Nswap2L: Transparently Managing Heterogeneous Cluster Storage Resources for Fast Swapping

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Data Intensive Computing

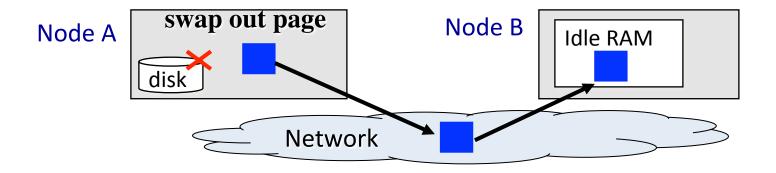
"The next 10 years will be shaped primarily by new algorithms that make sense of massive and diverse datasets and discover hidden value" Genevieve Bell, Intel Fellow, from SC'13 Keynote

- Require significant resources for good performance (or even feasibility)
- Stress the memory hierarchy
 - WS too large to fit in RAM → swapping
 - Out-of-core algorithms → temporary files

MEMSYS'16, newhall

General Purpose Clusters

- Variable WL lead to resource usage imbalances
 - Some nodes swapping, others have idle RAM
- Network RAM Storage is an option:
 - Use idle RAM of remote nodes for swap space



Cluster Storage Devices

Heterogeneous:

 HDD, Flash SSD, PCM(?), Network RAM, local and remote

Different Strengths:

- Network RAM: fast, volatile, variable capacity
- PCM: fast, expensive, capacity issues
- Flash SSD: faster R than W, erasure blocks, wear-out
- HDD: slow, cheap, still widely used
- → Cluster Storage likely to remain heterogeneous

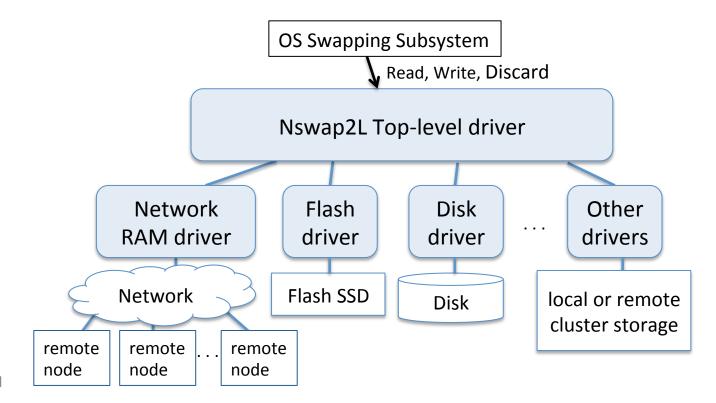
Storage and Performance

- Making best use of heterogeneous storage will have significant impact on application performance
- Problem: Tuning node OS's swapping subsystem for every possible combination of underlying storage is not possible
 - Still mostly optimized for disk
 - Some SSD support (TRIM/Discard)

Our Solution: Nswap2L

2 Level Device Driver Design:

- Transparent: to OS a single, fast, random access device
- Manages underlying heterogeneous storage devices
 Adaptable/Tunable Policies for Data Placement & Prefetching



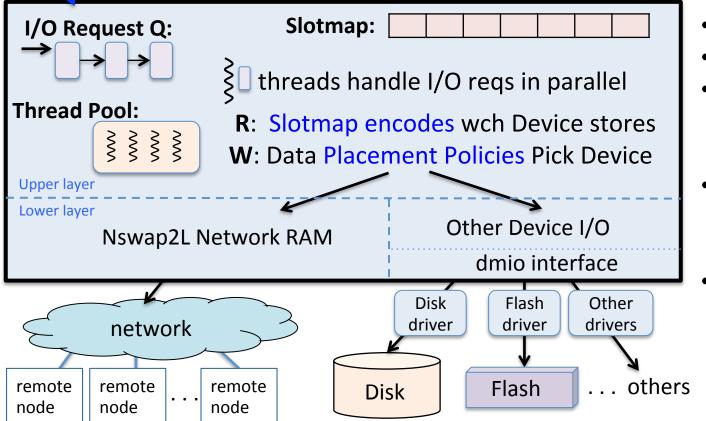
System Architecture

- Linux 4.0 driver, added as swap partition on nodes Underlying devices at load & added later (via /sys)
- OS sends Read/Write/Discard to swap in/out/free

Nswap2L: Top-level Device Driver

R/W/D Req.

From OS

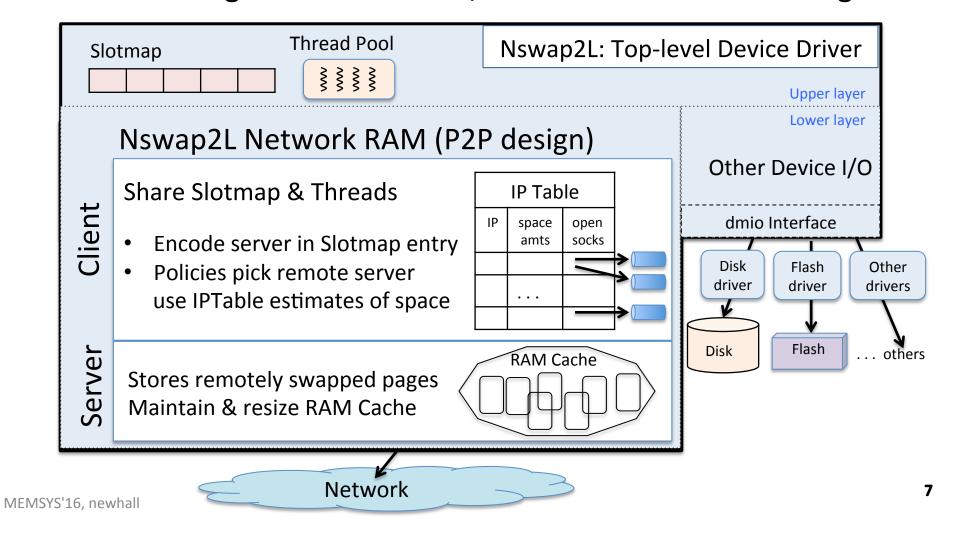


- Multi-threaded
- Slotmap
- Placement
 Polices on W
- dmio to most bottom-level devices
- Nework RAM part of Nswap2L

Nswap2L Network RAM

Scalable: P2P design, local estimate available Network RAM

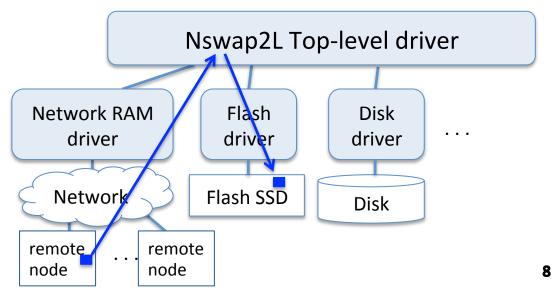
Adaptable: Amount of RAM available for Network RAM grows & shrinks w/cluster workload RAM usage



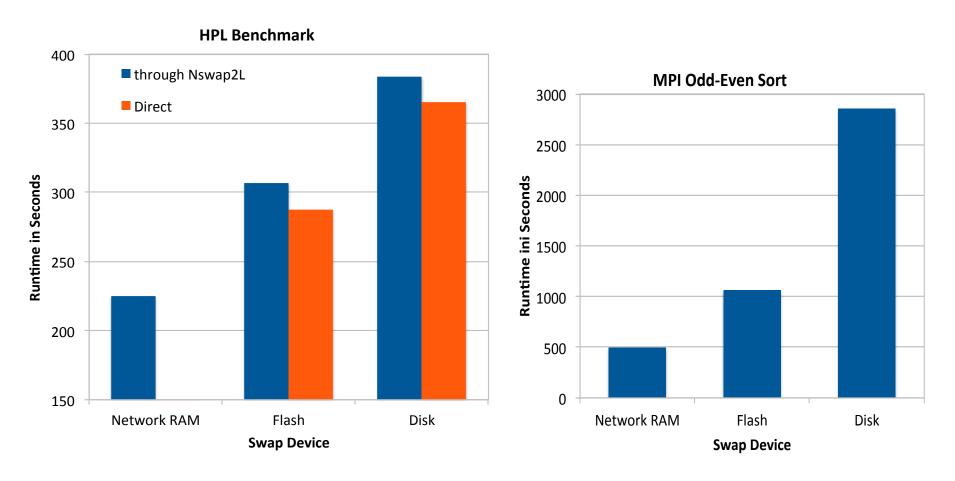
Prefetching

- Moving page data between underlying devices
 - Transparent to OS: I/O internal to Nswap2L
- Take advantage of strengths of different devices (ex) Prefetching from Network RAM to Flash:
 - free up NW RAM for future fast Writes
 - increase Read parallelism over Flash and NW RAM

Tunable Policies via /sys interface How much? How often? Which pages? To/From?



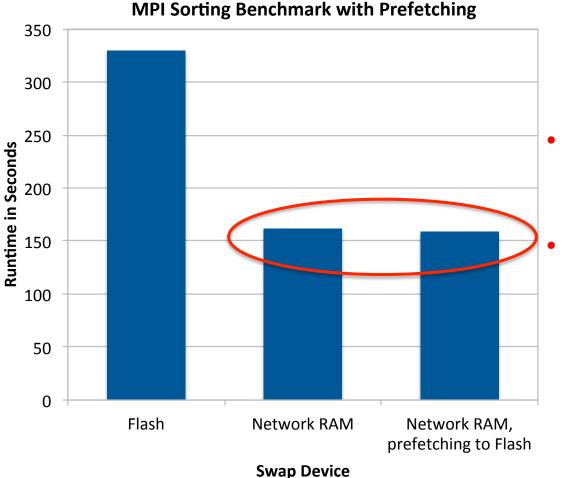
Nswap2L Performance



Nswap2L loaded as Linux 4.0.4 device driver, and added as swap partition on 16 node cluster, 10 Gbit Ethernet. Underlying Network RAM, Flash SSD & HDD.

Benefit of Prefetching

 Swapping out to faster Network RAM with prefetching some pages from Network RAM to slower Flash

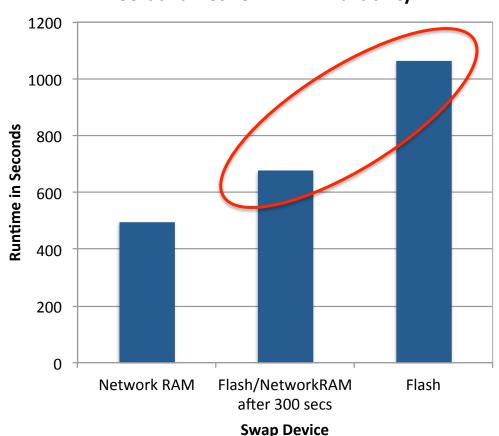


- Increase degree of Read parallelism over Flash and Network RAM
- Get slightly faster runtime than swapping to faster Network RAM alone

Benefit of Adaptable Policies

 Flash available for entire run, Network RAM becomes available after about 5 minutes of runtime

MPI Sort and Network RAM Availability



- Adaptable Data
 Placement Polices allow
 Nswap2L to discover
 faster underlying storage
 and make use of it as it
 becomes available
- Significantly faster runtime than statically swapping to Flash alone

Nswap2L:

- Fast Backing Store for Swap in Clusters
- Transparently Manages Heterogeneous Cluster Storage Devices, including its own Network RAM
 - Presents as single, fast swap device to node OSs
- Adaptable polices result in performance improvements over fastest single swap device

Current and Future Directions

- Expand Nswap2L to be used for other types of backing storage
 - File data, particularly targeting temporary files
- Further investigating adaptive policies
 - Data placement
 - Prefetching
 - Network RAM growing/shrinking
 - Implementing a more Extensible Policy interface
 - adding in new policy on the fly
- Evaluate on Larger Systems, Scalability

Thank You

Questions?

www.cs.swarthmore.edu/~newhall/nswap2L.html