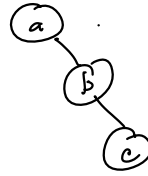
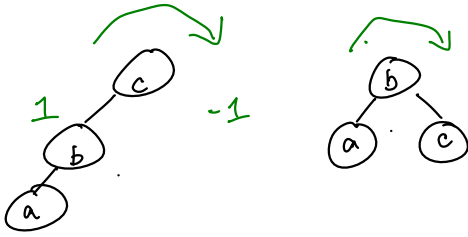


BST is a kind of Dictionary

Just the BST algorithms (w/o rebalancing) : slow ( $O(\text{height})$ )

"Balanced" binary search tree : AVL tree

BST where all nodes have children whose heights differ by no more than one



Function rotateRight(node):

If node.left == null: ||

a ← node

b ← node.left

Y ← b.right

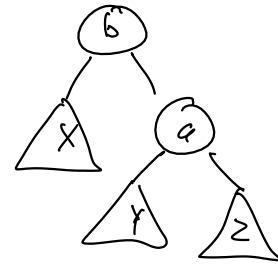
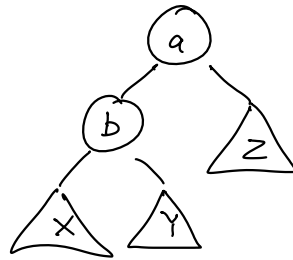
a.left ← Y

b.right ← a

Return b

EndFunction

$O(1)$



Function insertInSubtree(node, key, value):

If node == null:

Return new Node(key, value)

Else If key < node.key:

node.left ← insertInSubtree(node.left, key, value)

node ← rebalance(node)

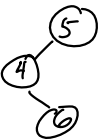
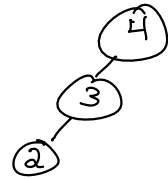
recalc Height(node)

Return node

Else ...

AVL tree

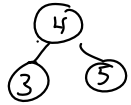
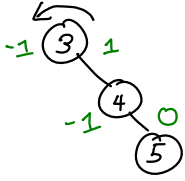
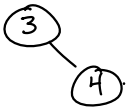
AVL free



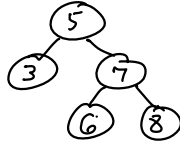
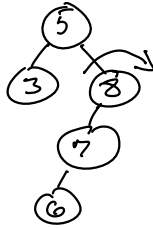
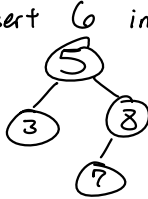
Insufficient

BST is a BT where left children are smaller and right children are larger

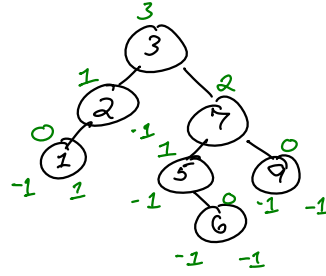
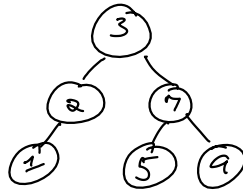
insert 5 into



insert 6 into

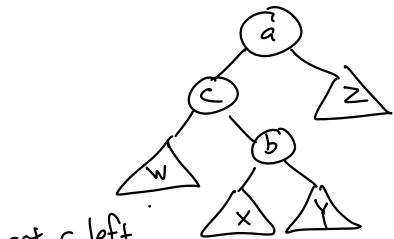
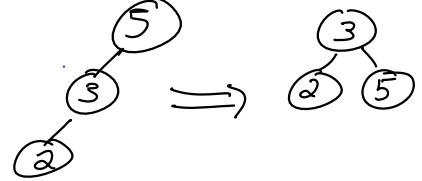
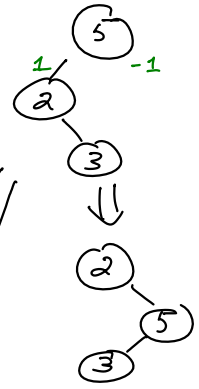


insert 6 into



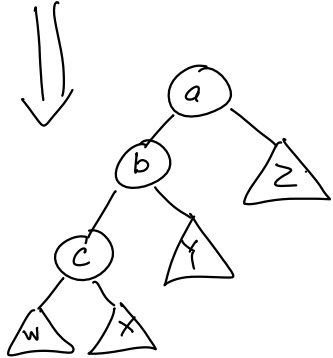
NO rotation

insert 3 into



Left-right

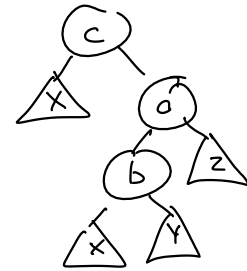
rot c left



Left-left

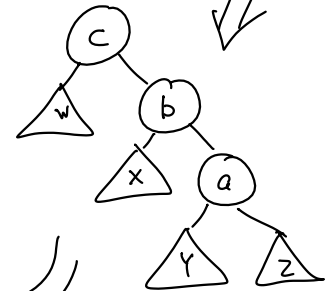
rot right  
a

Right-left



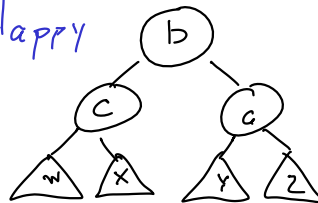
rot a right

Right-right



rot c left

Happy



Function rebalance(node):

$\text{delta} \leftarrow \text{node.right.height} - \text{node.left.height}$

If  $\text{delta} > 1$ :

If  $\text{node.right.right.height} < \text{node.right.left.height}$ :

$\text{node.right} \leftarrow \text{rotateRight}(\text{node.right})$

EndIf

$\text{node} \leftarrow \text{rotateLeft}(\text{node})$

Else If  $\text{delta} < -1$ :

If  $\text{node.left.left.height} < \text{node.left.right.height}$ :

$\text{node.left} \leftarrow \text{rotateLeft}(\text{node.left})$

EndIf

$\text{node} \leftarrow \text{rotateRight}(\text{node})$

EndIf

Return node

