

CS 43: Computer Networks

BGP

November 26, 2025



Adapted from Slides by: Kurose & Ross, D. Choffnes, K. Webb, J. Rexford

The Network Layer!

Application: the application (e.g., the Web, Email)

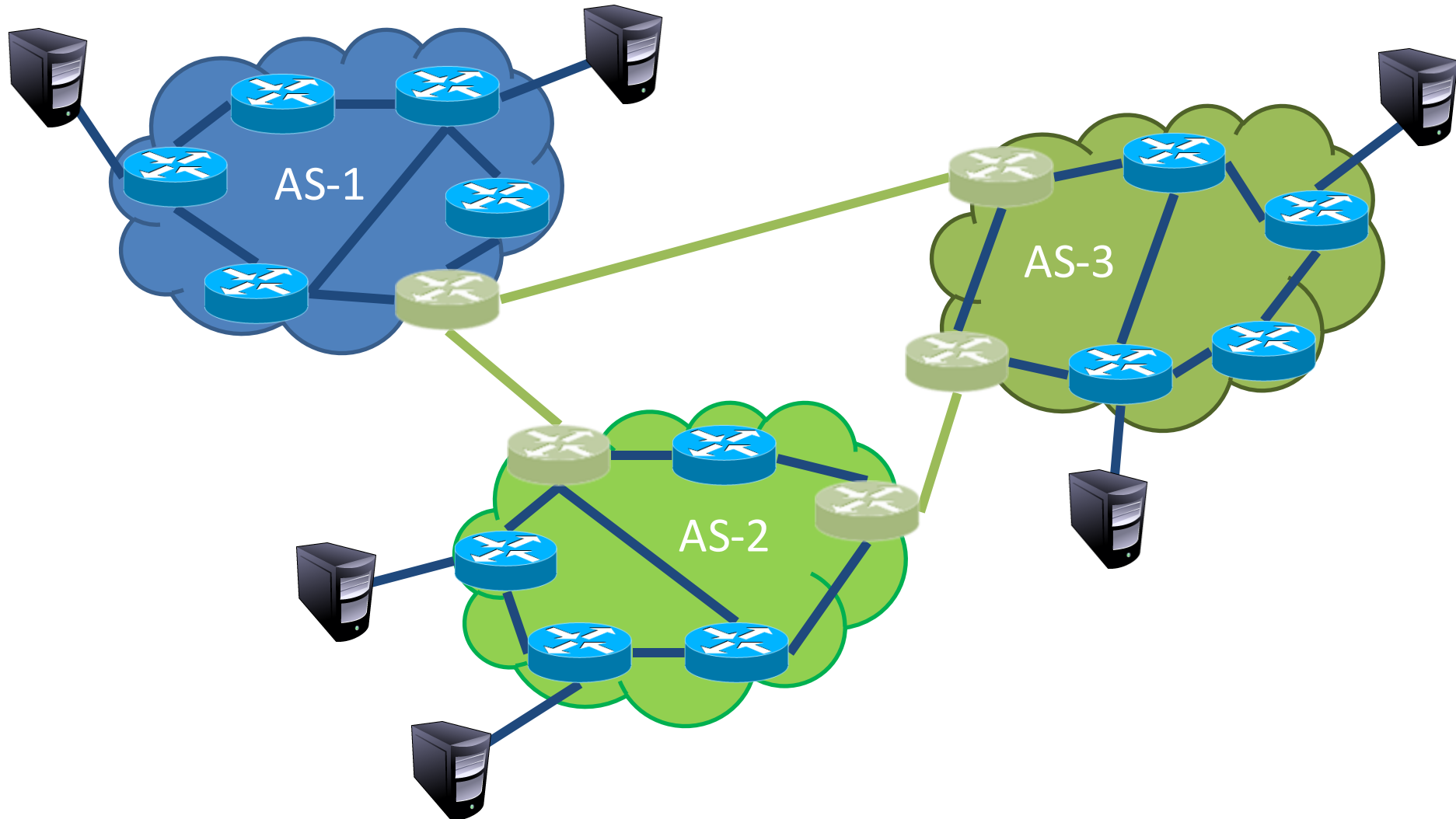
Transport: end-to-end connections, reliability

Network: routing

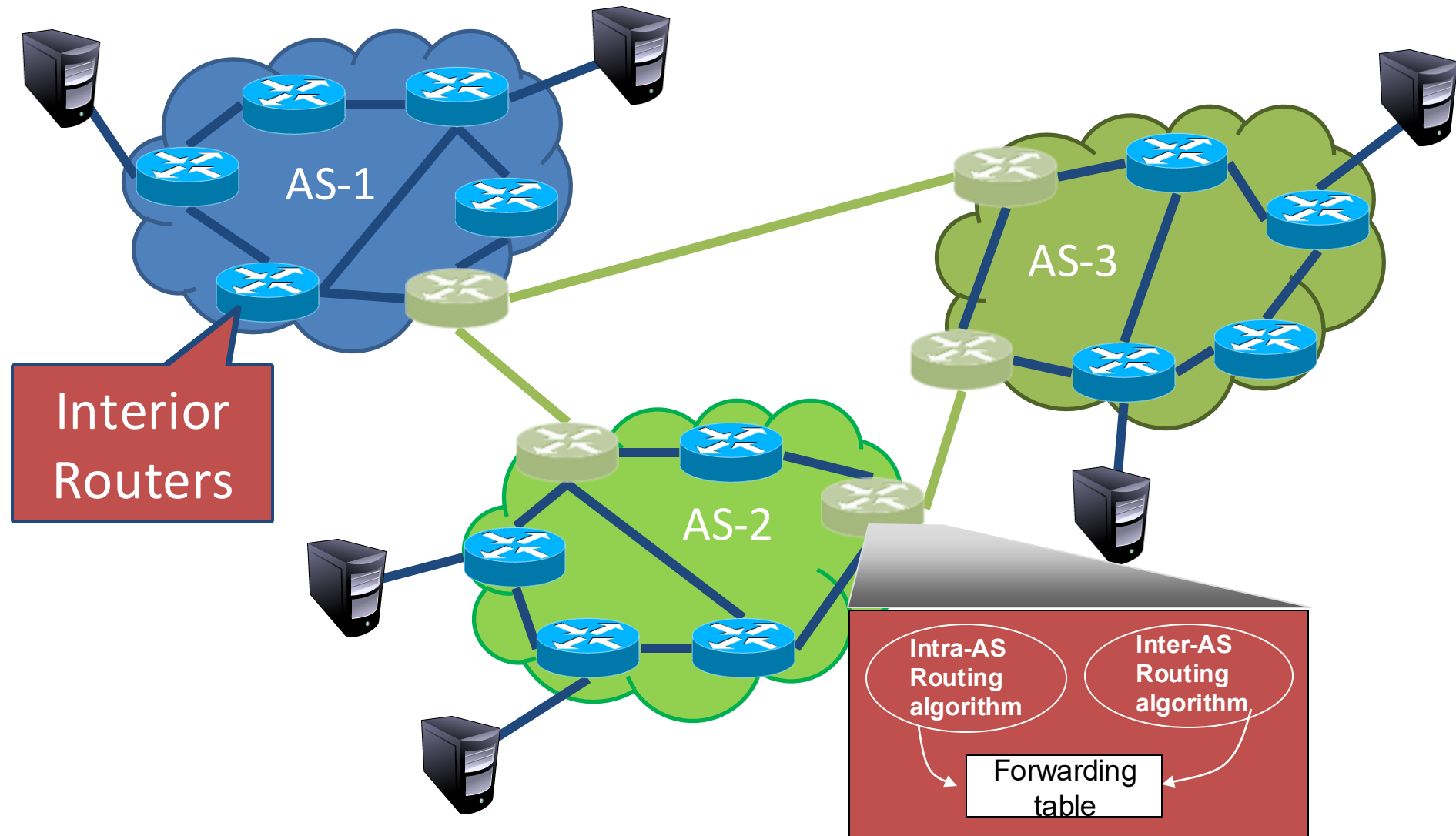
Link (data-link): framing, error detection

Physical: 1's and 0's/bits across a medium (copper, the air, fiber)

Hierarchical routing: Autonomous Systems

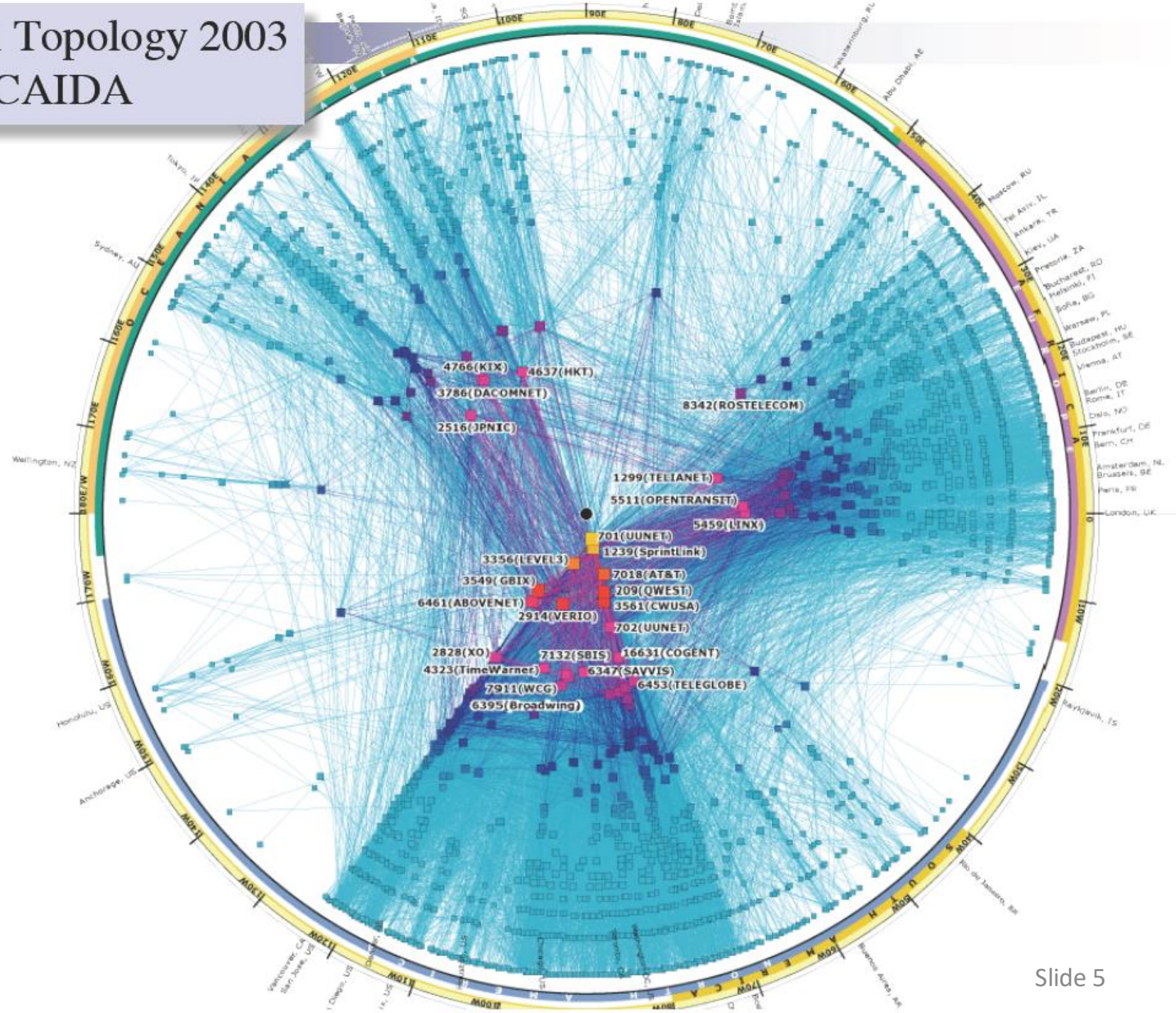


Hierarchical routing: Interconnected ASes

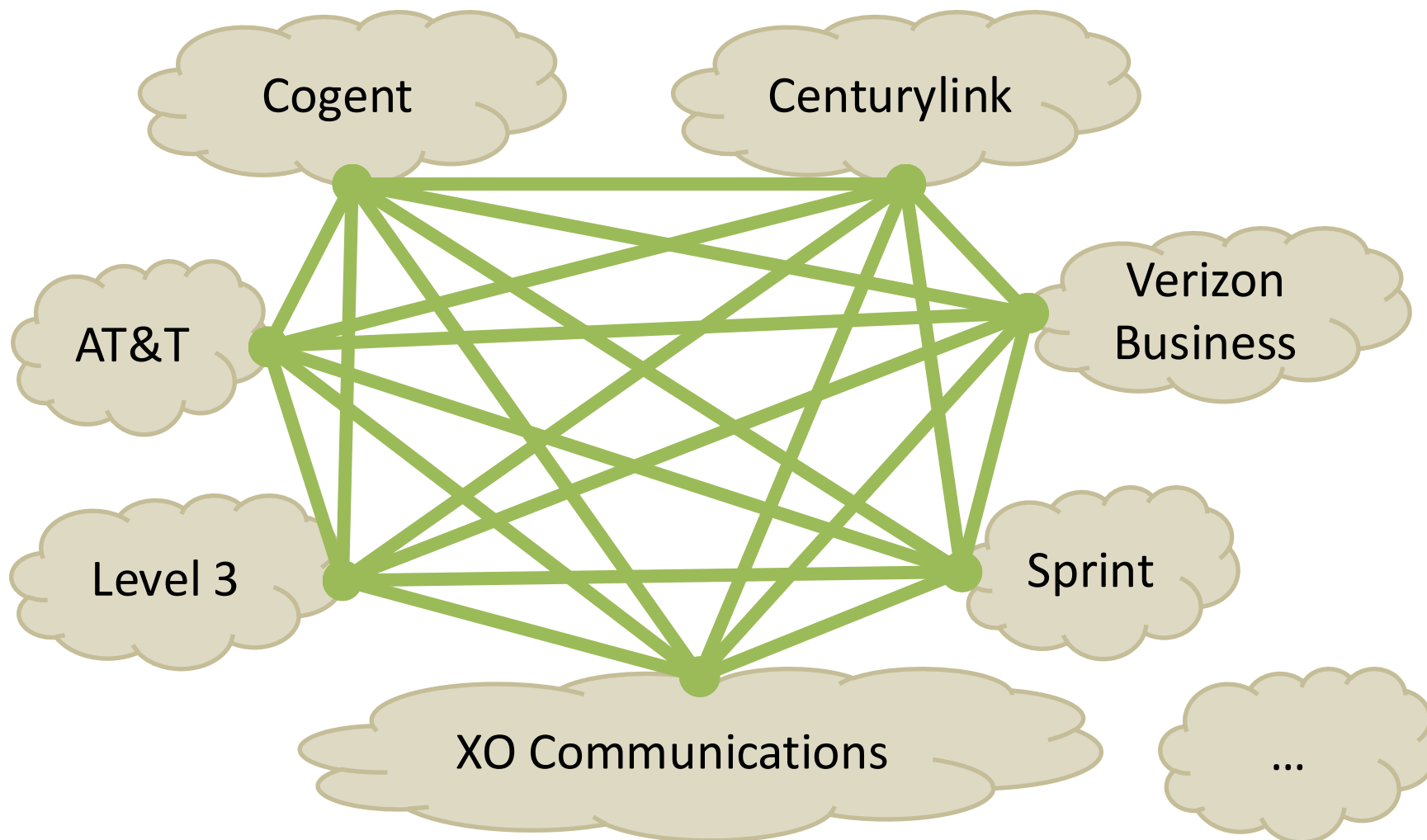


AS-level Topology 2003

Source: CAIDA

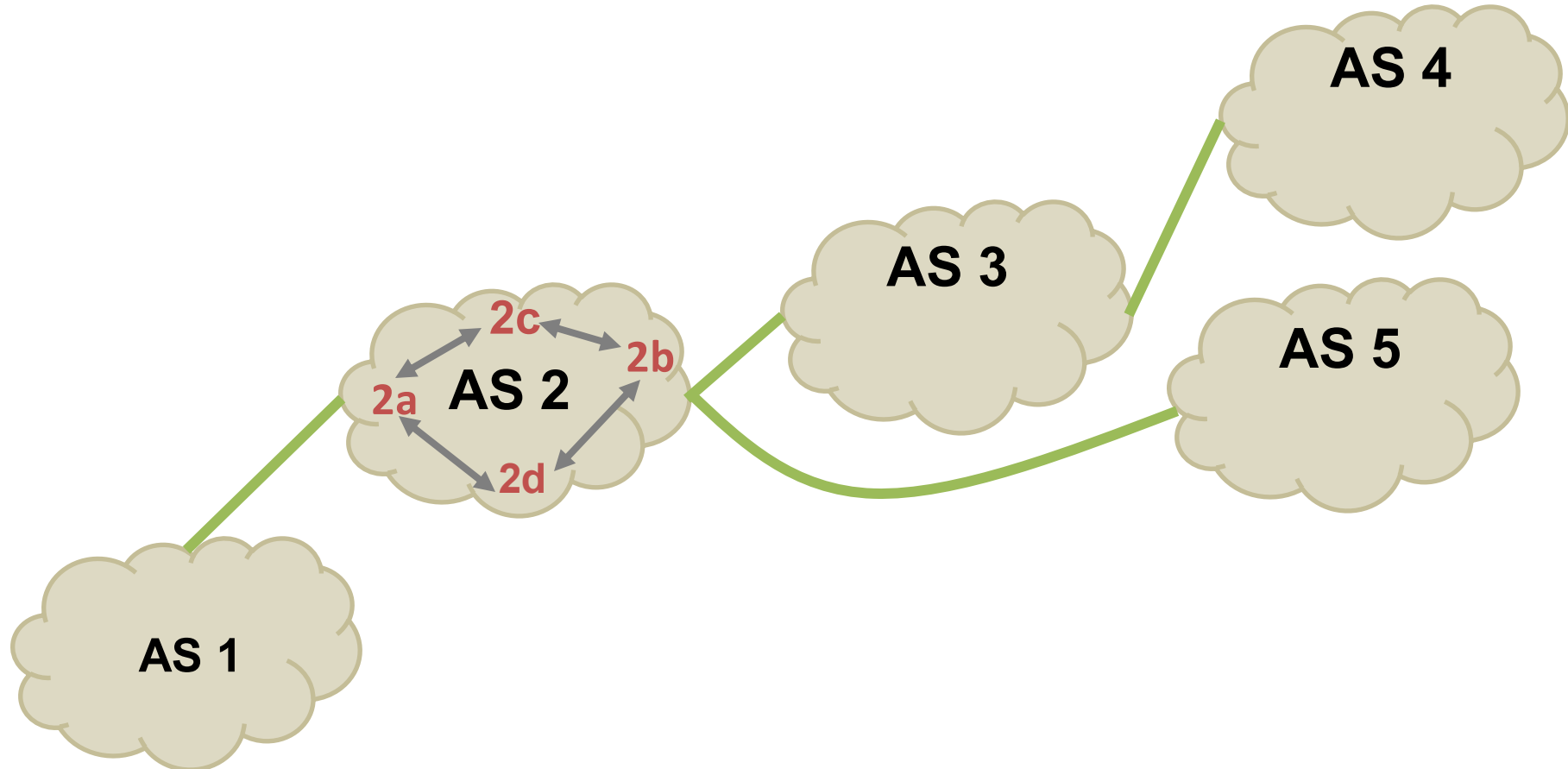


Tier-1 ISP Peering



Path Vector Protocol

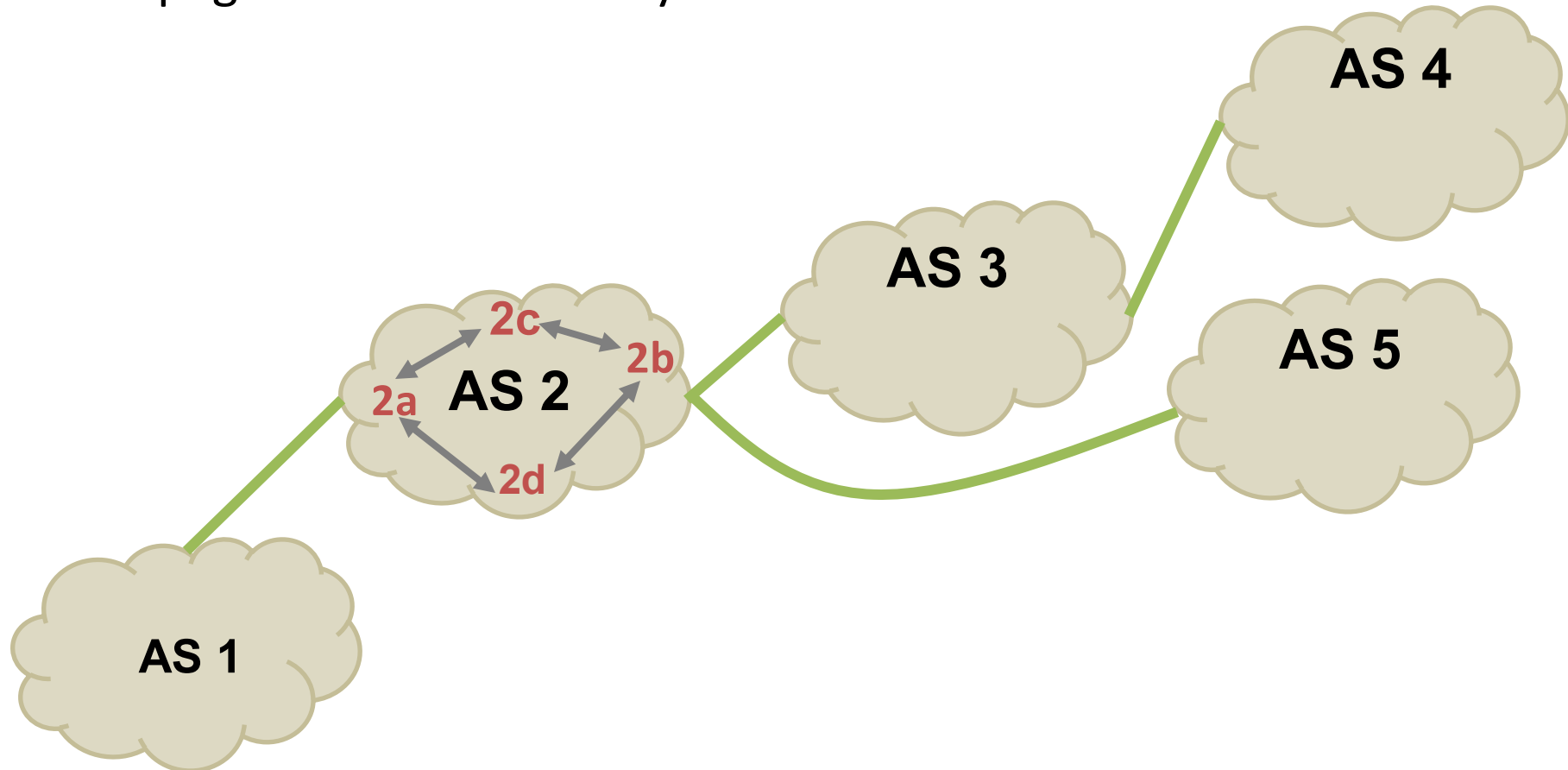
- Key idea: advertise the entire path
 - Distance vector: send *distance metric* per dest d
 - Path vector: send the *entire path* for each dest d



Inter-domain (Inter-ISP) Routing

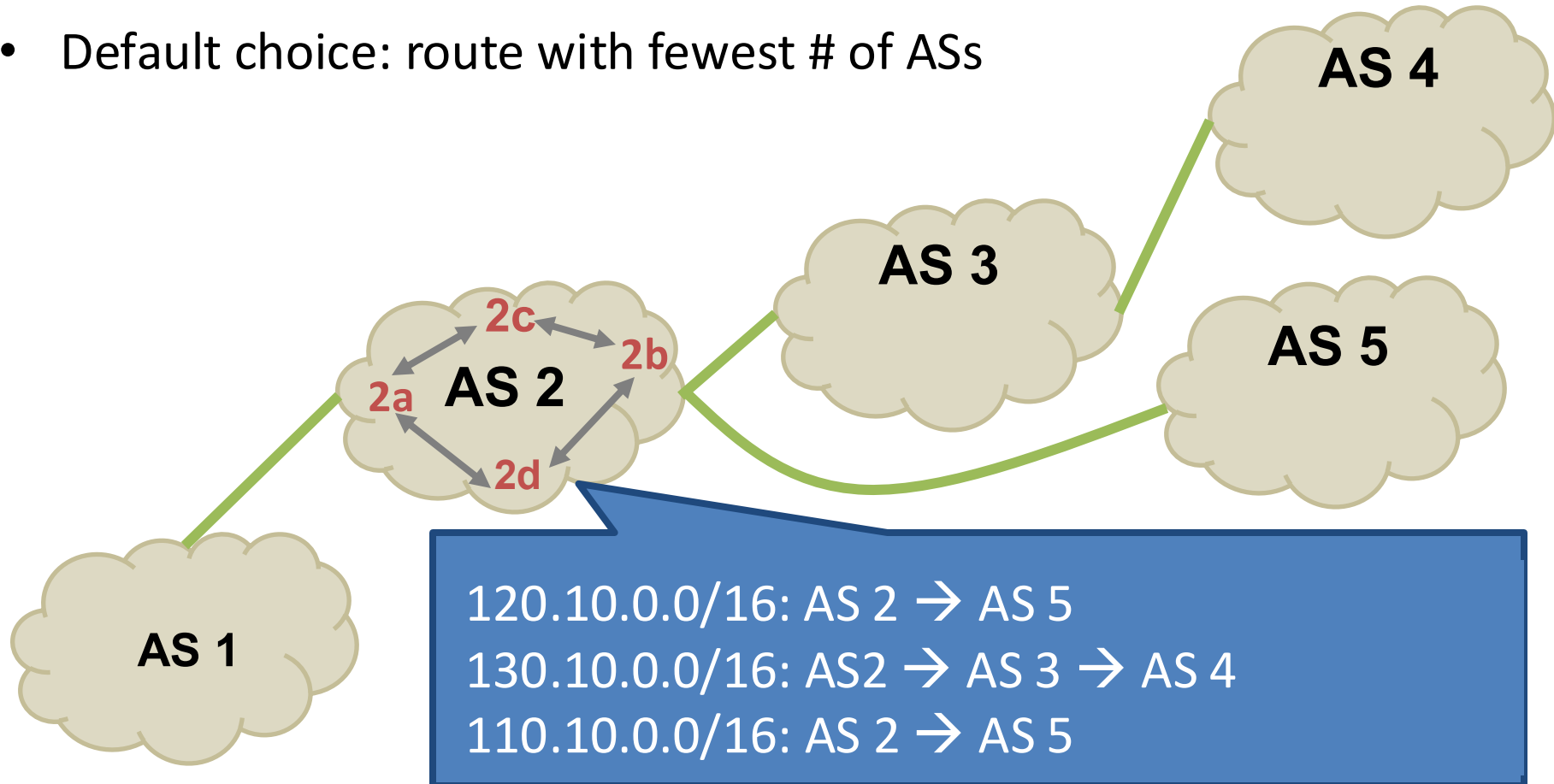
AS2 must:

1. Learn destinations reachable through AS2
2. Propagate this reachability info to all routers in AS2



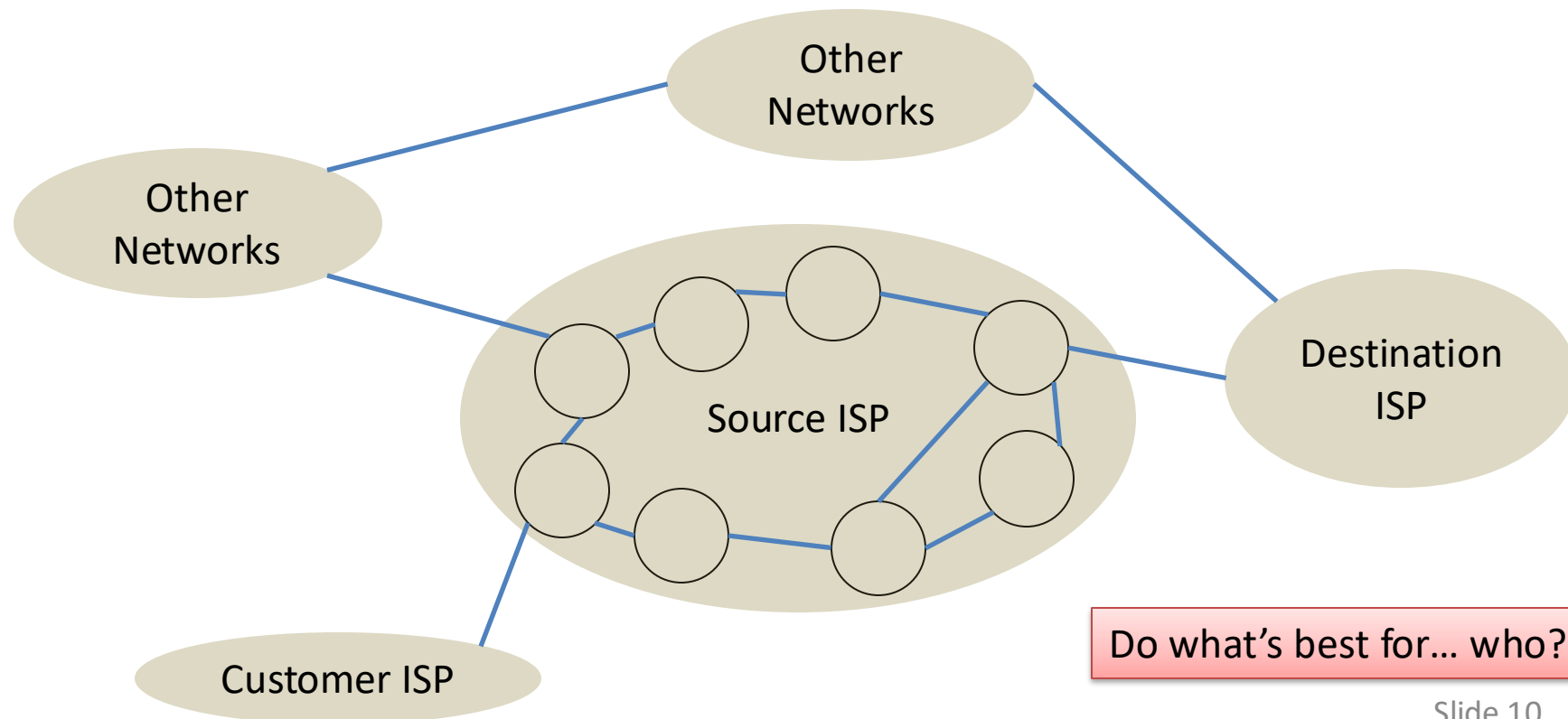
Path Vector Protocol

- AS-path: sequence of ASs a route traverses
 - Like distance vector, plus additional information
- Used for loop detection and to apply policy
- Default choice: route with fewest # of ASs



Routing Policy

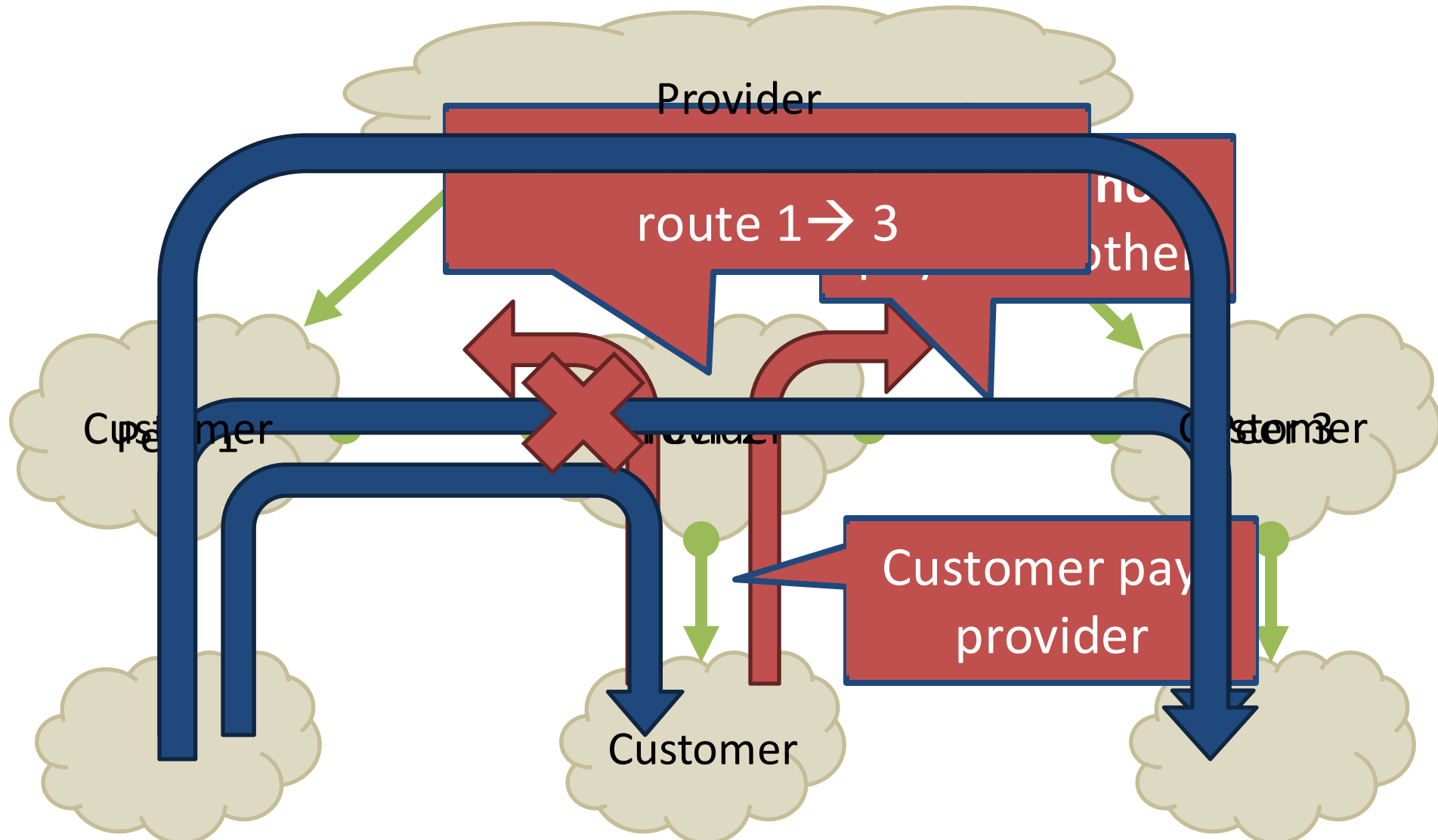
- How should the ISP route the customer's traffic to the destination?



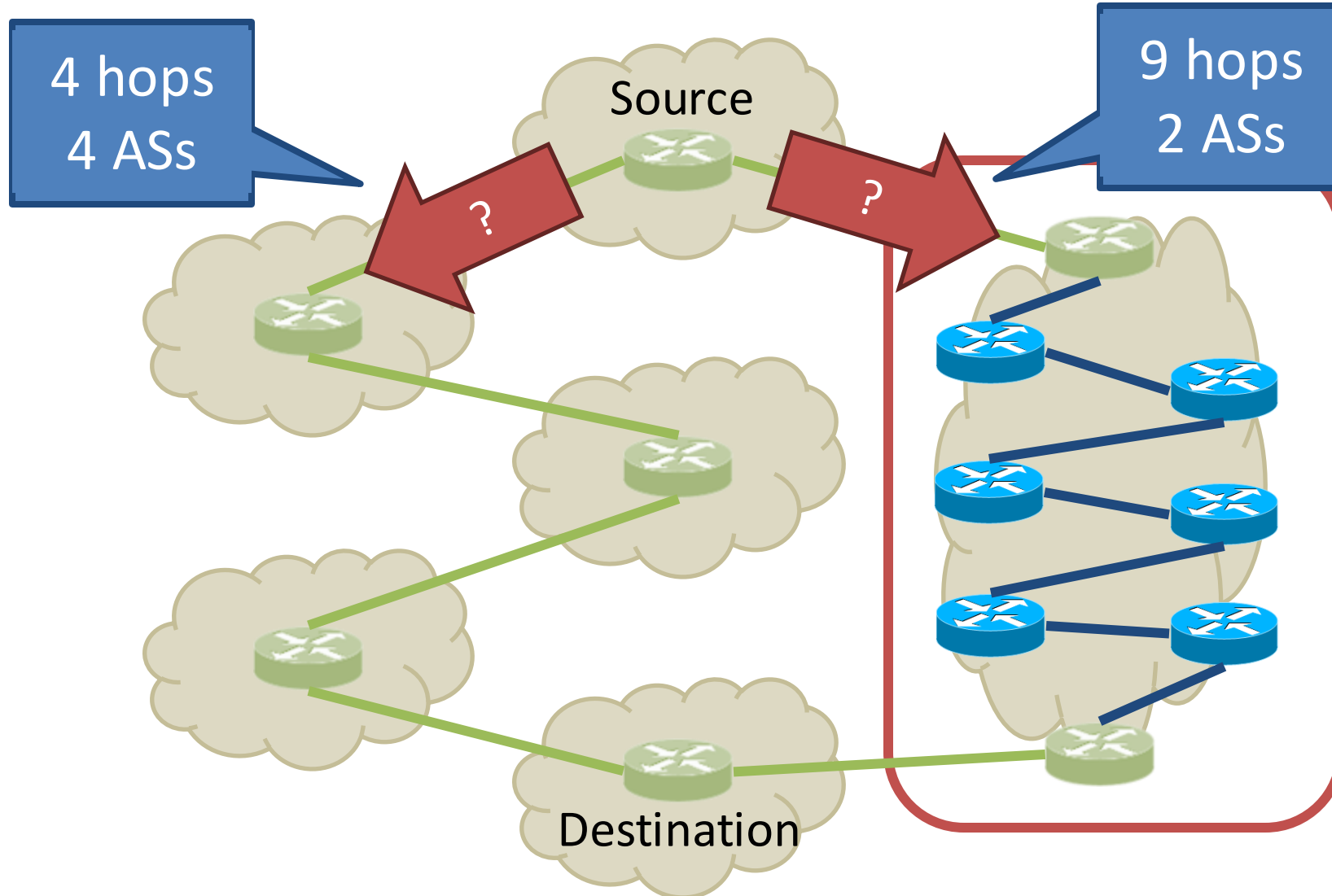
Which routes a BGP router advertises will depend on...

- A. which ISPs have contractual agreements.
- B. the shortest path to a subnet/prefix.
- C. which subnets are customers of an ISP.
- D. More than one of the above. (which?)

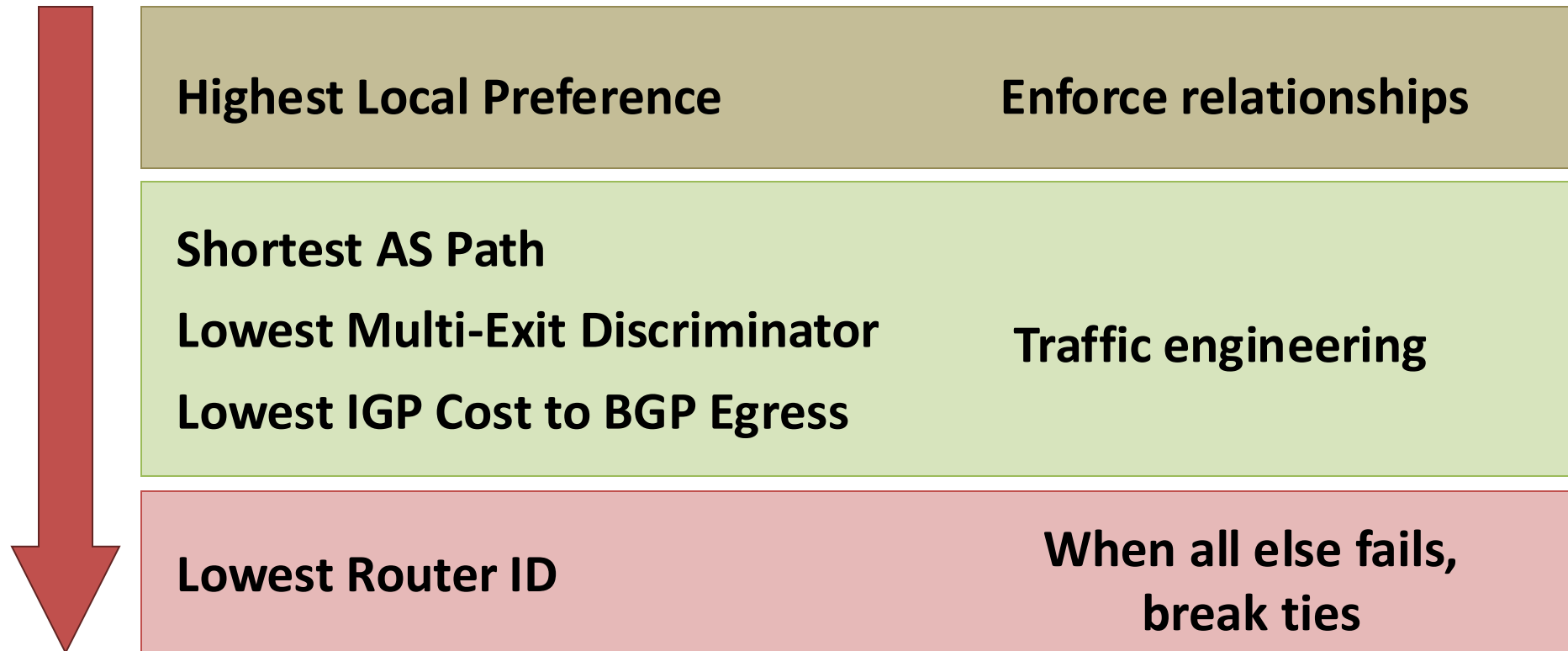
BGP Relationships



Shortest AS Path $\neq$ Shortest Path



Route Selection Summary



Peering/Interconnection Wars

- Peer

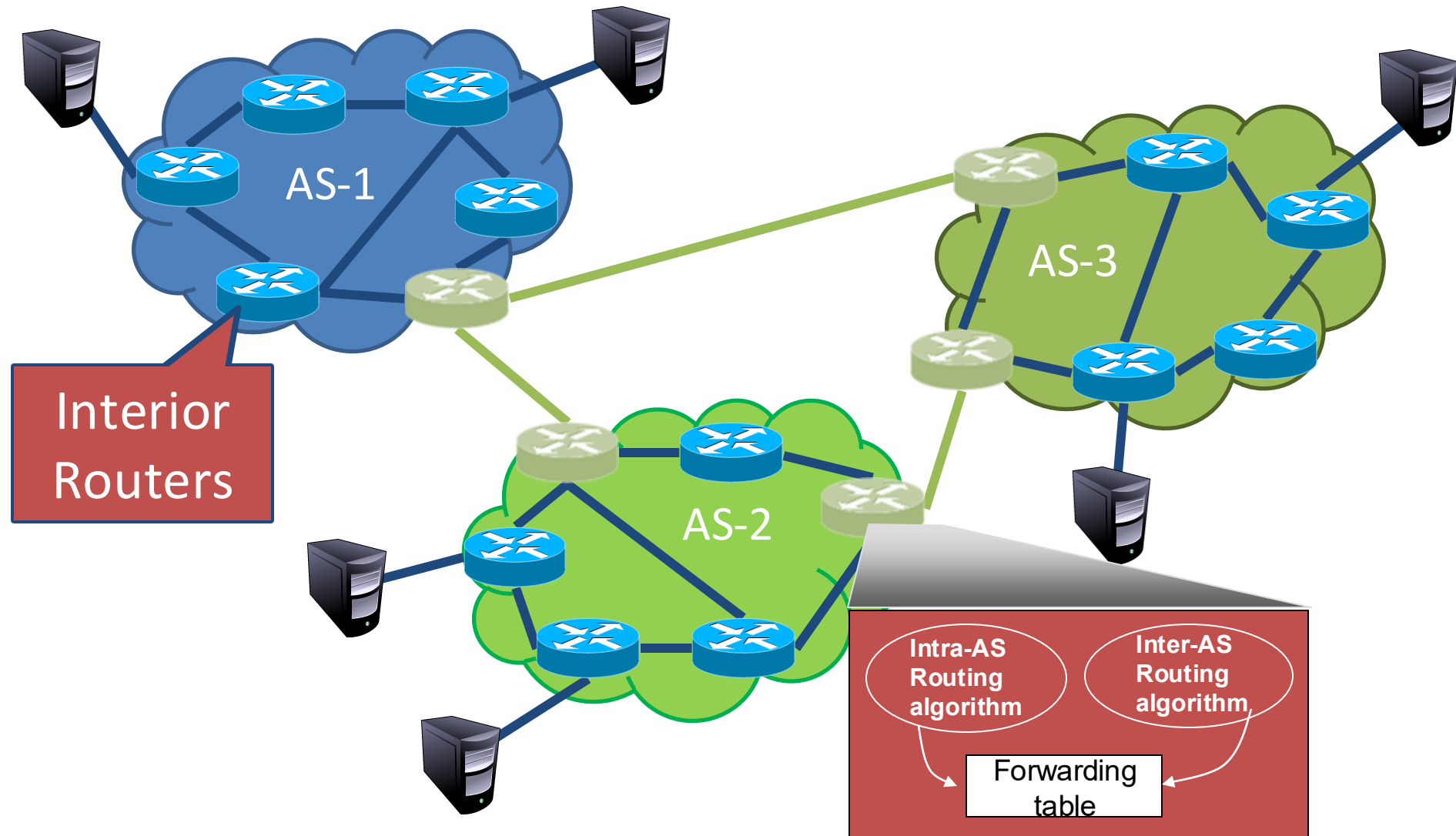
- Reduce upstream costs
- Improve end-to-end performance
- May be the only way to connect to parts of the Internet

- Don't Peer

- You would rather have customers
- Peers are often competitors
- Peering agreements require periodic renegotiation

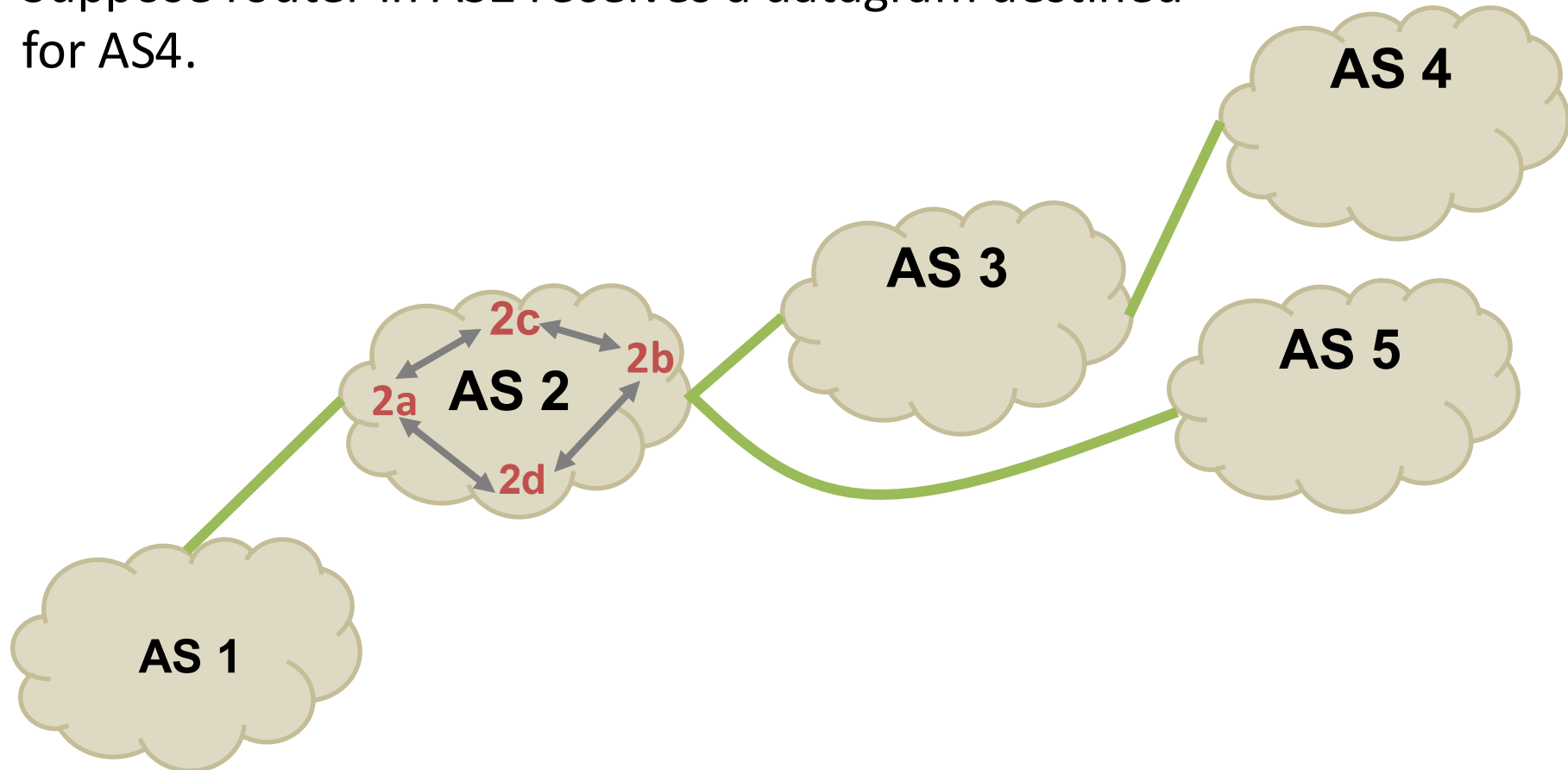
Peering struggles in the ISP world are extremely contentious, agreements are usually confidential

Hierarchical routing: Interconnected ASes



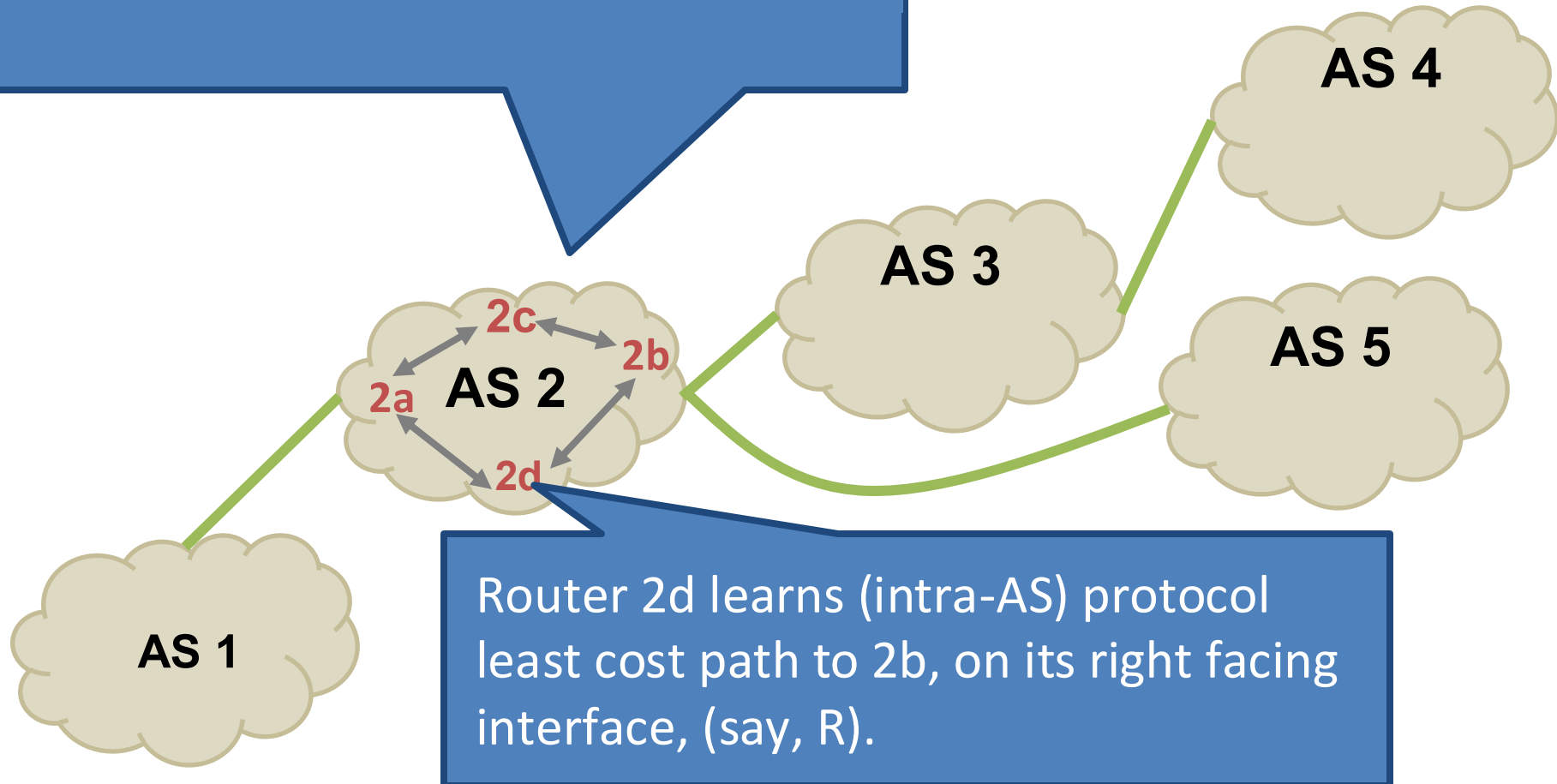
Building the forwarding table in router 2d, for path to AS4

Suppose router in AS2 receives a datagram destined for AS4.



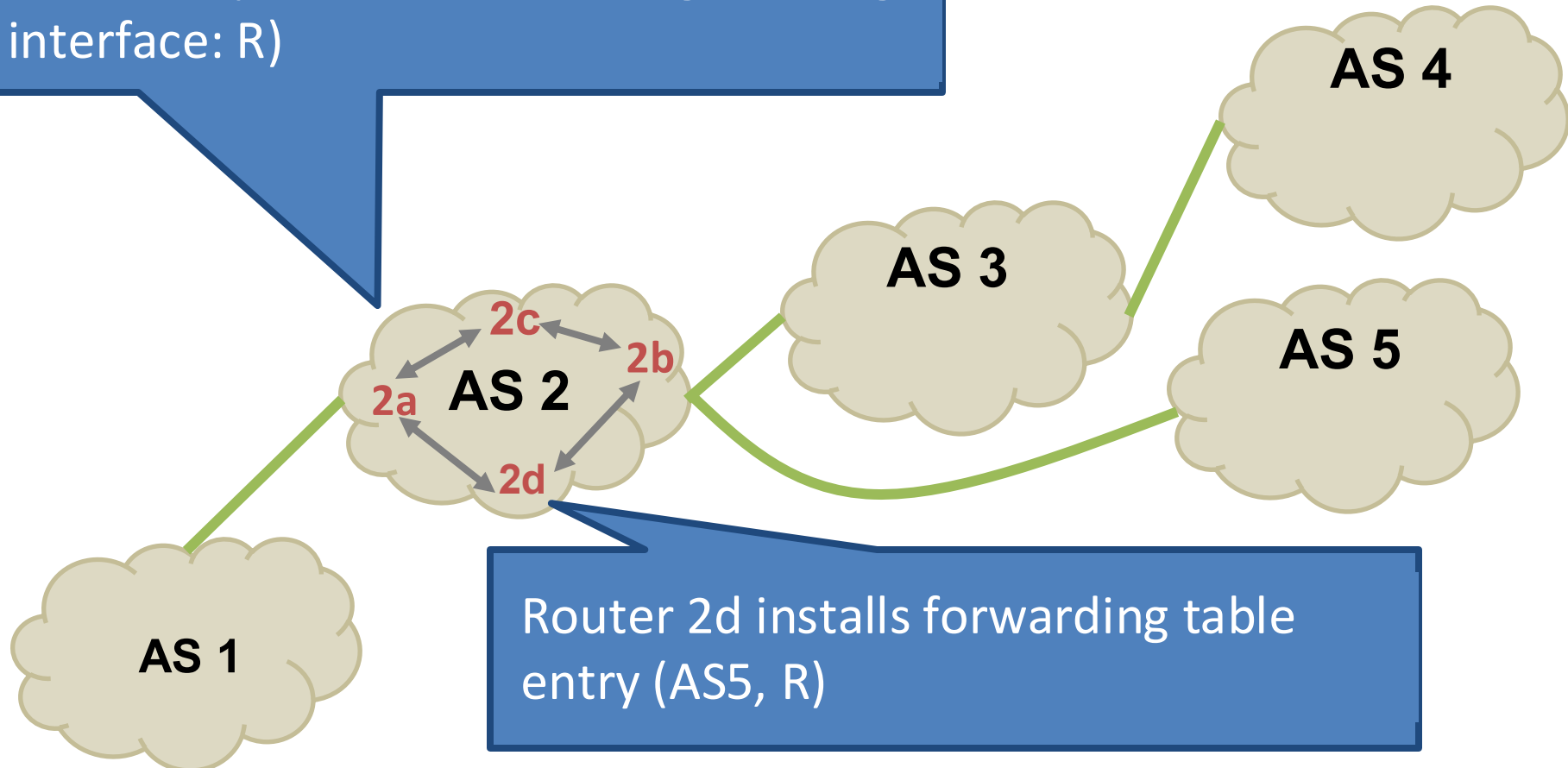
Building the forwarding table in router 2d, for path to AS4

AS2 learns (Inter-AS) protocol that AS4 is reachable through AS3 (via gateway 2b).



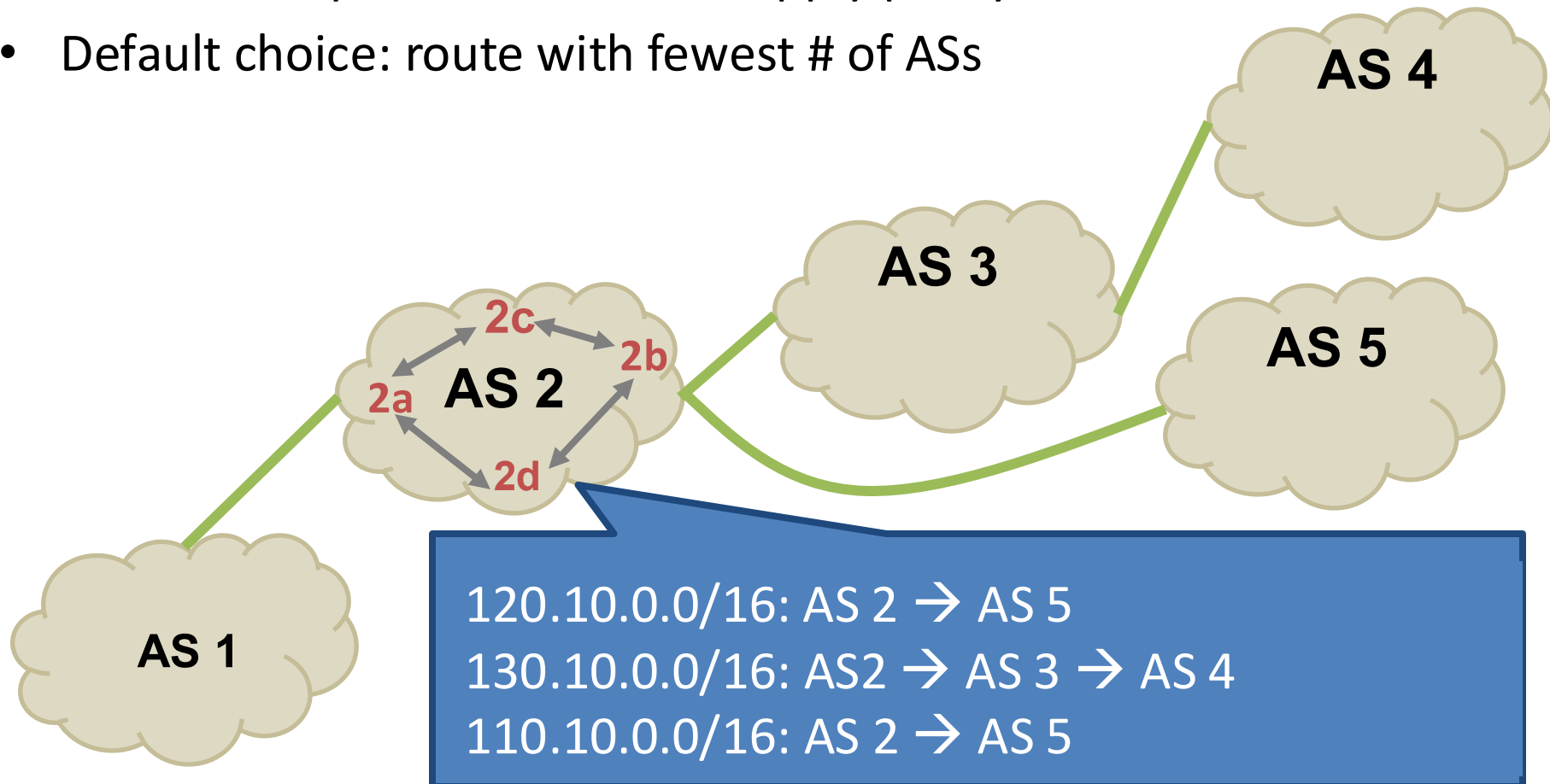
Building the forwarding table in router 2d, for path to AS 5

Router 2d learns (intra-AS) protocol least cost path to 2b (on it's right facing interface: R)



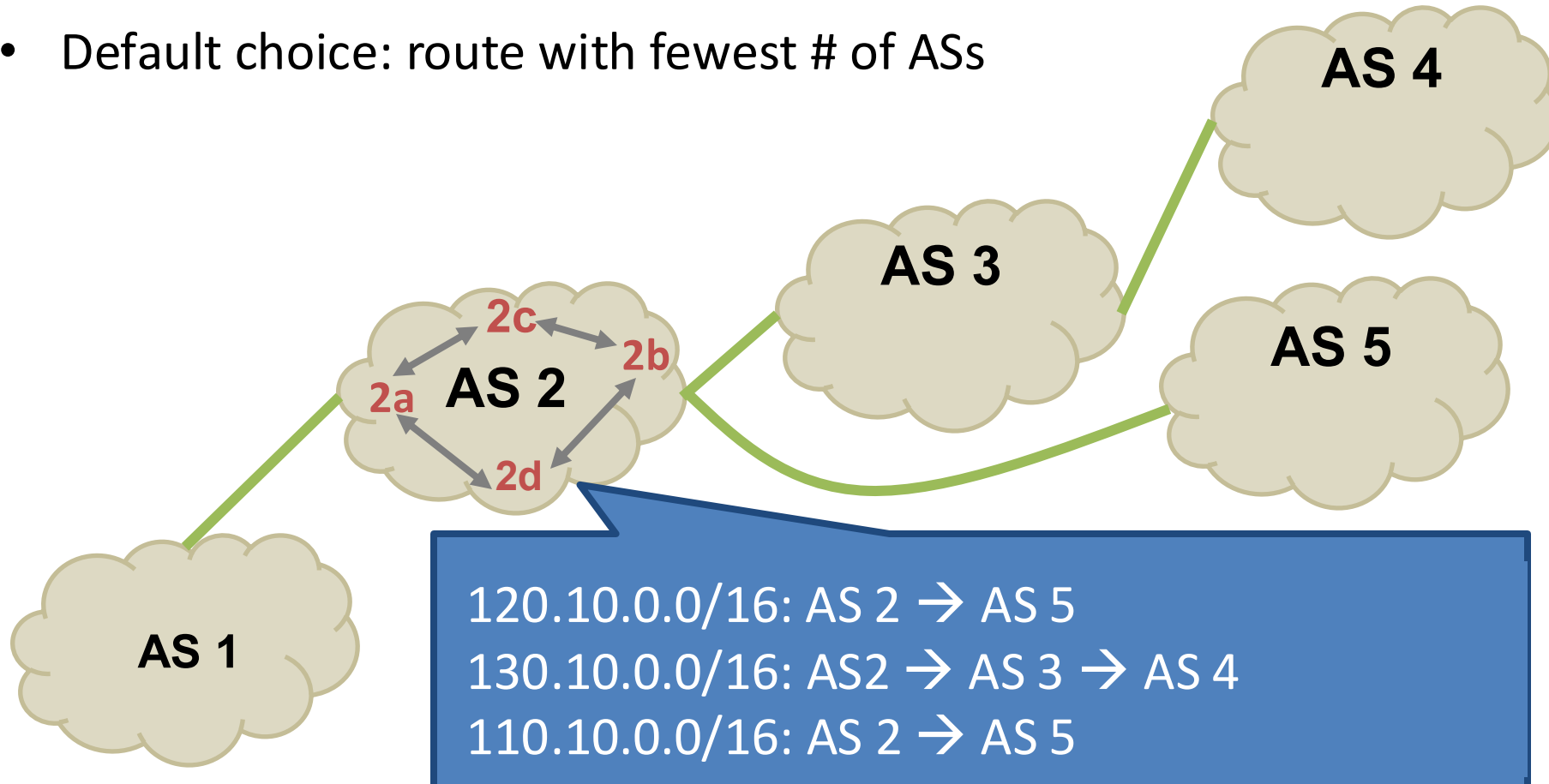
Path Vector Protocol

- AS-path: sequence of ASs a route traverses
 - Like distance vector, plus additional information
- Used for loop detection and to apply policy
- Default choice: route with fewest # of ASs

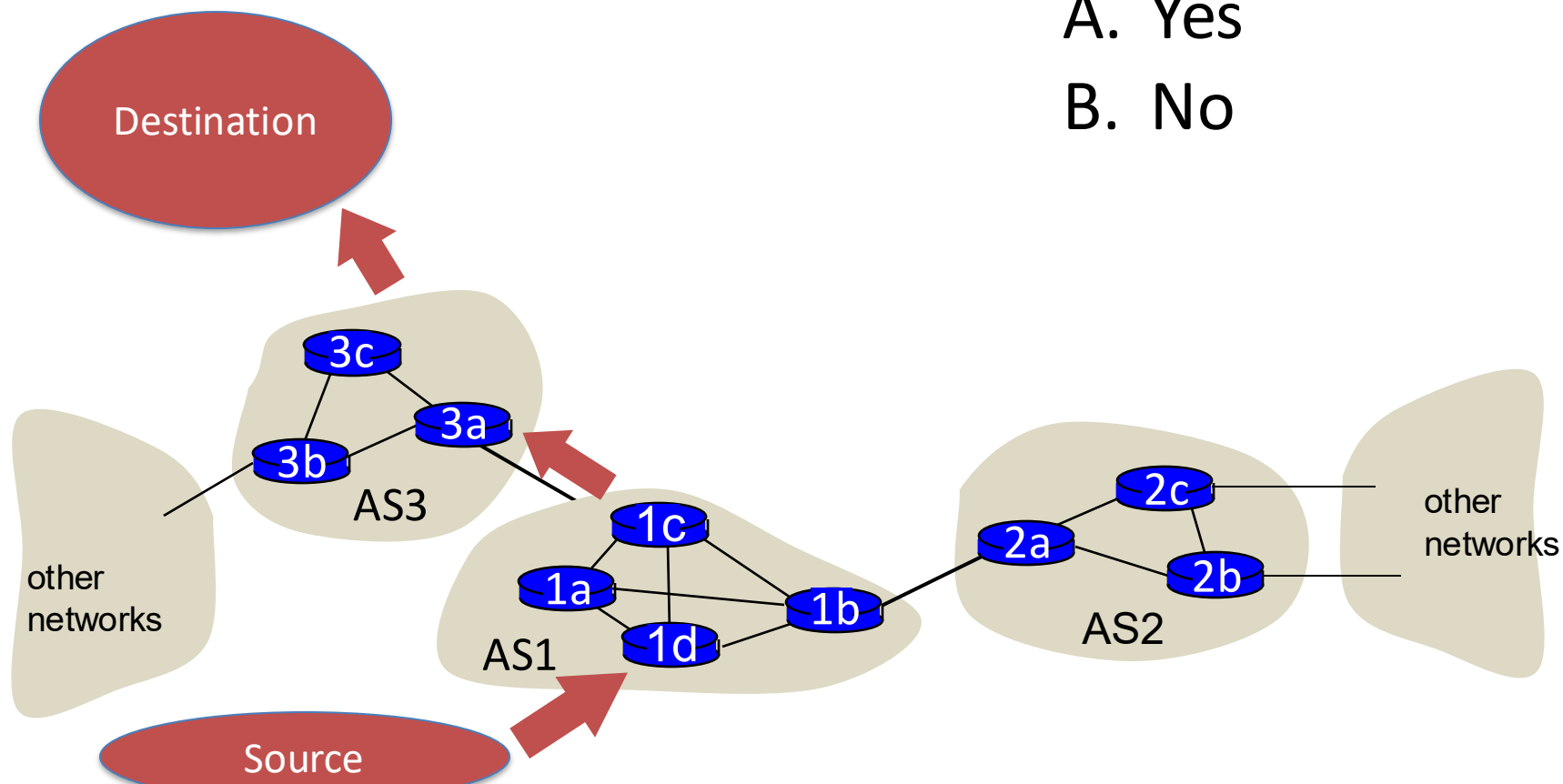


Path Vector Protocol

- AS-path: sequence of ASs a route traverses
 - Like distance vector, plus additional information
- Used for loop detection and to apply policy
- Default choice: route with fewest # of ASs

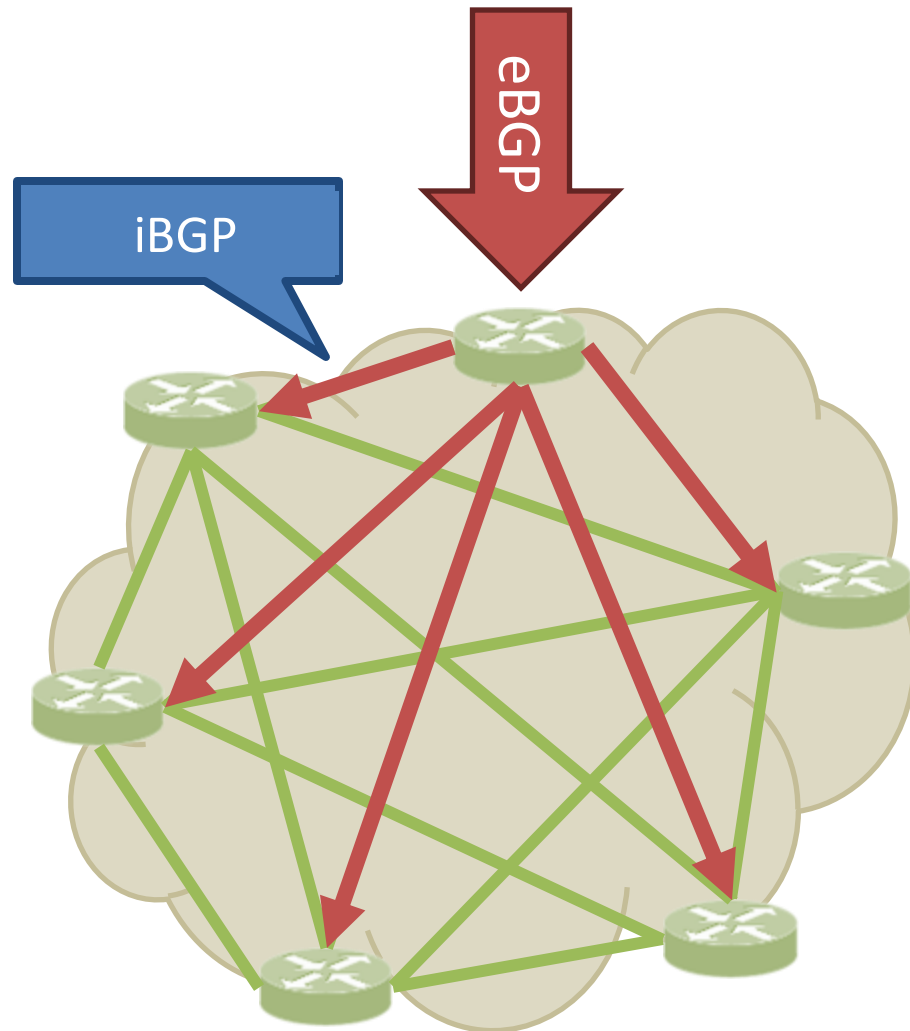


Border routers: exchange AS reachability, Internal routers: exchange intra-AS reachability., Is this sufficient to route from source to destination?



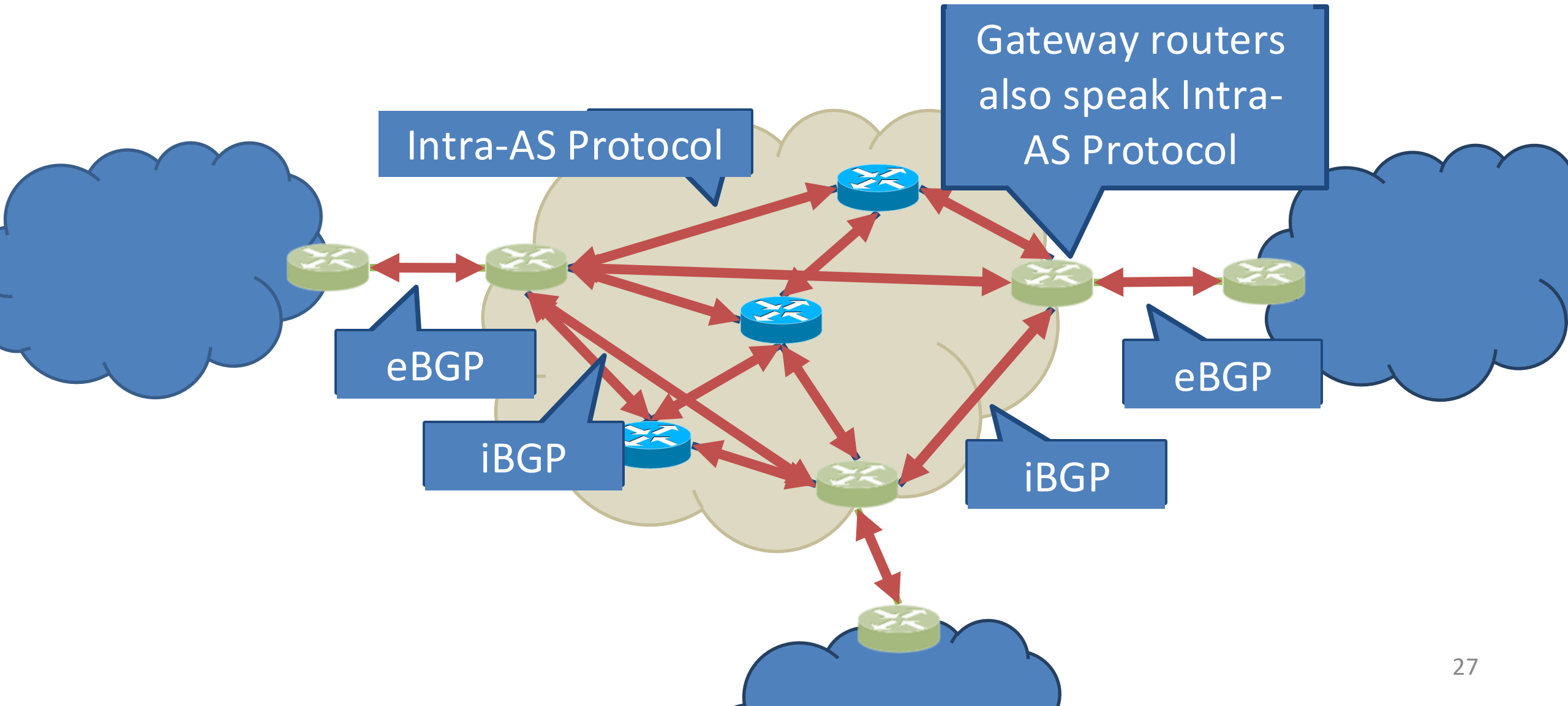
- A. Yes
- B. No

Internet inter-AS routing: BGP



- Question: why do we need iBGP?
 - OSPF does not include BGP policy info
 - Prevents routing loops within the AS
- iBGP updates do not trigger announcements

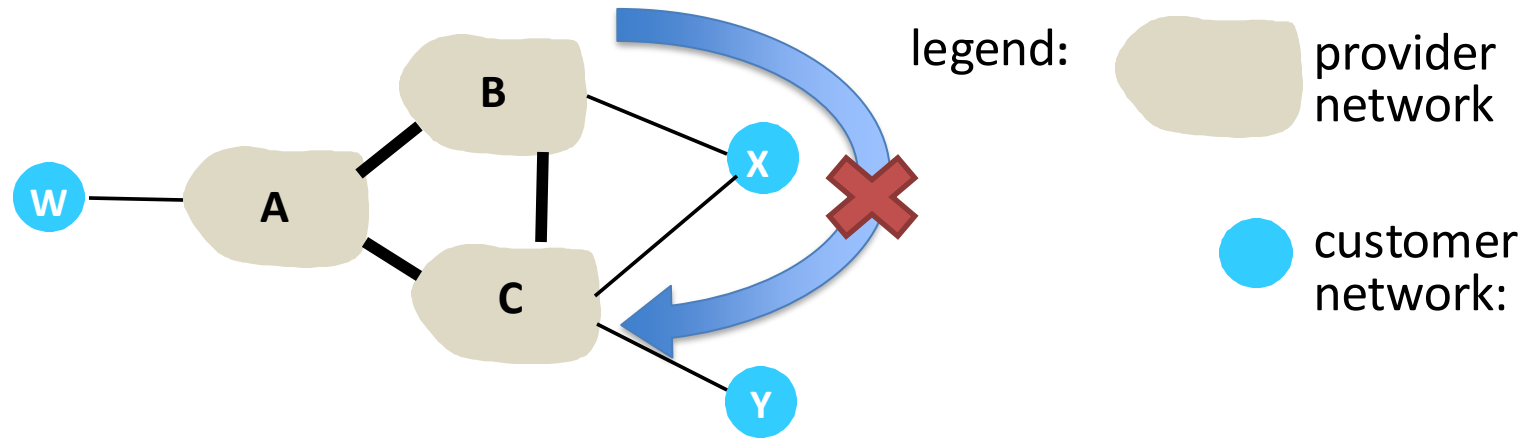
Internet inter-AS routing: BGP



Internet inter-AS routing: BGP

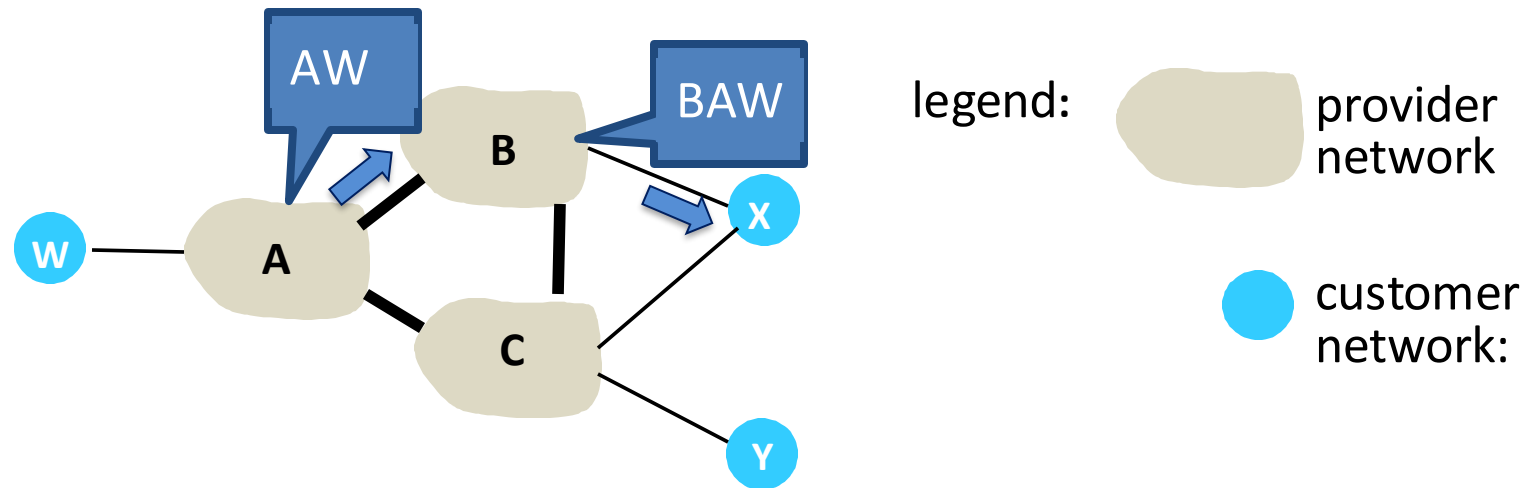
- **BGP (Border Gateway Protocol):**
The de facto inter-domain routing protocol
- BGP provides each AS a means to:
 - **external BGP:** obtain subnet reachability information from neighboring ASs.
 - **internal BGP:** propagate reachability information to all AS-internal routers.
 - determine “good” routes to other networks based on reachability information and policy.
- Allows a subnet to advertise its prefix to the rest of the Internet

BGP routing policy



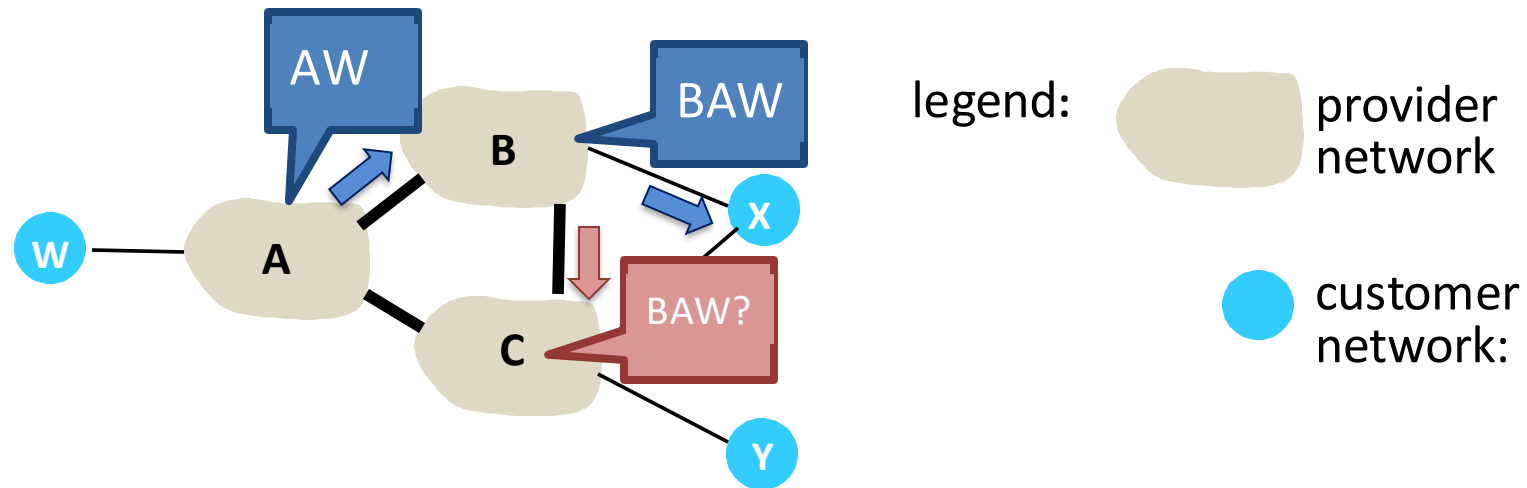
- A,B,C are **provider networks**
- X,W,Y are customers of the providers
- X is **dual-homed**: attached to two networks (B and C)
 - X does not want to route from B via X to C
 - .. so X will not advertise to B a route to C

BGP routing policy



- A advertises path AW to B
- B advertises path BAW to X

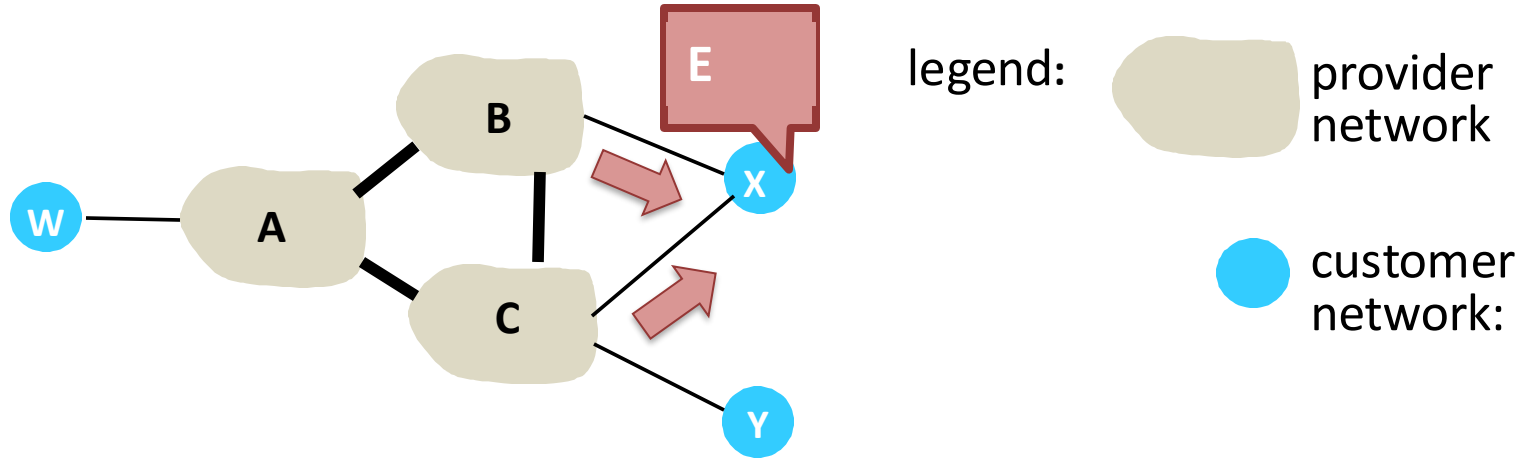
BGP routing policy: Should B advertise path BAW to C?



Should B advertise path BAW to C?

- B gets no “revenue” for routing CBAW since neither W nor C are B’s customers
- B wants to force C to route to w via A
- B wants to route *only* to/from its customers!

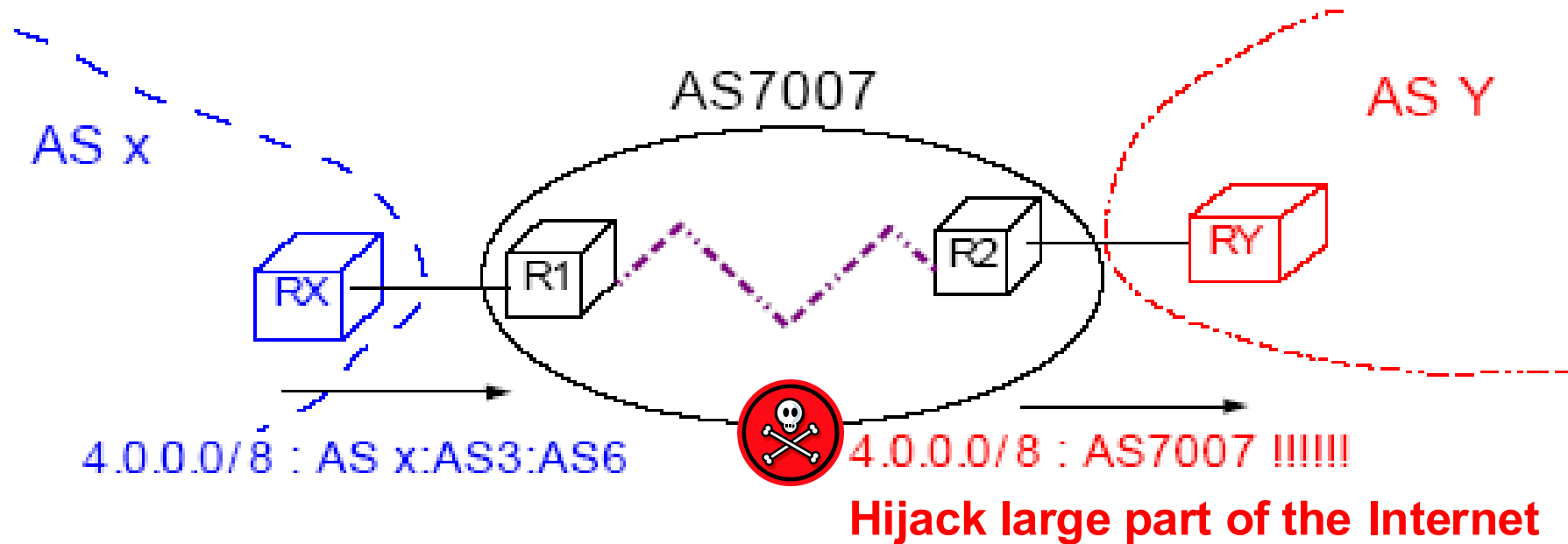
BGP routing policy gone wrong



- x advertises a path to E (that it is not connected to).
- all traffic starts to flow into x from B and C!

Faulty redistribution can be dangerous!

- AS7007 incident (April, 1997):

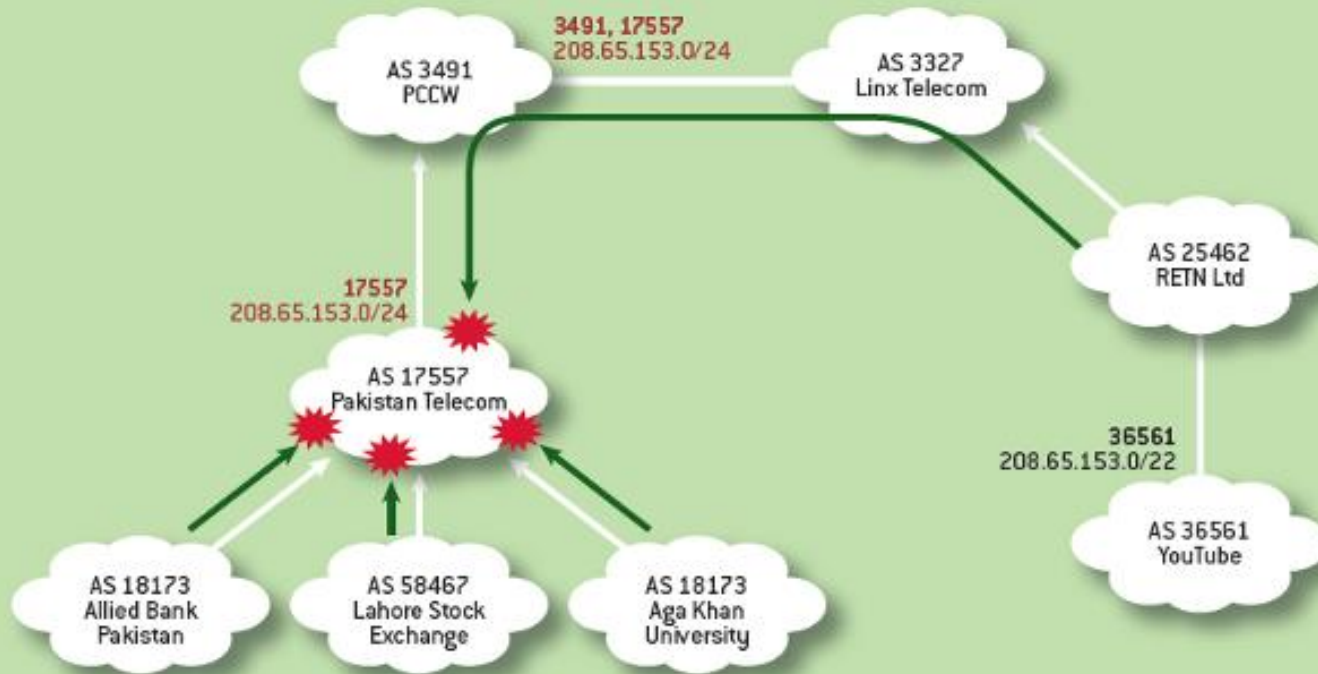


<https://www.kentik.com/blog/a-brief-history-of-the-internets-biggest-bgp-incidents/>

Some “incorrect” route announcements are intentional: BGP Hijacking

FIGURE 3

Pakistan Telecom Hijacks YouTube³⁷



- Pakistan Telecom’s AS 17557 launched a **subprefix hijack** by originating the subprefix 208.65.153.0/24 of YouTube’s prefix 208.65.153.0/22 to its customer ASes in Pakistan
- Traffic destined for YouTube’s servers in AS 36561 would instead be **forwarded to the longer IP prefix** originated by Pakistan Telecom’s AS 17557
- Bogus BGP announcement **leaked out of Pakistan** to large ISP PCCW

Summary

- As we've seen before (DNS), a hierarchy can help manage state storage constraints.
 - intra-AS routing: lots of info about local routes
 - inter-AS routing: less info about far away routes
- BGP: the inter-AS routing protocol for the Internet
 - Decisions often contractual
- BGP advertises AS prefixes, including:
 - entire path of ASes along the way
 - which border router heard the advertisement (Next Hop)

Additional Info:

Inter-Domain Routing Challenges

- BGP4 is the only inter-domain routing protocol currently in use world-wide
- Issues?
 - Lack of security
 - Ease of misconfiguration
 - Poorly understood interaction between local policies
 - Poor convergence
 - Lack of appropriate information hiding
 - Non-determinism
 - Poor overload behavior

Additional Info:

Lots of research into how to fix this

- Security
 - BGPSEC, RPKI
- Misconfigurations, inflexible policy
 - SDN
- Policy Interactions
 - PoiRoot (root cause analysis)
- Convergence
 - Consensus Routing
- Inconsistent behavior
 - LIFEGUARD, among others

Additional Info

Why are these still issues?

- Backward compatibility
- Buy-in / incentives for operators
- Stubbornness

Very similar issues to IPv6 deployment

Additional Info:

Why Network Reliability Remains Hard

- Visibility
 - IP provides no built-in monitoring
 - Economic disincentives to share information publicly
- Control
 - Routing protocols optimize for policy, not reliability
 - Outage affecting your traffic may be caused by distant network
- Detecting, isolating and repairing network problems for Internet paths remains largely a slow, manual process