \vdash (\text{Function } r:(\{a:\text{Int}\}) \rightarrow 4) : \{a:\text{Int}\} \rightarrow \text{Int}}$$

$$\frac{\frac{R}{\text{Int} <:\text{Int}} \quad \frac{S_{\text{Rec}}}{\{a:\text{Int}; b:\text{Int}\} <:\{a:\text{Int}\}} \quad \frac{R}{\text{Int} <:\text{Int}}}{\frac{S_{\text{Fun}}}{\{a:\text{Int}\} \rightarrow \text{Int} <:\{a:\text{Int}; b:\text{Int}\} \rightarrow \text{Int}}}$$

$$\emptyset \vdash (\text{Function } r:(\{a:\text{Int}\}) \rightarrow 4) : \{a:\text{Int}; b:\text{Int}\} \rightarrow \text{Int}$$



$$\Gamma ::= \{x:\tau, \dots\}$$

$$\tau ::= \text{Int} \mid \text{Bool} \mid \tau \rightarrow \tau \mid \alpha$$

$$\alpha ::= 'a' \mid 'b' \mid \dots$$

$$E ::= \{\tau = \tau, \dots\}$$

$$\text{Int} \setminus \emptyset$$

$$'a' \setminus \{'a' = \text{Int}\}$$

$$'b' \setminus \{\text{Int} = 'a', 'b' = 'a'\}$$

$$'a' \setminus \{'a' = \text{Int}, \text{Bool} = 'a'\}$$

$$\Gamma \vdash e : \tau \setminus E$$

Assuming the type environment Γ , e has type τ under the assumptions of the equations E .

1. Infer a type and equation set for expression.
2. Deductively close over equation set.
3. Check for inconsistencies.
4. Type substitution

$$\boxed{\Gamma \vdash e : \tau \setminus E}$$

Function $q \rightarrow \text{Not } q$

$$\frac{v \in \mathbb{Z}}{\Gamma \vdash v : \text{Int} \setminus \emptyset}$$

$$\frac{(x:\tau) \in \Gamma}{\Gamma \vdash x : \tau \setminus \emptyset}$$

$$\frac{\Gamma \{x:\alpha\} \vdash e : \tau \setminus E \quad \alpha \text{ is fresh}}{\Gamma \vdash (\text{Function } x \rightarrow e) : \alpha \rightarrow \tau \setminus E}$$

$$\frac{\Gamma \vdash e : \tau \setminus E}{\Gamma \vdash \text{Not } e : \text{Bool} \setminus E \cup \{\tau = \text{Bool}\}}$$

$$\frac{(q:'a) \in \{q:'a\}}{\{q:'a\} \vdash q : 'a \setminus \emptyset}$$

$$'a \rightarrow \text{Bool} \setminus \{ 'a = \text{Bool} \}$$

$\Downarrow$

$$\text{Bool} \rightarrow \text{Bool}$$

$$\{q:'a\} \vdash \text{Not } q : \text{Bool} \setminus \{ 'a = \text{Bool} \}$$

$$\emptyset \vdash (\text{Function } q \rightarrow \text{Not } q) : 'a \rightarrow \text{Bool} \setminus \{ 'a = \text{Bool} \}$$

$$\frac{\Gamma \vdash e_1 : \tau_1 \setminus E_1 \quad \Gamma \vdash e_2 : \tau_2 \setminus E_2}{\Gamma \vdash e_1 + e_2 : \text{Int} \setminus E_1 \cup E_2 \cup \{\tau_1 = \text{Int}, \tau_2 = \text{Int}\}}$$

$$\frac{\Gamma \vdash e_1 : \tau_1 \setminus E_1 \quad \Gamma \vdash e_2 : \tau_2 \setminus E_2 \quad \alpha \text{ is fresh}}{\Gamma \vdash e_1 \ e_2 : \alpha \setminus E_1 \cup E_2 \cup \{\tau_1 = \tau_2 \rightarrow \alpha\}}$$