

# CS 43: Computer Networks

Naming and DNS

September 24, 2025



Slides adapted from Kurose & Ross, Vasanta Chaganti, Kevin Webb

# Where we are

Application: the application (e.g., HTTP, DNS)

Transport: end-to-end connections, reliability

Network: routing

Link (data-link): framing, error detection

Physical: 1's and 0's/bits across a medium  
(copper, the air, fiber)

# Today

- Domain Name System
  - Telephone directory of the Internet
  - Protocol format
  - Caching: Load balancing

# What's the biggest challenge for DNS?

- A. It's old.
- B. The fact that the Internet is global.
- C. The fact that DNS is now critical infrastructure.
- D. The sheer number of name lookups happening at any given time.
- E. How and when the name -> IP address mapping should change.

## In the old days...

- Pre-1982, everyone downloads a “hosts.txt” file from SRI
- Pre-1998, Jon Postel, researcher at USC, runs the **Internet Assigned Numbers Authority (IANA)**
  - RFCs 882 & 883 in 1983
  - RFCs 1034 & 1035 in 1987



Emailed 8/12 root DNS servers, asked change to his authority. They did.

<http://www.wired.com/wiredenterprise/2012/10/joe-postel/>


# Who should control DNS?

- A. US government
- B. UN / International government
- C. Private corporation
- D. Someone else

## Since 1998...

- Control of Internet Assigned Numbers Authority (IANA) transferred to **Internet Corporation for Assigned Names and Numbers (ICANN)**
  - ICANN is a private non-profit (formerly) blessed by US DOC
  - Global advisory committee for dealing with international issues
  - 2000's: Many efforts for UN control, US resisted
  - 2016: ICANN no longer partnered with DOC

# DNS Root Server and Politics




BIZ & IT TECH SCIENCE POLICY CARS GAMING & CULTURE

REVOKE .RU DOMAIN, UKRAINE SAYS —

## Ukraine asks ICANN to revoke Russian domains and shut down DNS root servers

Expert: Cutting DNS links would harm Russian people but have little impact on gov't.

JON BRODKIN - 3/2/2022, 2:33 PM



BIZ & IT TECH SCIENCE POLICY CARS GAMING & CULTURE

THE INTERNET MUST KEEP WORKING —

## ICANN won't revoke Russian Internet domains, says effect would be "devastating"

ICANN's mission: Make sure the Internet works "regardless of the provocations."

JON BRODKIN - 3/4/2022, 1:13 PM



## Uses of DNS

### Hostname to IP address translation

- Reverse lookup: IP address to hostname translation

### Host name aliasing: other DNS names for a host

- Alias hostnames point to canonical hostname

### Email: look up domain's mail server by domain name

# Different DNS Mappings

1-1 mapping  
between domain  
name and IP addr

www.cs.cornell.edu  
maps to  
132.236.207.20

Multiple domain  
names maps to the  
same IP addr

eecs.mit.edu and  
cs.mit.edu both  
map to 18.62.1.6

Single domain  
name maps to  
multiple IP addrs

aol.com and  
www.aol.com map  
to multiple IP addrs

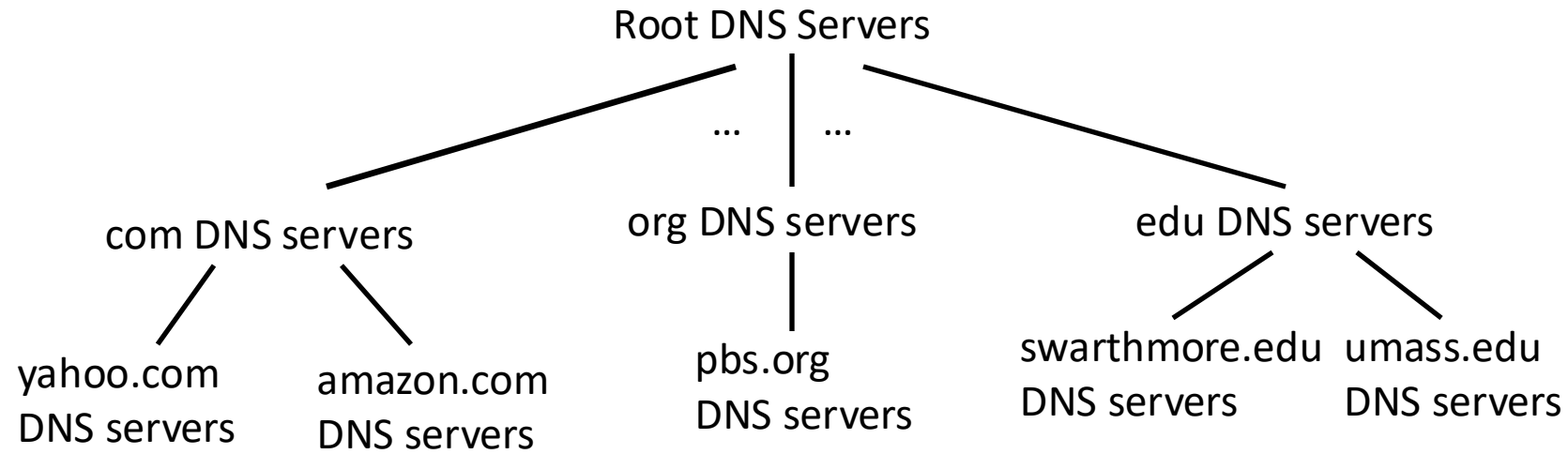
Some valid domain  
names don't map  
to any IP addr

cmcl.cs.cmu.edu

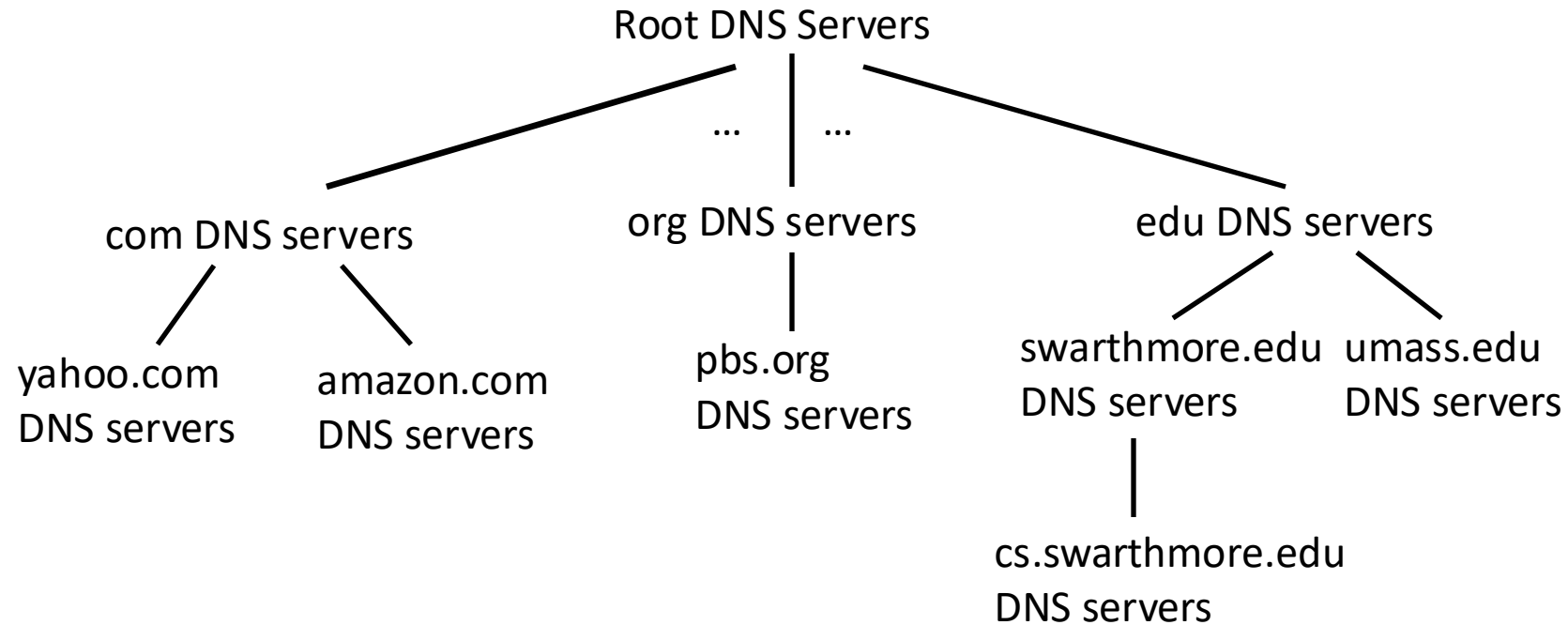
# Domain Name System (DNS)

- Distributed administrative control
  - Hierarchical name space divided into zones
  - Distributed over a collection of DNS servers
- Hierarchy of DNS servers
  - Root servers
  - Top-level domain (TLD) servers
  - Authoritative DNS servers
- Performing the translations
  - Local DNS servers
  - Resolver software

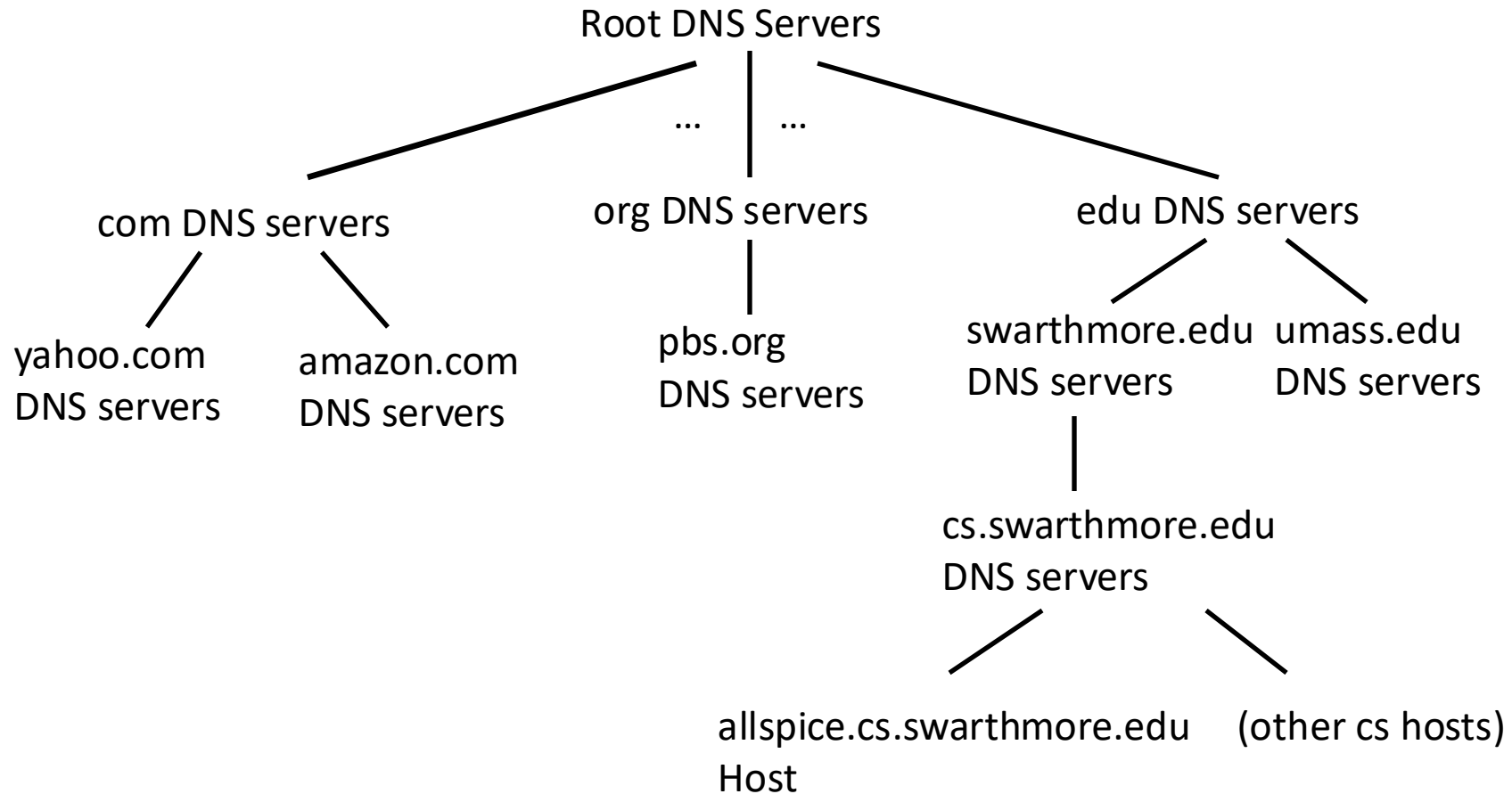
# DNS: a distributed, hierarchical database



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- allspice.cs.swarthmore.edu.

Nameless root,  
Usually implied.

Why do we structure DNS like this? Which of these helps the most?  
Drawbacks?

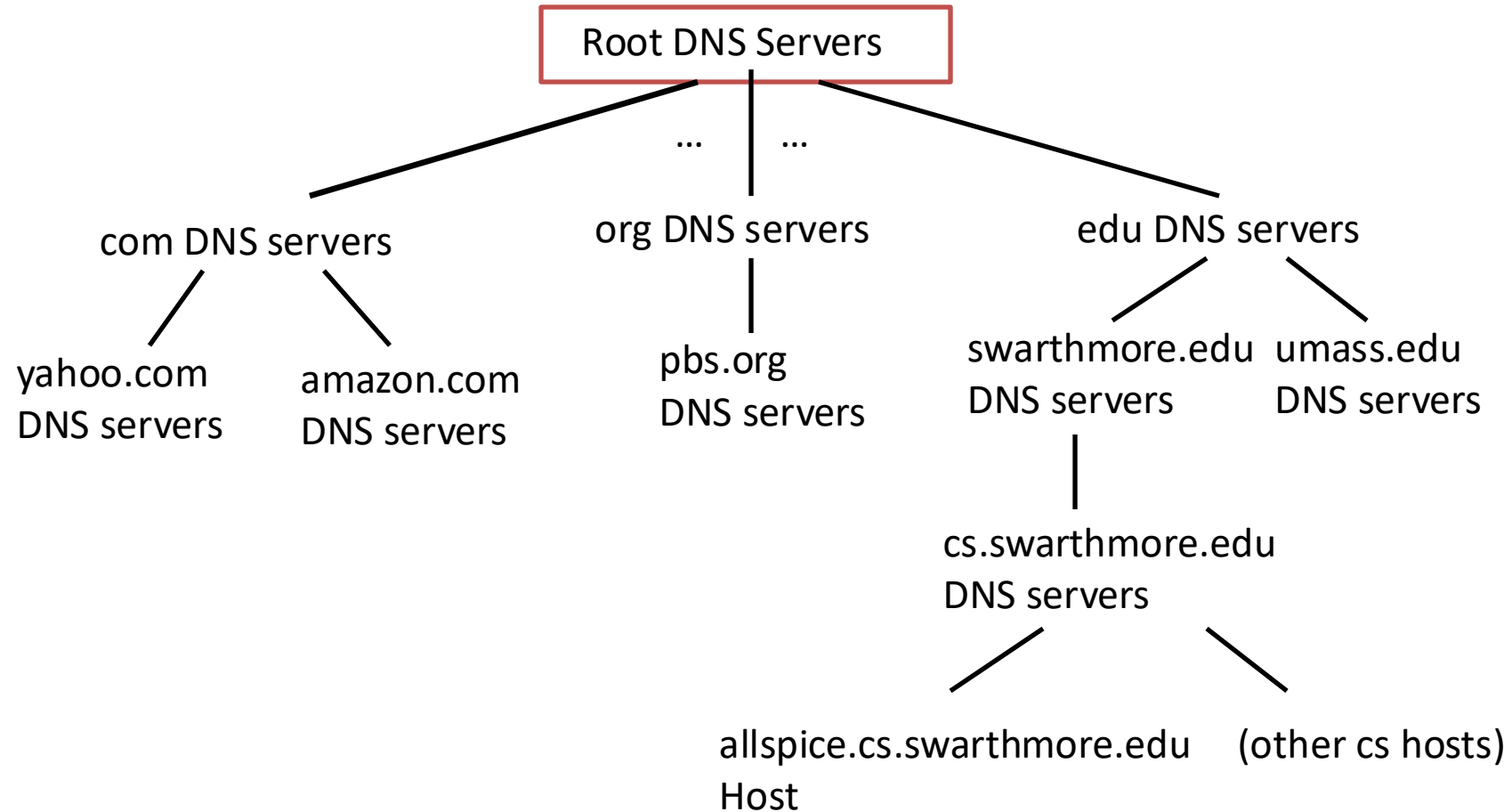
- A. It divides up responsibility among parties.
- B. It improves performance of the system.
- C. It reduces the size of the state that a server needs to store.
- D. Some other reason.

Why do we structure DNS like this? Which of these helps the most? Drawbacks?

- A. It divides up responsibility among parties.
- B. It improves performance of the system overall but individual end hosts (assuming no caching) have a look-up overhead of traversing the hierarchy .
- C. It reduces the size of the state that a server needs to store.
- D. Some other reason.

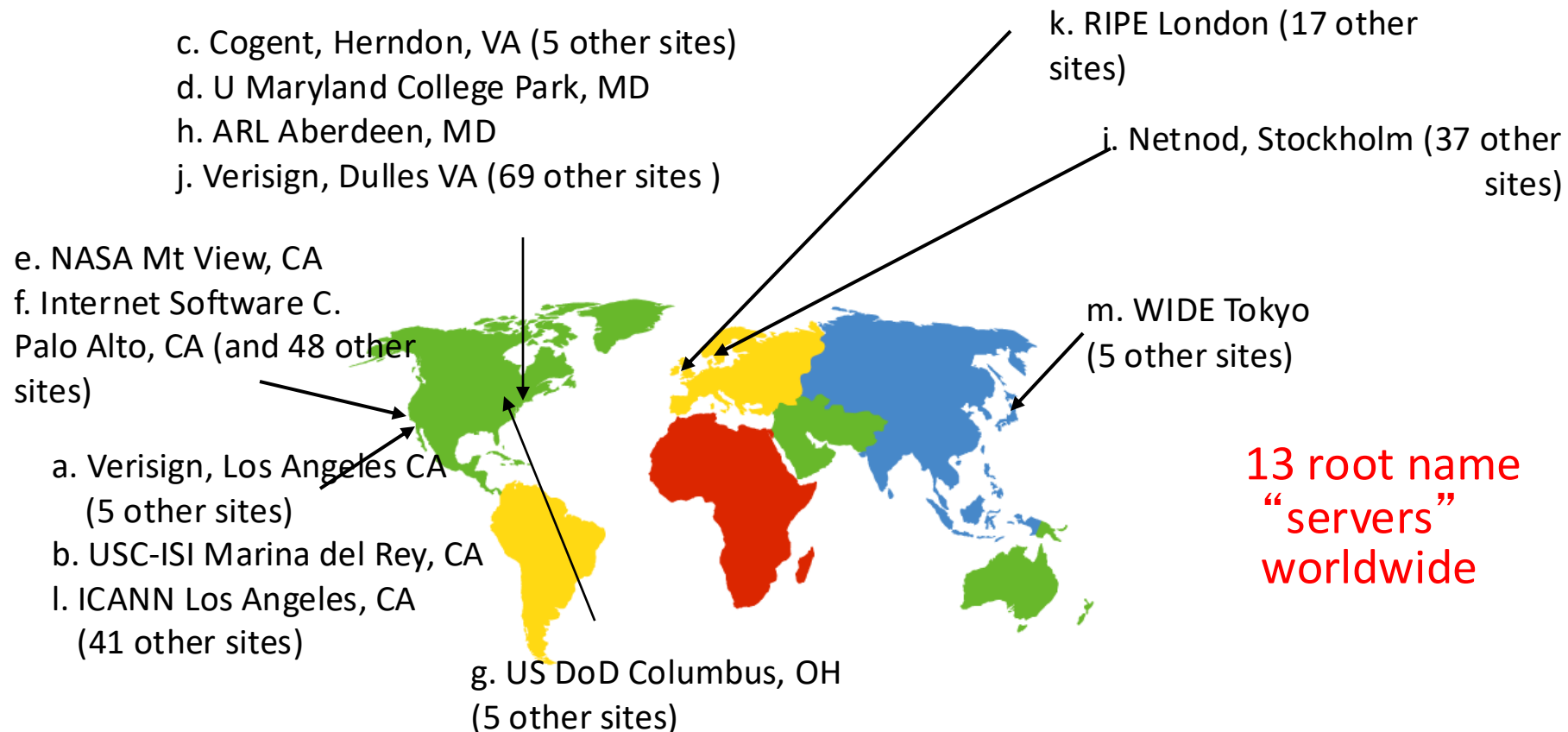


# DNS: a distributed, hierarchical database



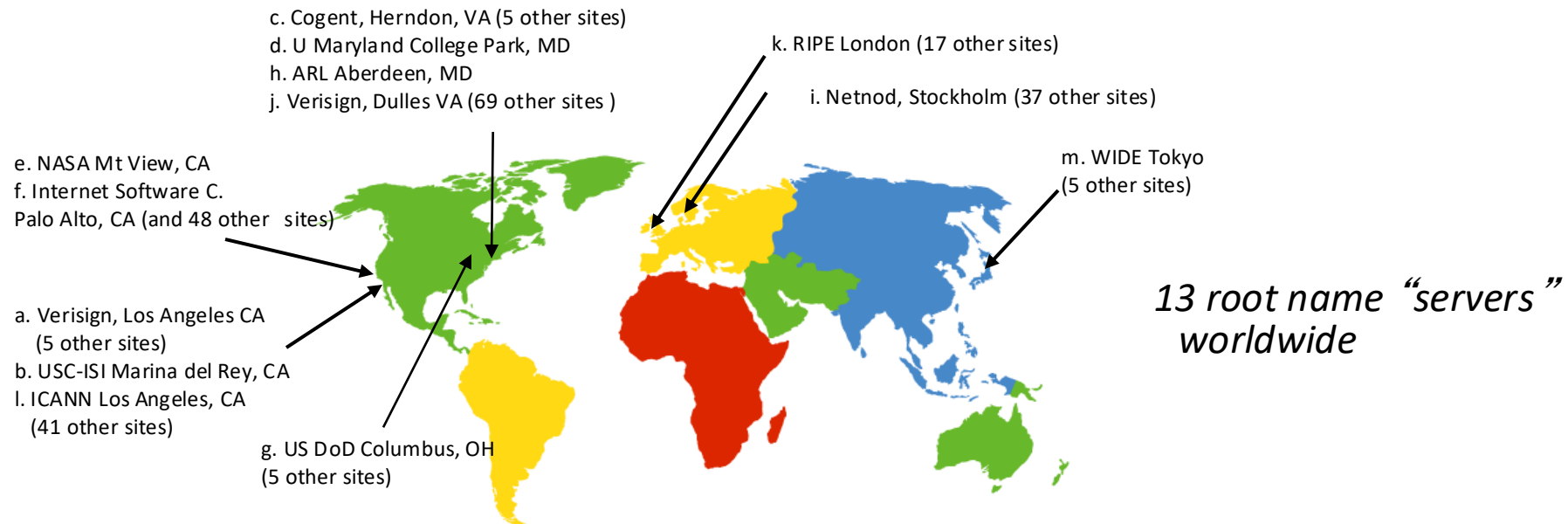
# DNS: Root Name Servers

- Root name server:
  - Knows how to find top-level domains (.com, .edu, .gov, etc.)
  - How often does the location of a TLD change?



# DNS: Root Name Servers

- Root name server:
  - Knows how to find top-level domains (.com, .edu, .gov, etc.)
  - How often does the location of a TLD change?
  - approx. 1700 total root servers
  - Significant amount of traffic is not legitimate



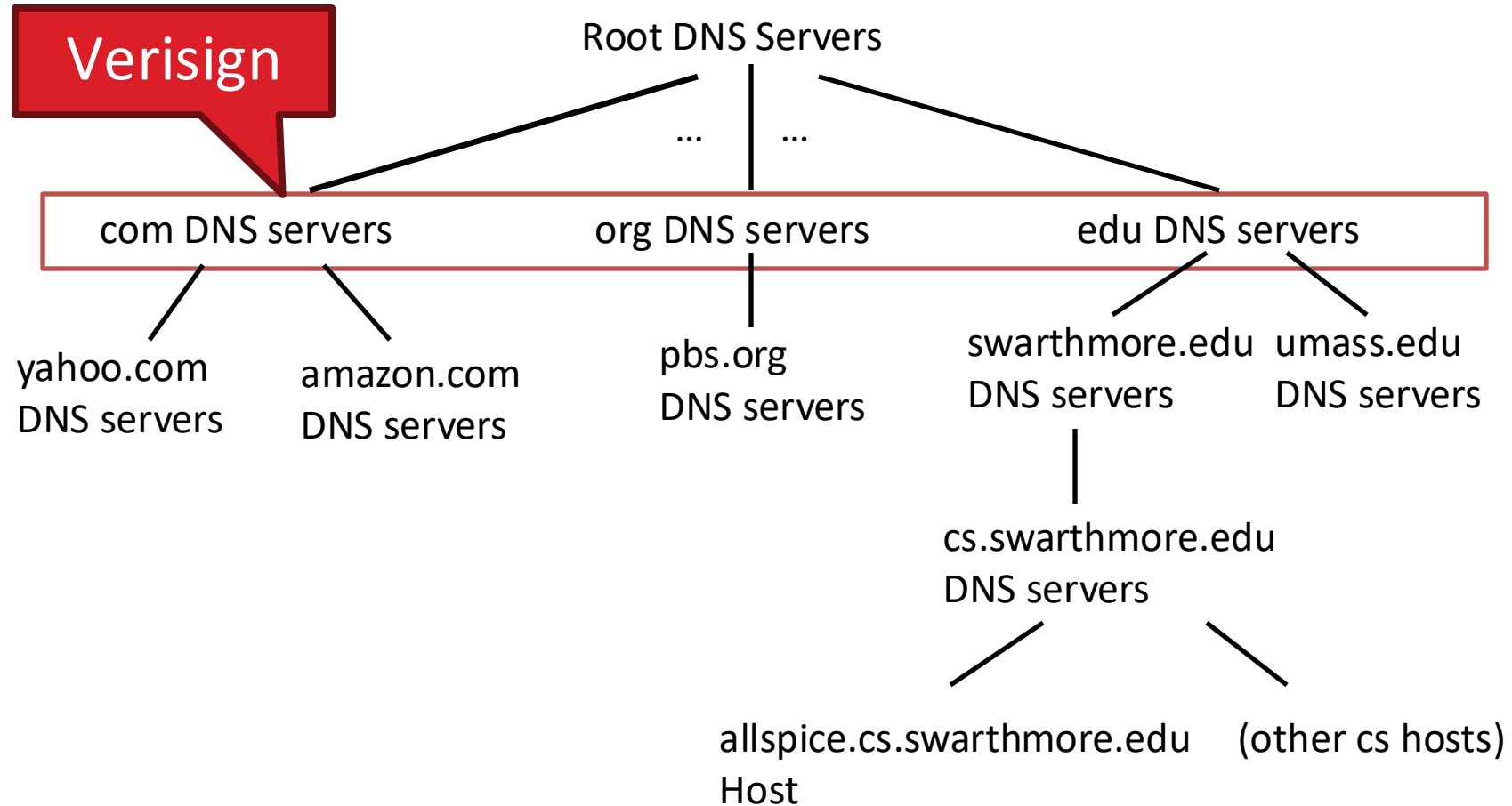
# Root Name Servers

- Responsible for the Root Zone File

com.	172800 IN	NS	a.gtld-servers.net.
com.	172800 IN	NS	b.gtld-servers.net.
com.	172800 IN	NS	c.gtld-servers.net.

- In practice, most systems cache this information
- Lists the TLDs and who controls them
- ~272KB in size

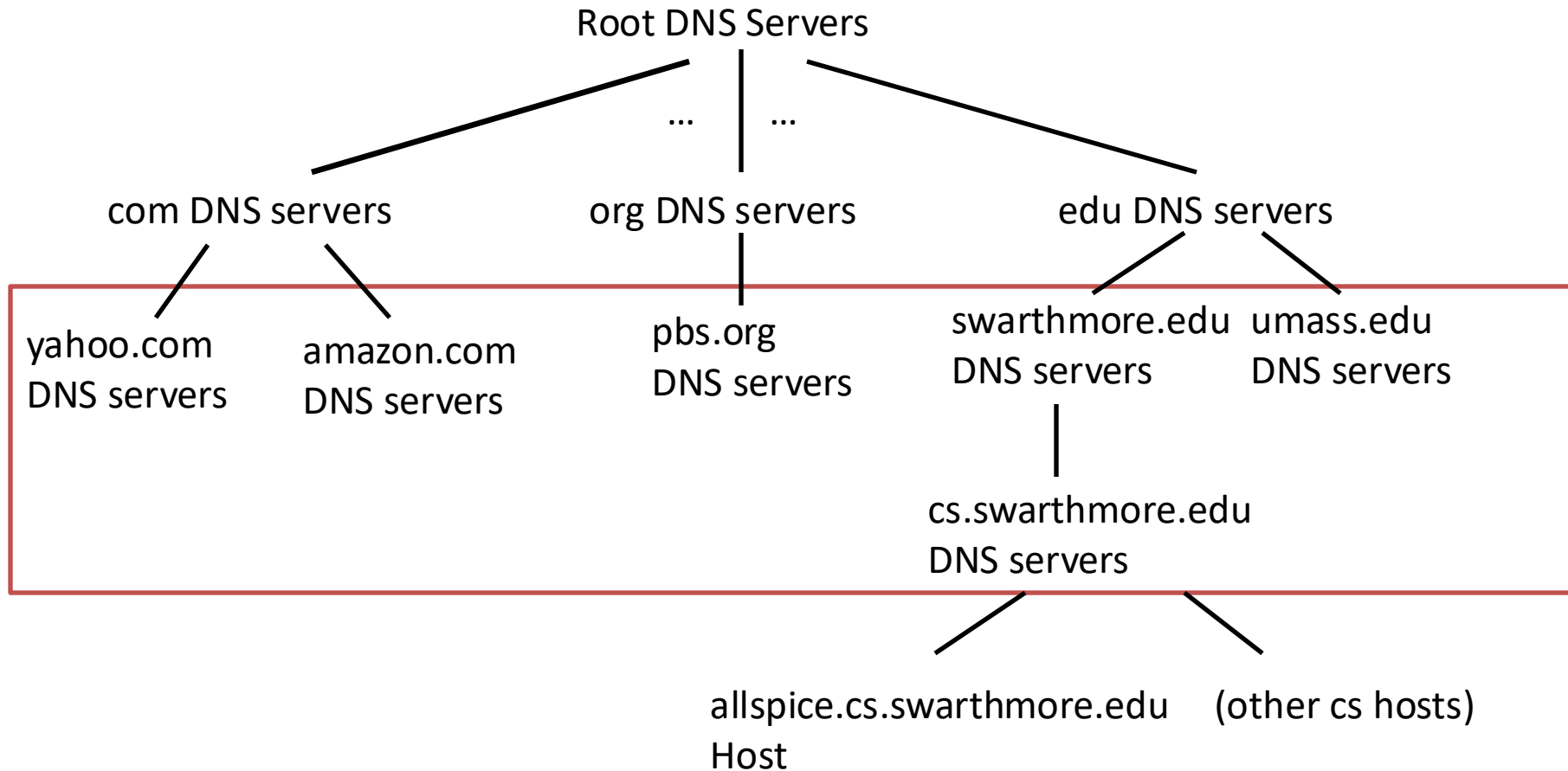
# Top Level Domain (TLD) servers



# Top Level Domain (TLD) servers

- who maintains the servers?:
  - Verisign: .com, .net
  - Educause: .edu (Verisign backend)
  - local governments or companies
- Responsible for:
  - com, org, net, edu, gov, aero, jobs, museums,
  - all top-level country domains, e.g.: uk, fr, de, ca, jp, etc

# Authoritative Servers



# Authoritative Servers

## Authoritative DNS servers:

- Organization's own DNS server(s), providing authoritative hostname to IP mappings for organization's named hosts
- Can be maintained by organization or service provider, easily changing entries
- Often, but not always, acts as organization's local name server (for responding to look-ups)



# Resolution Process

- End host wants to look up a name, who should it contact?
  - It could traverse the hierarchy, starting at a root
  - More efficient for ISP to provide a local server
- ISP's local server for handling queries not necessarily a part of the pictured hierarchy

# Local DNS Name Server

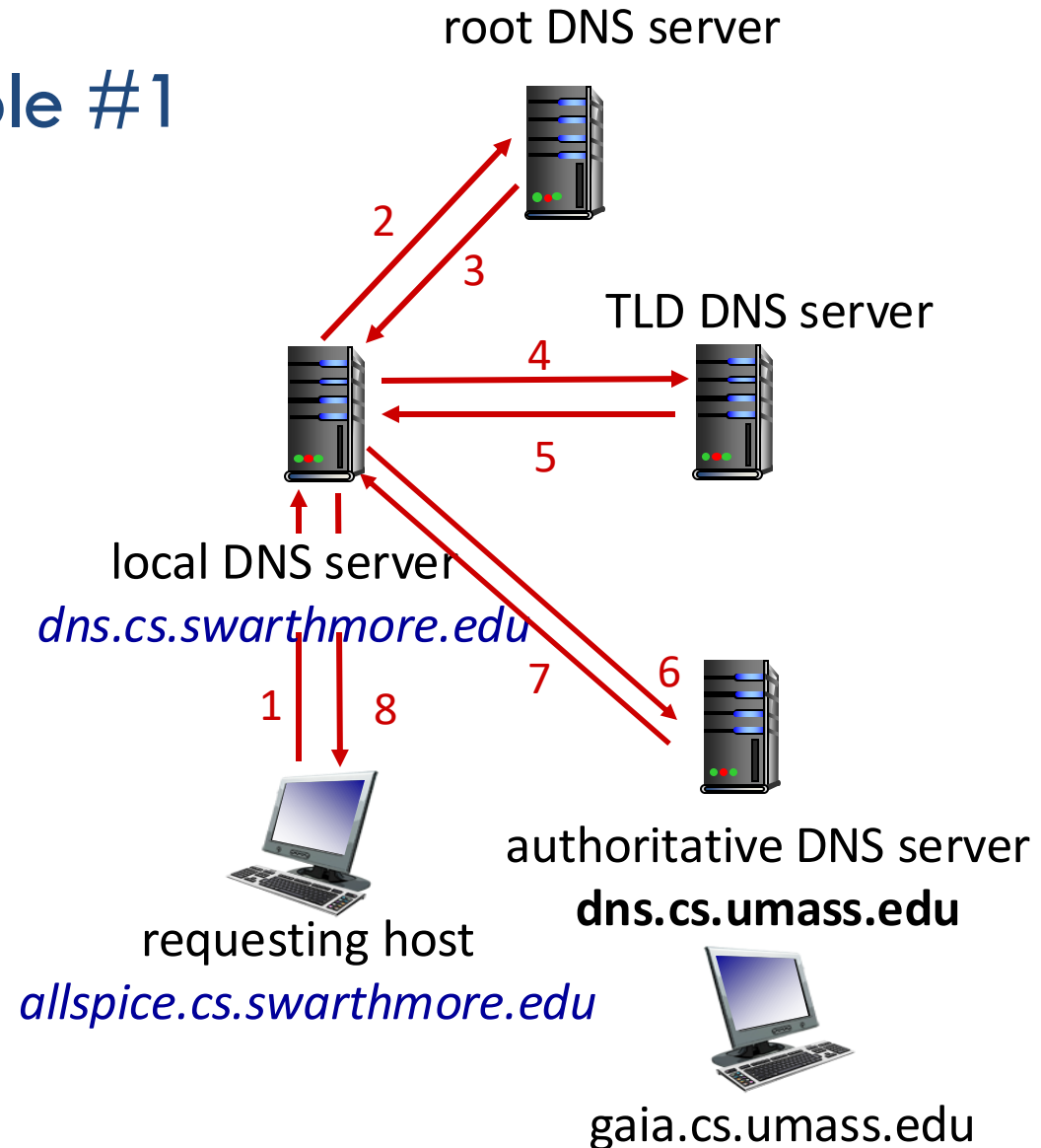
- Each ISP (residential ISP, company, university) has (at least) one
  - also called “default name server”
- When host makes DNS query, query is sent to its local DNS server
  - has local cache of recent name-to-address translation pairs (but may be out of date!)
  - acts as proxy, forwards query into hierarchy

# DNS name resolution example #1

- allspice wants IP address for gaia.cs.umass.edu

## iterative query:

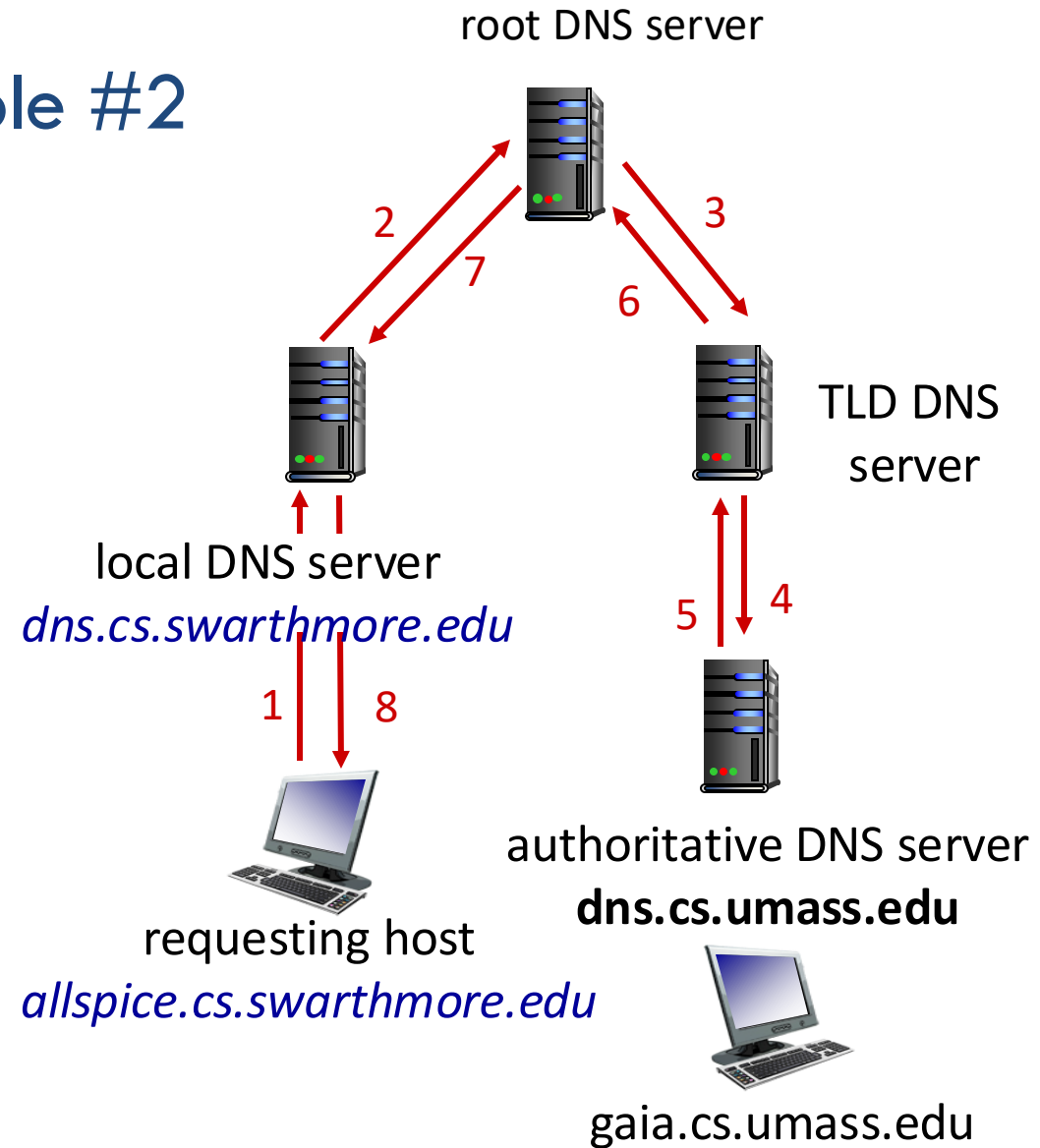
- contacted server replies with name of server to contact
- “I don’t know this name, but ask this server”



# DNS name resolution example #2

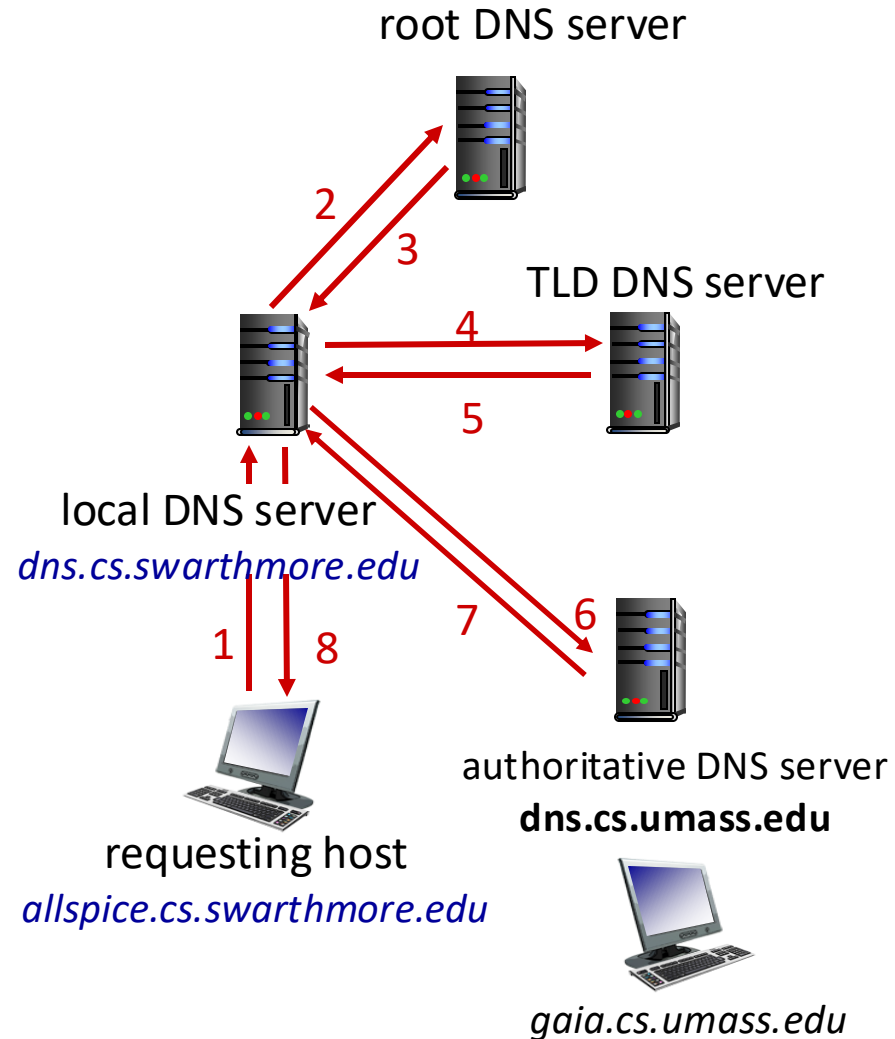
## recursive query:

- each server asks the next one, in a chain

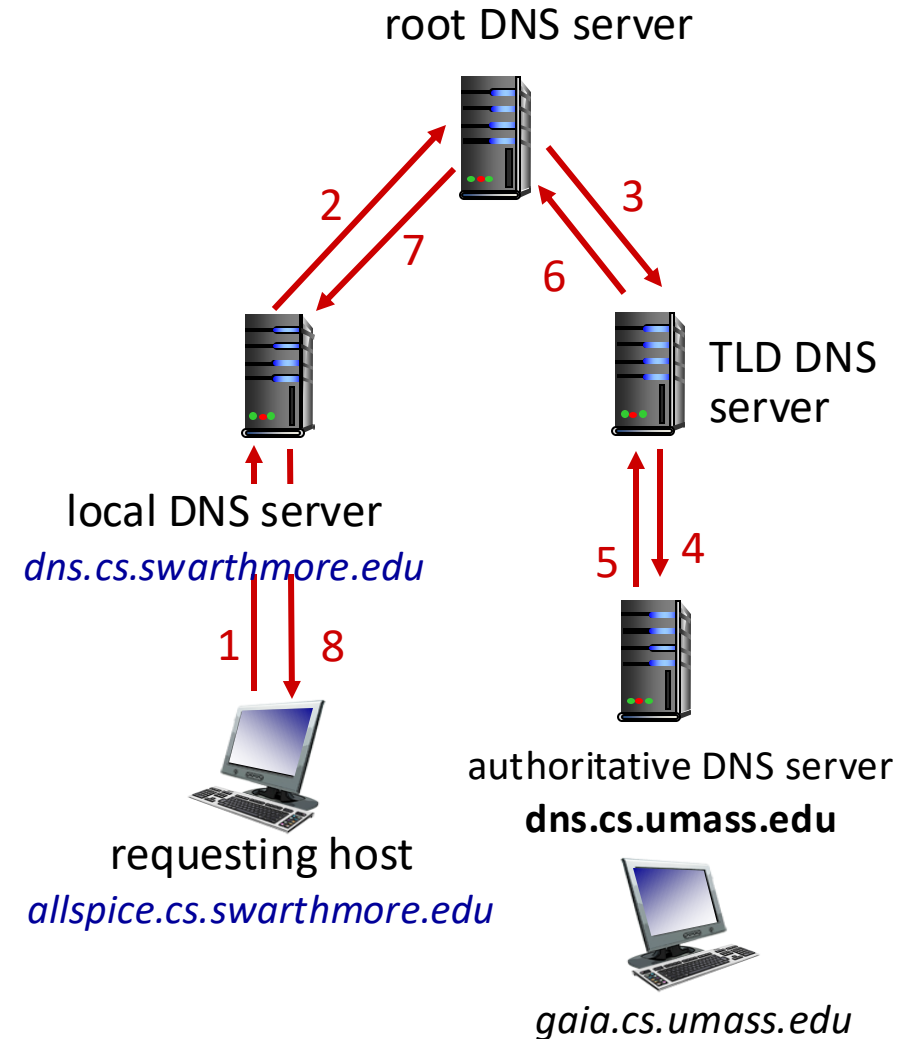


# Which would you use? Why?

## A. Iterative



## B. Recursive

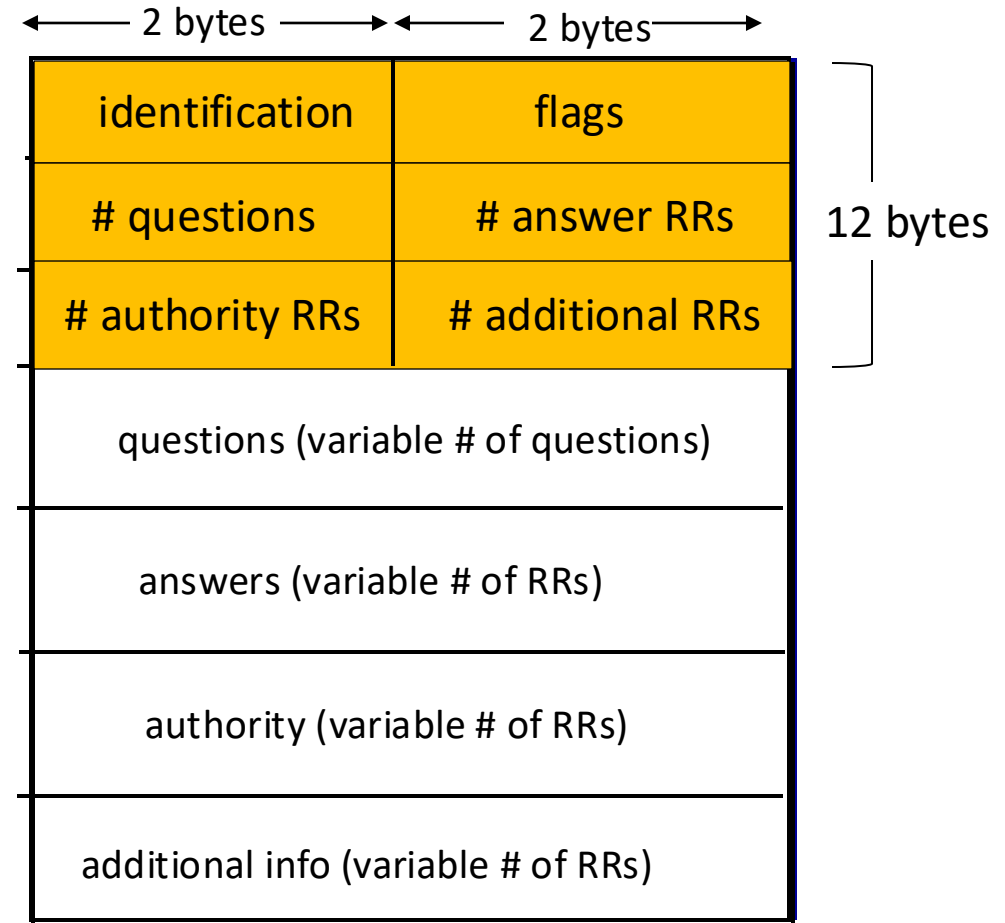


# DNS protocol, messages

- **query** and **reply** messages, both with same **message format**

## Message header

- **identification**: 16 bit id for query, reply to query uses same id.
- **flags**: recursion, query/reply
- # Resource Records to follow



# DNS protocol, messages

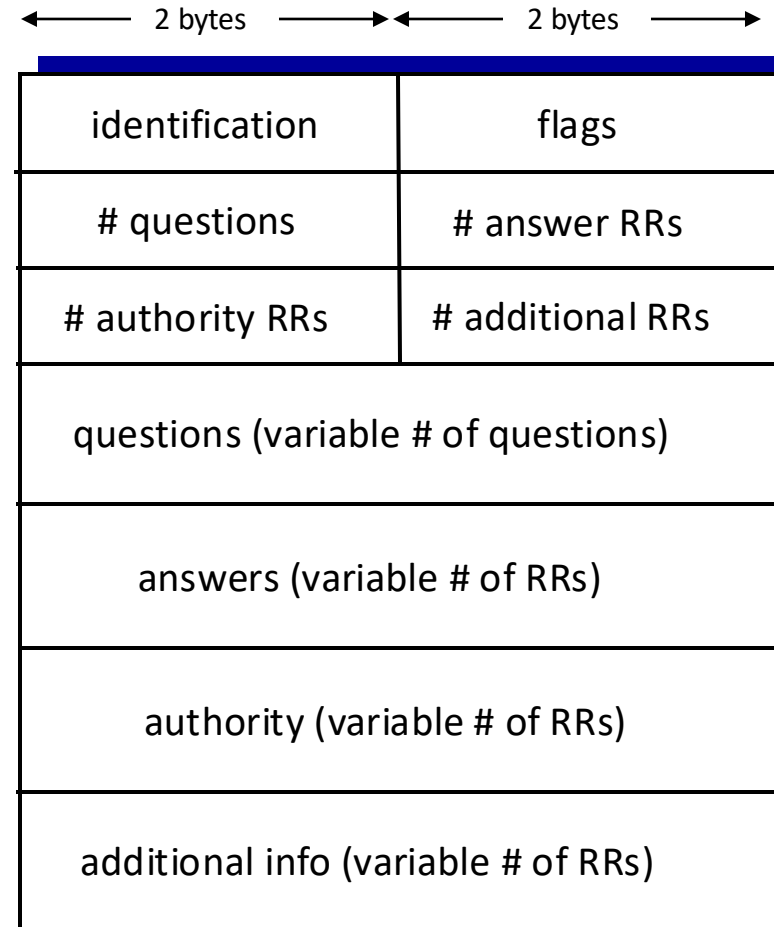
- query and reply messages, both with same message format!

## Binary Protocol!

- Delimiters: pre-defined lengths/field
- Names: <len><name>

## Sent via UDP (User Datagram Protocol)

- No connection established
- Not reliable!



# DNS Types

RR format: (name, value, type, ttl)

- Type = A / AAAA
  - Name = domain name
  - Value = IP address
  - A is IPv4, AAAA is IPv6
- Type = NS
  - Name = partial domain
  - Value = name of DNS server for this domain
  - “Go send your query to this other server”

Query

Name: cs.swarthmore.edu  
Type: A

Resp.

Name: cs.swarthmore.edu  
Value: 130.58.68.9

Query

Name: cs.swarthmore.edu  
Type: NS

Resp.

Name: cs.swarthmore.edu  
Value: dns.cs.swarthmore.edu



# DNS Types, Continued

RR format: (name, value, type, ttl)

- Type = CNAME
  - Name = hostname
  - Value = canonical hostname
  - Useful for aliasing
  - CDNs use this
- Type = MX
  - Name = domain in email address
  - Value = canonical name of mail server

Query

Name: [foo.mysite.com](http://foo.mysite.com)  
Type: CNAME

Resp.

Name: [foo.mysite.com](http://foo.mysite.com)  
Value: [bar.mysite.com](http://bar.mysite.com)

Query

Name: cs.umass.edu  
Type: MX

Resp.

Name: cs.umass.edu  
Value: barramail.cs.umass.edu.

# DNS Records

**DNS:** distributed DB storing resource records (**RR**)

RR format: (**name**, **value**, **type**, **ttl**)

## type=A

- **name** is hostname
- **value** is IP address

## type=NS

- **name** is domain (e.g., foo.com)
- **value** is hostname of authoritative name server for this domain

## type=CNAME

- **name** is alias name for some “canonical” (the real) name
- **www.ibm.com** is really servereast.backup2.ibm.com
- **value** is canonical name

## type=MX

- **value** is name of mailserver associated with name

# Example: iterative query using dig()

```
dig . ns
```

```
dig +norec demo.cs.swarthmore.edu @a.root-servers.net
```

```
dig +norec demo.cs.swarthmore.edu @a.edu-servers.net
```

```
dig +norec demo.cs.swarthmore.edu @ibext.its.swarthmore.edu
```

```
demo.cs.swarthmore.edu. 259200 IN A 130.58.68.26
```

*How many answers*  
*Time to live in seconds*  
*How many additional records?*

```
$ dig @a.root-servers.net www.freebsd.org +norecurse
;; Got answer:
;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 57494
;; QUERY: 1, ANSWER: 0, AUTHORITY: 2, ADDITIONAL: 2

;; QUESTION SECTION:
;www.freebsd.org.      IN A

;; AUTHORITY SECTION:
org.      172800 IN NS b0.org.afilias-nst.org.
org.      172800 IN NS d0.org.afilias-nst.org.

;; ADDITIONAL SECTION:
b0.org.afilias-nst.org. 172800 IN A 199.19.54.1
d0.org.afilias-nst.org. 172800 IN A 199.19.57.1
```

*How many answers*  
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org.      172800 IN NS d0.org.afiliast-nst.org.

;; ADDITIONAL SECTION:
b0.org.afiliast-nst.org. 172800 IN A 199.19.54.1
d0.org.afiliast-nst.org. 172800 IN A 199.19.57.1
```

*Glue records*

*How many answers?*  
*How many additional records?*

 (authoritative for org.)

```
$ dig @199.19.54.1 www.freebsd.org +norecurse
;; Got answer:
;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 39912
;; QUERY: 1, ANSWER: 0, AUTHORITY: 3, ADDITIONAL: 0

;; QUESTION SECTION:
;www.freebsd.org.      IN  A

;; AUTHORITY SECTION:
freebsd.org.          86400 IN  NS  ns1.isc-sns.net.
freebsd.org.          86400 IN  NS  ns2.isc-sns.com.
freebsd.org.          86400 IN  NS  ns3.isc-sns.info.
```

*How many answers?  
How many additional records?*

 (authoritative for org.)

```
$ dig @199.19.54.1 www.freebsd.org +norecurse
;; Got answer:
;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 39912
;; QUERY: 1, ANSWER: 0, AUTHORITY: 3, ADDITIONAL: 0

;; QUESTION SECTION:
;www.freebsd.org.      IN  A

;; AUTHORITY SECTION:
freebsd.org.          86400 IN  NS  ns1.isc-sns.net.
freebsd.org.          86400 IN  NS  ns2.isc-sns.com.
freebsd.org.          86400 IN  NS  ns3.isc-sns.info.
```

 (authoritative for freebsd.org.)

```
$ dig @ns1.isc-sns.net www.freebsd.org +norecurse
```

```
;; Got answer:
```

```
;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 17037
```

```
;; QUERY: 1, ANSWER: 1, AUTHORITY: 3, ADDITIONAL: 3
```

```
;; QUESTION SECTION:
```

```
;www.freebsd.org.      IN  A
```

*How many answers?*

*How many authoritative records?*

*How many additional records?*

```
;; ANSWER SECTION:
```

```
www.freebsd.org. 3600  IN  A   69.147.83.33
```

```
;; AUTHORITY SECTION:
```

```
freebsd.org.      3600  IN  NS  ns2.isc-sns.com.
```

```
freebsd.org.      3600  IN  NS  ns1.isc-sns.net.
```

```
freebsd.org.      3600  IN  NS  ns3.isc-sns.info.
```

```
;; ADDITIONAL SECTION:
```

```
ns1.isc-sns.net. 3600  IN  A   72.52.71.1
```

```
ns2.isc-sns.com. 3600  IN  A   38.103.2.1
```

```
ns3.isc-sns.info. 3600  IN  A   63.243.194.1
```



 (authoritative for freebsd.org.)  
\$ dig @ns1.isc-sns.net www.freebsd.org +norecurse

;; Got answer:

;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 17037

;; QUERY: 1, ANSWER: 1, AUTHORITY: 3, ADDITIONAL: 3

;; QUESTION SECTION:

;www.freebsd.org. IN A

*How many answers?*

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;; ANSWER SECTION:

www.freebsd.org. 3600 IN A 69.147.83.33

;; AUTHORITY SECTION:

freebsd.org. 3600 IN NS ns2.isc-sns.com.

freebsd.org. 3600 IN NS ns1.isc-sns.net.

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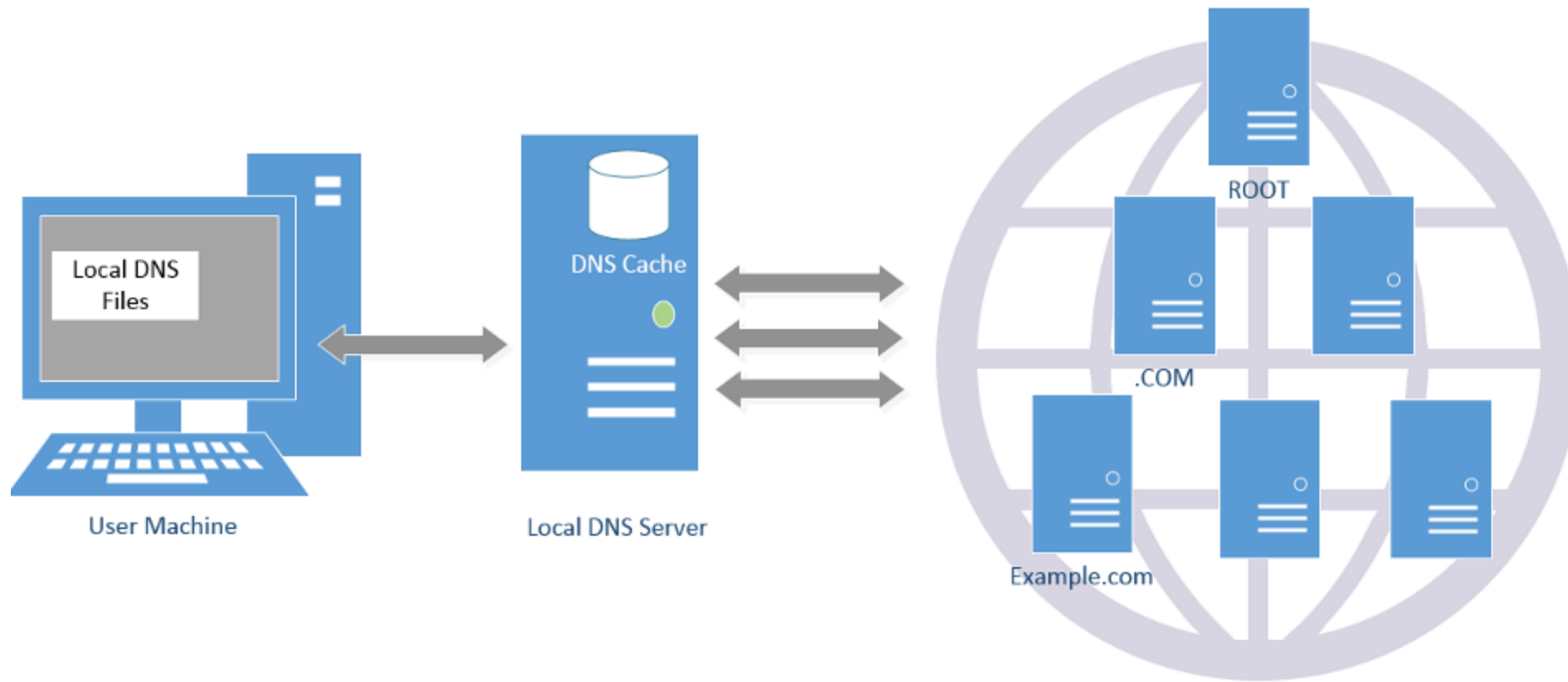
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ns3.isc-sns.info. 3600 IN A 63.243.194.1

# DNS Query Process and Cache

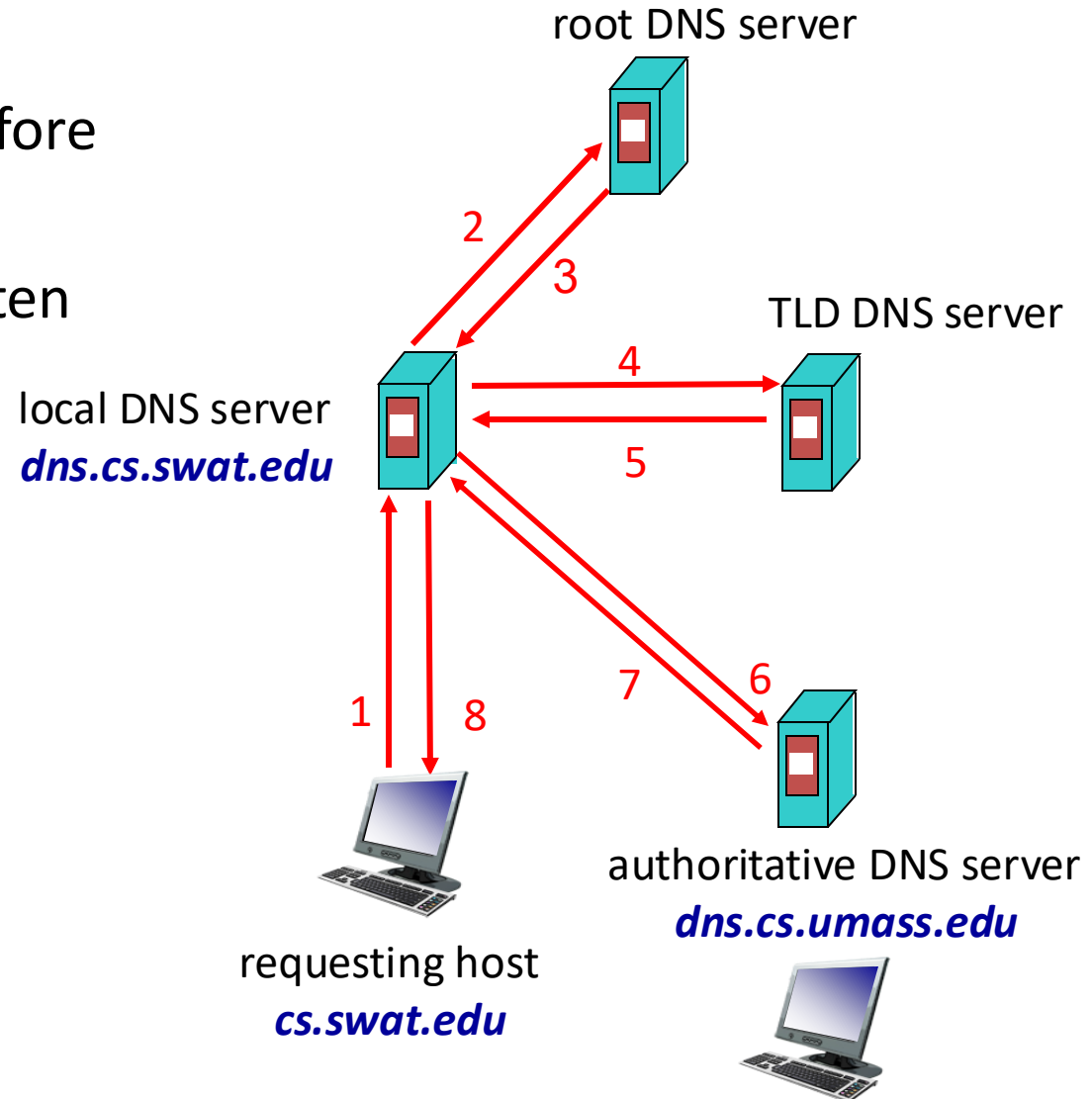


# Caching

- Once (any) name server learns a mapping, it **caches** mapping
  - cache entries timeout (disappear) after some time (TTL: time to live)
  - TLD servers typically cached in local name servers
  - Thus root name servers not often (legitimately) visited

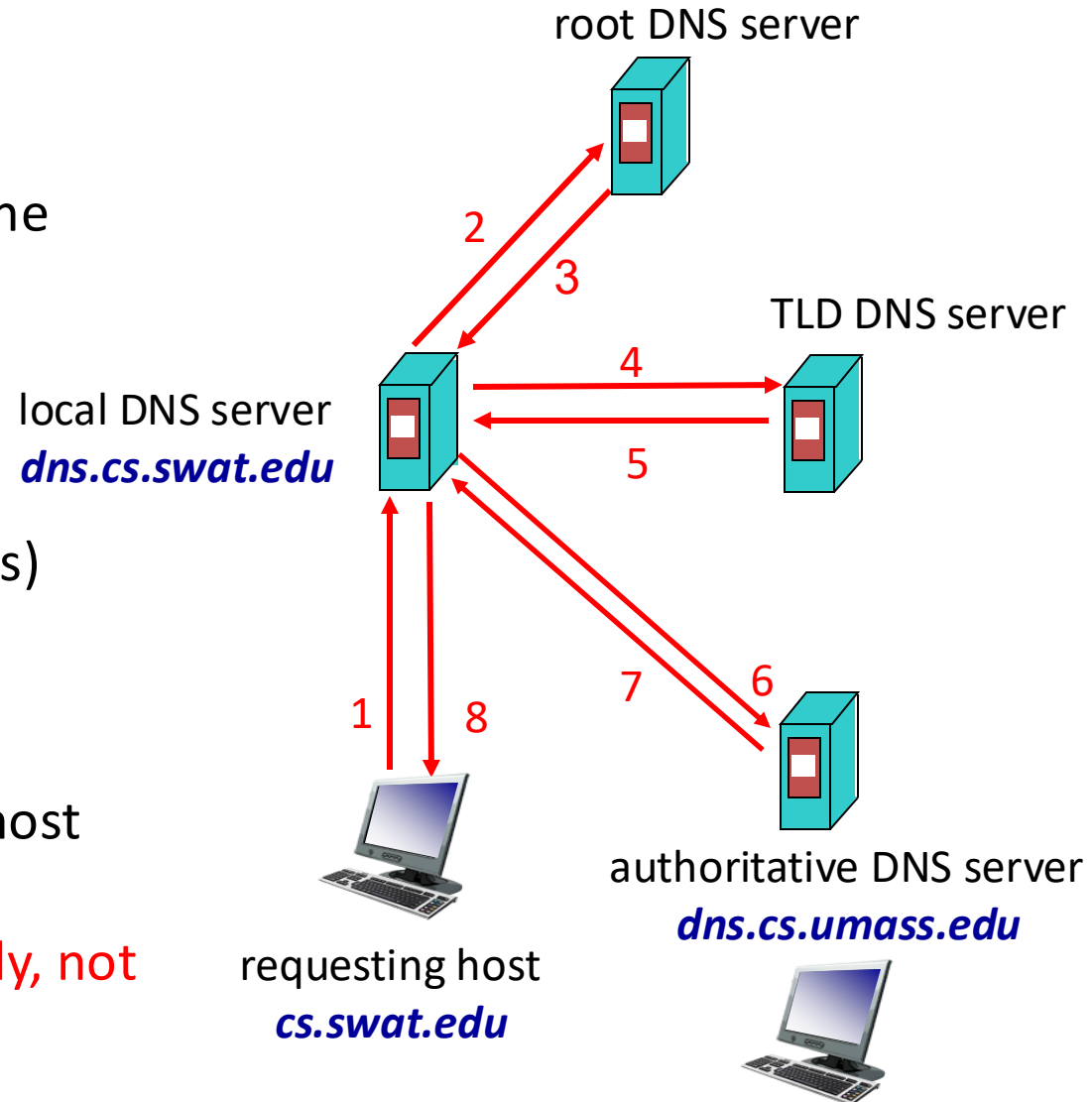
# DNS Caching

- Why cache?
  - apprx. 1 sec latency before starting a download
  - Popular sites visited often
- Where to cache?
  - Local DNS server
  - Browser



# DNS Caching

- When to cache?
  - learn a mapping? cache!
  - any name server can cache
- For how long?
  - until Time To Live (expires)
- What to cache?
  - TLD servers cached – almost never change
  - **Root name servers usually, not visited legitimately**



# Caching

- Once (any) name server learns a mapping, it **caches** mapping
  - cache entries timeout (disappear) after some time (TTL: time to live)
  - TLD servers typically cached in local name servers.
  - Root name servers not often (legitimately) visited
- (+) Subsequent requests need not burden DNS
- (-) Cached entries may be **out-of-date** (best effort!)
  - If host's name or IP address changes, it may not be known Internet-wide until all TTLs expire

## The TTL value should be...

- A. Short, to make sure that changes are accurately reflected
- B. Long, to avoid re-queries of higher-level DNS servers
- C. Something else

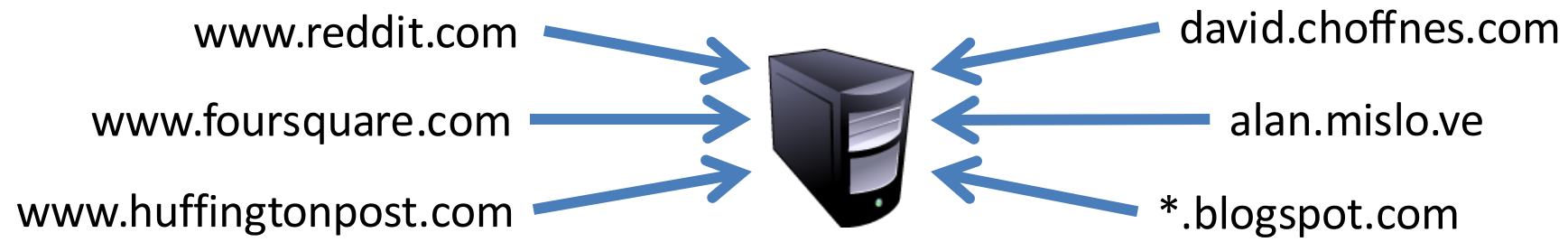
# DNS as Indirection Service

- DNS gives us very powerful capabilities
  - Not only easier for humans to reference machines!
- Changing the IPs of machines becomes trivial
  - e.g. you want to move your web server to a new host
  - Just change the DNS record!

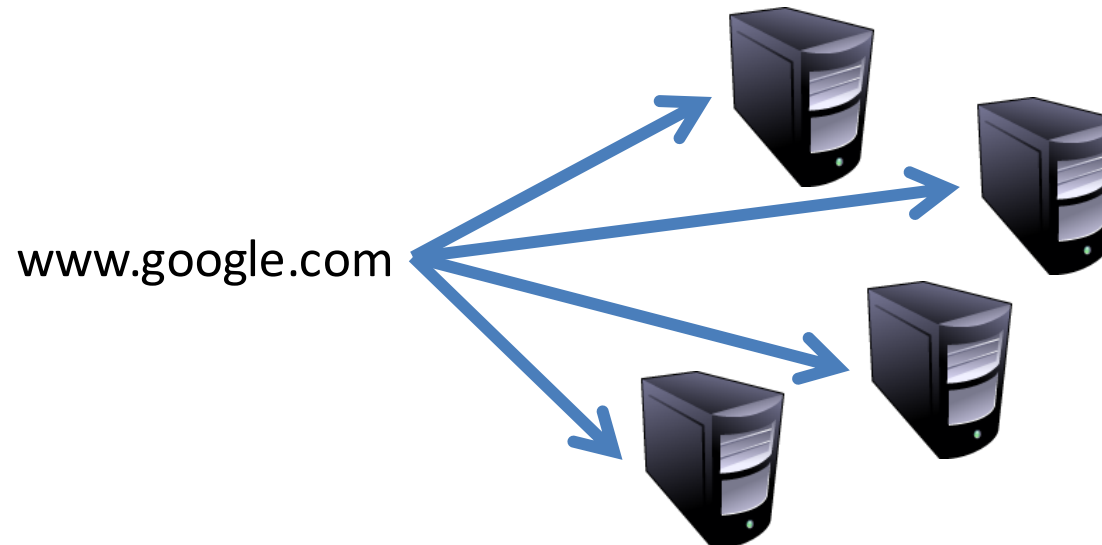


# Aliasing and Load Balancing

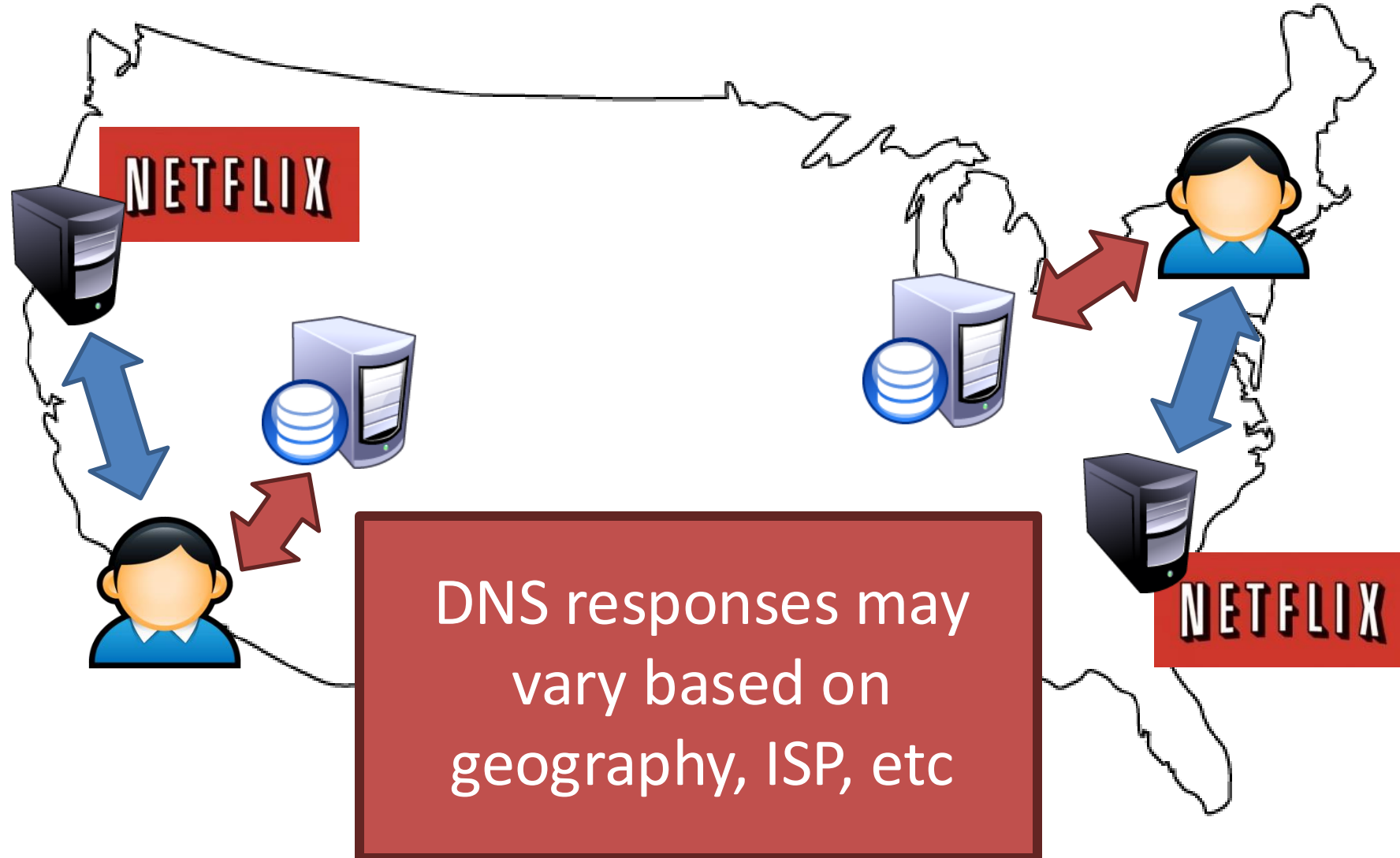
- One machine can have many aliases



- One domain can map to multiple machines



# Content Delivery Networks



# Inserting (or changing) records

Example: new startup “Network Utopia”

- Step 1: Register networkutopia.com at **DNS registrar**
  - provide names, IP addresses of authoritative name server (primary and secondary)

# Inserting (or changing) records

Example: new startup “Network Utopia”

- Step 2: Registrar inserts two RRs into .com TLD server
  - (networkutopia.com, dns1.networkutopia.com, NS)
  - (dns1.networkutopia.com, 212.212.212.1, A)

# Inserting (or changing) records

Example: new startup “Network Utopia”

- Step 3: Set up **authoritative server** at that name/address
  - Create records for the services:

# Inserting (or changing) records

Example: new startup “Network Utopia”

- Step 3: Set up **authoritative server** at that name/address
  - Create records for the services:
    - **type A record** for www.networkutopia.com
    - **type MX record** for @networkutopia.com email

# Inserting (or changing) records

- Example: new startup “Network Utopia”
- Register networkutopia.com at **DNS registrar**
  - provide names, IP addresses of authoritative name server (primary and secondary)
  - registrar inserts two RRs into .com TLD server
    - (networkutopia.com, dns1.networkutopia.com, NS)
    - (dns1.networkutopia.com, 212.212.212.1, A)
- Set up **authoritative server** at that name/address
  - Create records for the services:
    - **type A record** for www.networkutopia.com
    - **type MX record** for @networkutopia.com email

# Summary

- DNS maps human readable names to IP addresses
- DNS arranged into a hierarchy
  - Scalability / distributed responsibility
  - Autonomous control of local name servers
- Caching crucial for performance