

# CS 43: Computer Networks

## Internet Routing

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# Hierarchical routing

Our routing study thus far - idealization

- all routers identical
- network “flat”

... *not* true in practice

*Scale:* with 600 million destinations:

- can't store all dest's in routing tables!
- routing table exchange would swamp links!

*Administrative autonomy*

- internet = network of networks
- each network admin may want to control routing in its own network

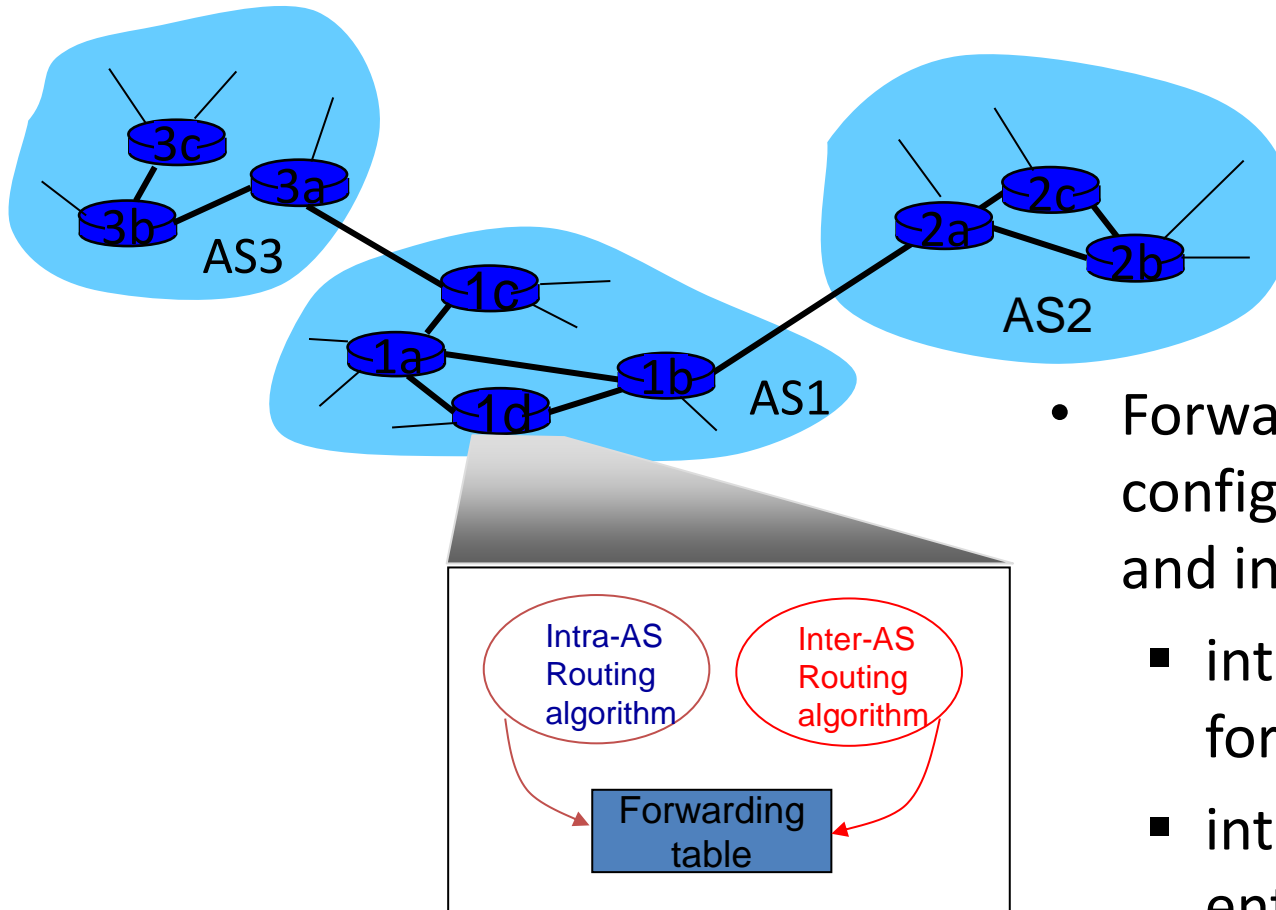
# Hierarchical routing

- We aggregate routers into regions, “autonomous systems” (AS)
- Routers in same AS run same routing protocol
  - “intra-AS” or “interior” routing protocol
  - routers in different AS can run different intra-AS routing protocol

## *Gateway (or border) router:*

- at “edge” of its own AS
- has link to router in another AS

# Interconnected ASes



- Forwarding table configured by both intra- and inter-AS routing algs
  - intra-AS sets entries for internal dests
  - inter-AS & intra-AS sets entries for external dests

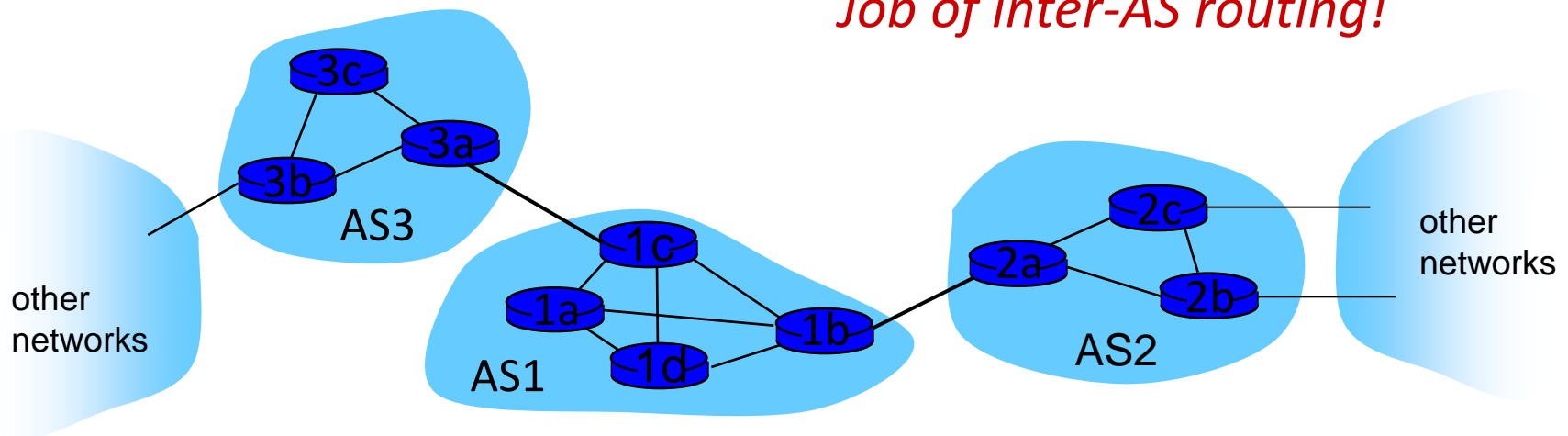
# Inter-AS tasks

- Suppose router in AS1 receives a datagram destined outside of AS1:
  - Router should forward packet to gateway router, but which one?

*AS1 must:*

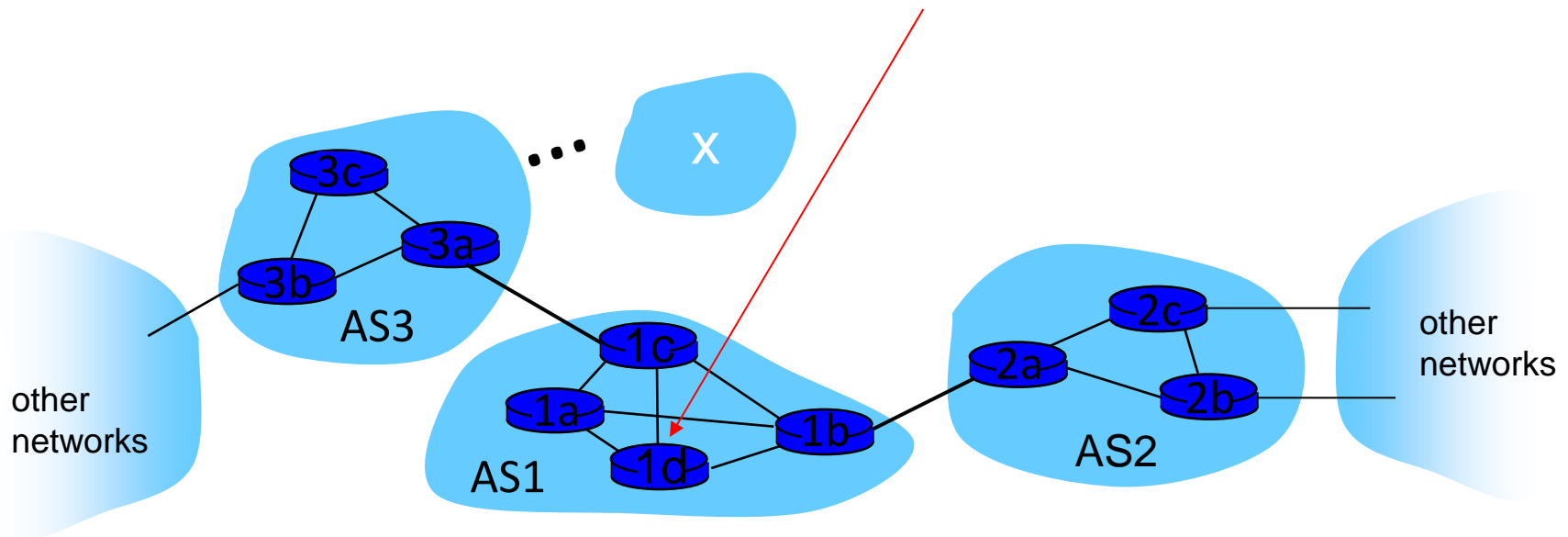
1. Learn which destds are reachable through AS2, which through AS3
2. Propagate this reachability info to all routers in AS1

*Job of inter-AS routing!*



# Example: setting forwarding table in router 1d

- Suppose AS1 learns (via inter-AS protocol) that AS with prefix  $x$  is reachable via AS3 (gateway 1c), but not via AS2
  - inter-AS protocol propagates reachability to all internal routers
- Router 1d determines from intra-AS routing info that its interface  $I$  is on the least cost path to 1c
  - Installs forwarding table entry  $(x, I)$

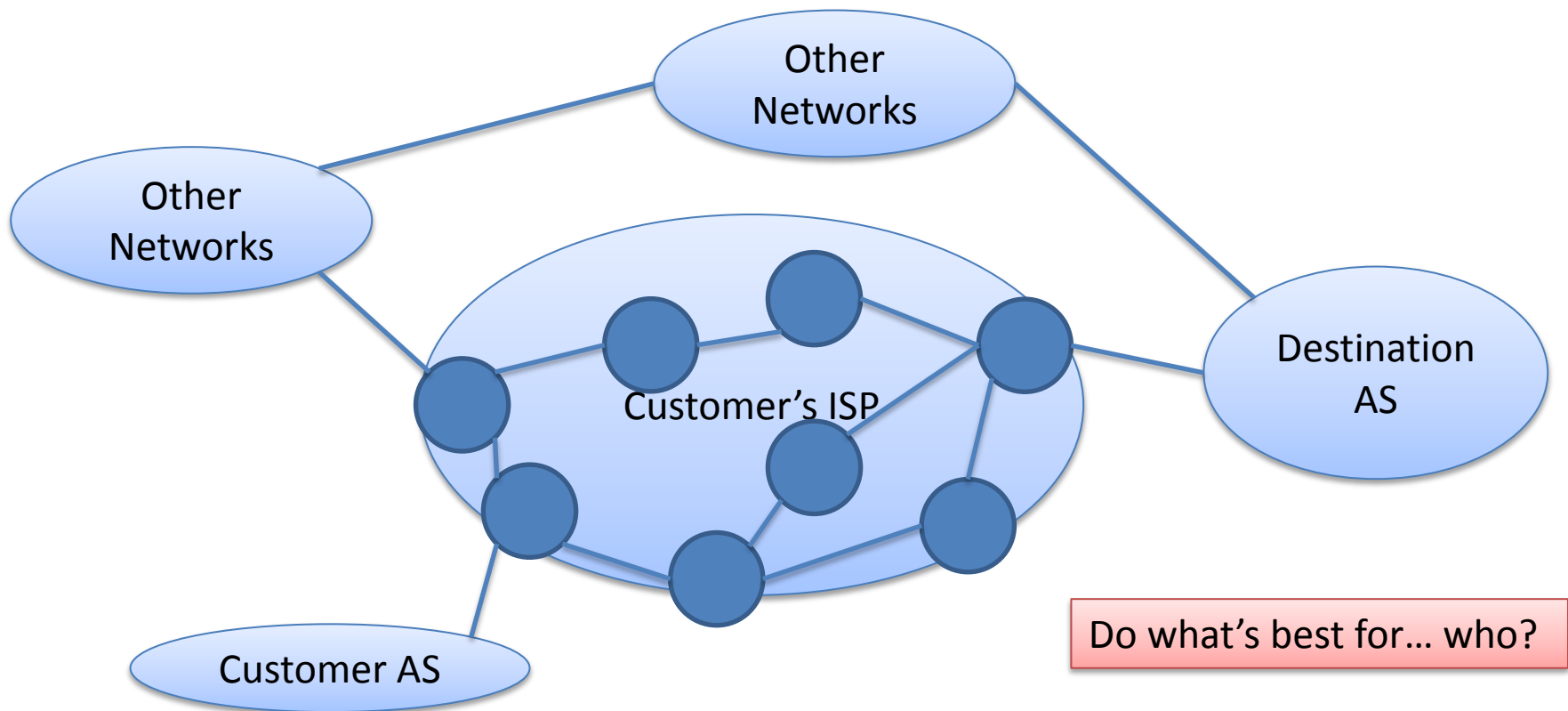


If an external destination is reachable from multiple gateways, a router inside the AS should forward packets for that destination to

- A. The closest gateway that can reach the destination.
- B. The gateway that has the least-cost external path to the destination.
- C. The gateway that has the least-cost path for both the internal and external path.
- D. Somewhere else.

# Routing Policy

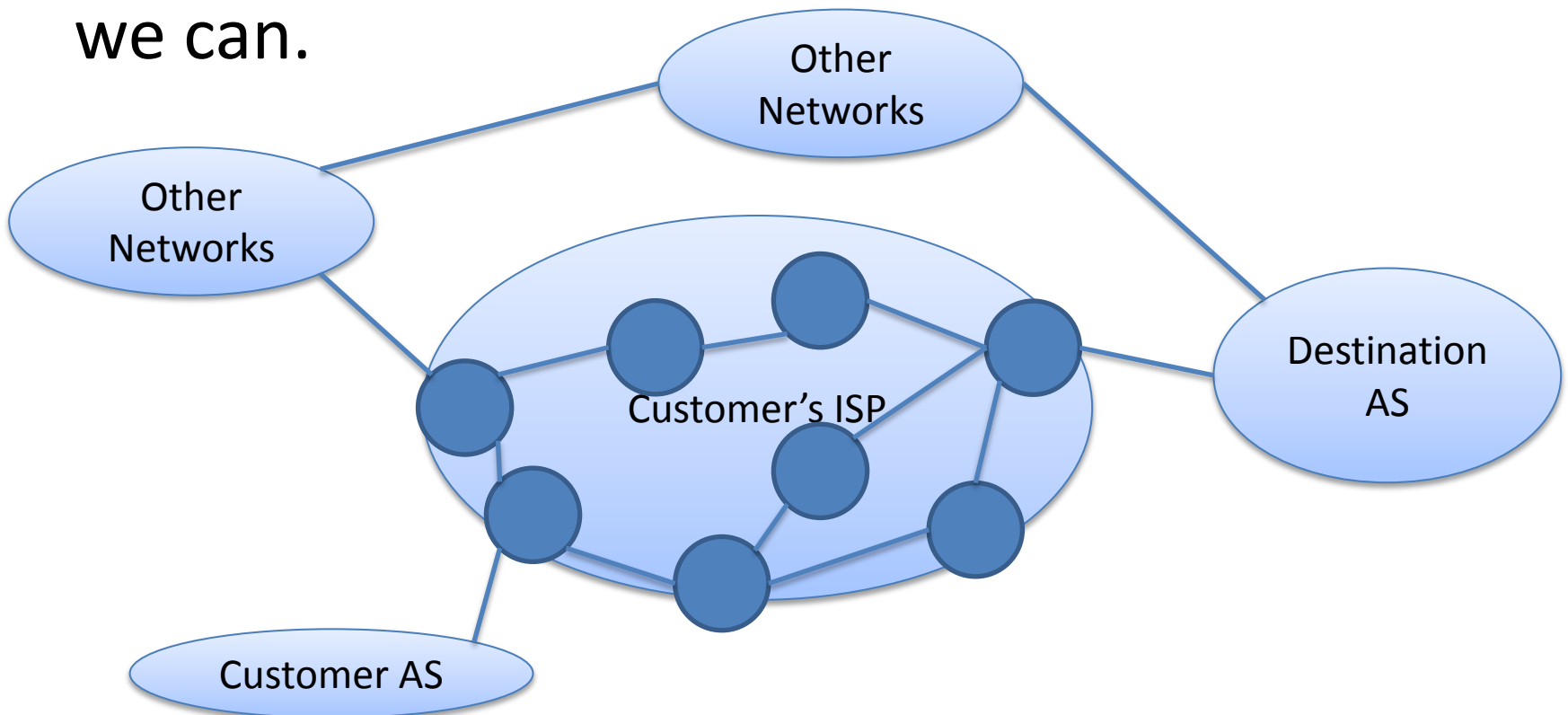
- How should the ISP route the customer's traffic to the destination?





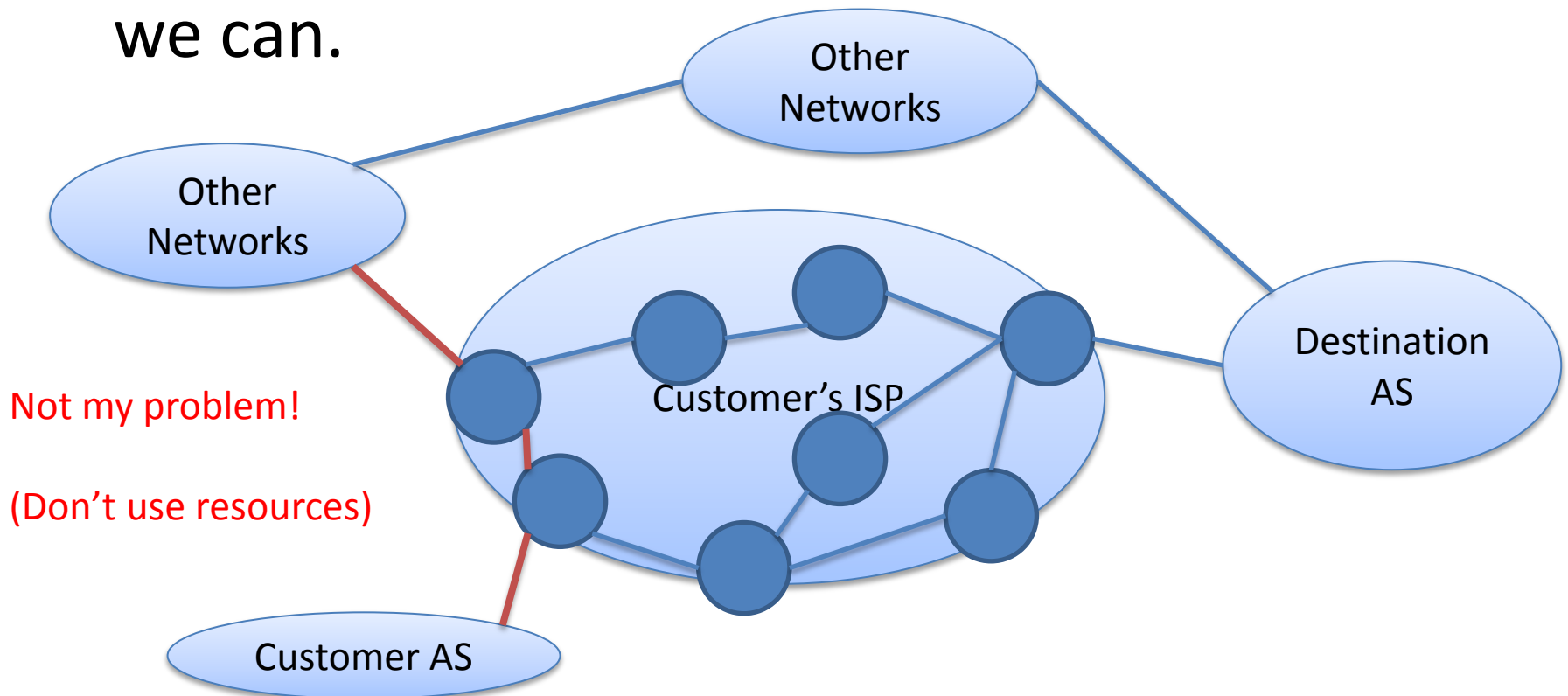
# Hot Potato Routing

- Hot Potato: get rid of packets ASAP!
- Best path: get it as close to the destination as we can.



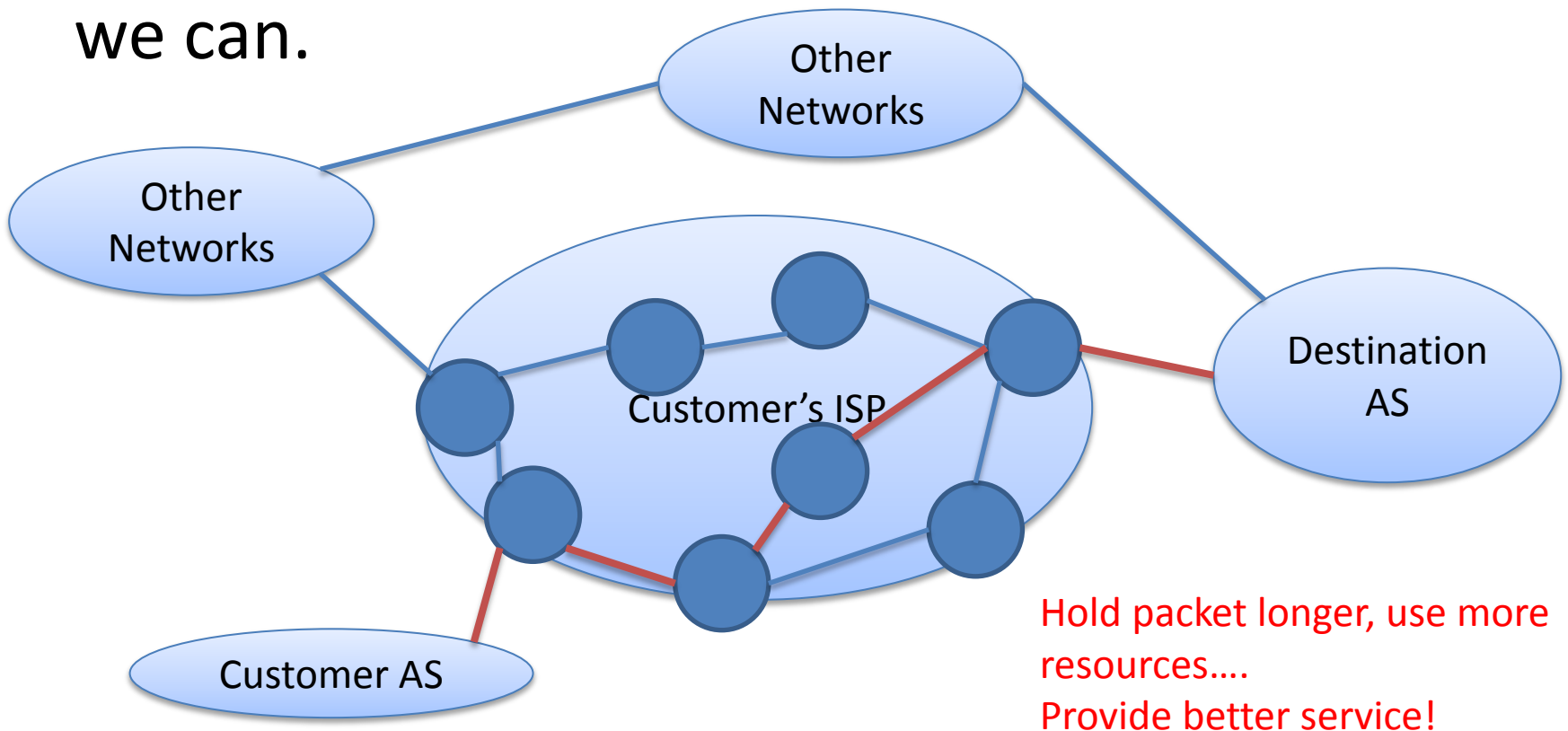
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# Route Selection

- Often dictated by non-technical factors
- When governed by protocols, two categories:
  - Intra-AS / Interior gateway protocols
  - Inter-AS / Exterior gateway protocols

# Why different Intra-, Inter-AS routing ?

## *Policy:*

- inter-AS: admin wants control over how its traffic routed, who routes through its net.
- intra-AS: single admin, so no policy decisions needed

## *Scale:*

- hierarchical routing saves table size, reduced update traffic

## *Performance:*

- intra-AS: can focus on performance
- inter-AS: policy may dominate over performance

# Intra-AS Routing

- Also known as *interior gateway protocols (IGP)*
- *Distance Vector:*
  - RIP: Routing Information Protocol
  - (E)IGRP: Interior Gateway Routing Protocol  
(Cisco proprietary)
- *Link State:*
  - OSPF: Open Shortest Path First
  - IS-IS: Intermediate system to Intermediate system

OSPF and IS-IS are deployed most commonly today!

# Intra-AS Routing

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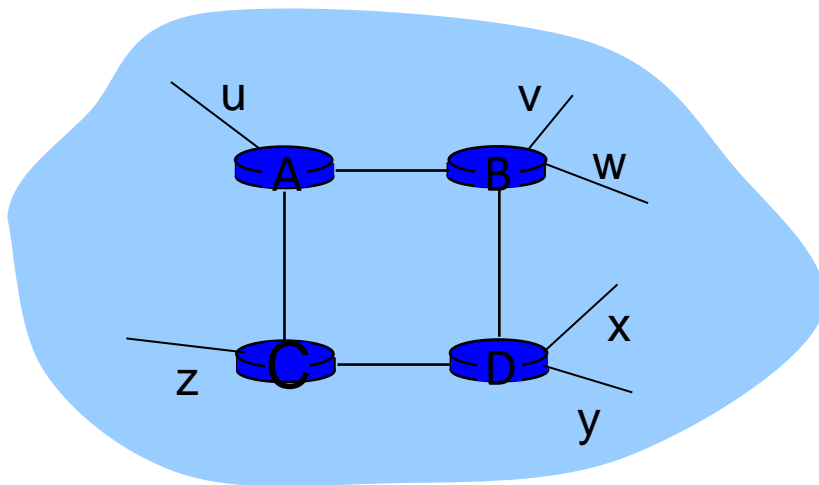
Goal:

Get traffic that is already in an AS to a destination inside that same AS.

OSPF and IS-IS are deployed most commonly today!

# RIP (Routing Information Protocol)

- Included in BSD-UNIX distribution in 1982
  - distance metric: # hops (max = 15 hops), each link has cost 1
  - hops = number of subnets traversed
  - Distance vectors exchanged with neighbors every 30 sec
  - Each advertisement: list of up to 25 destination *subnets*



from router A to destination *subnets*:

<u>subnet</u>	<u>hops</u>
u	1
v	2
w	2
x	3
y	3
z	2



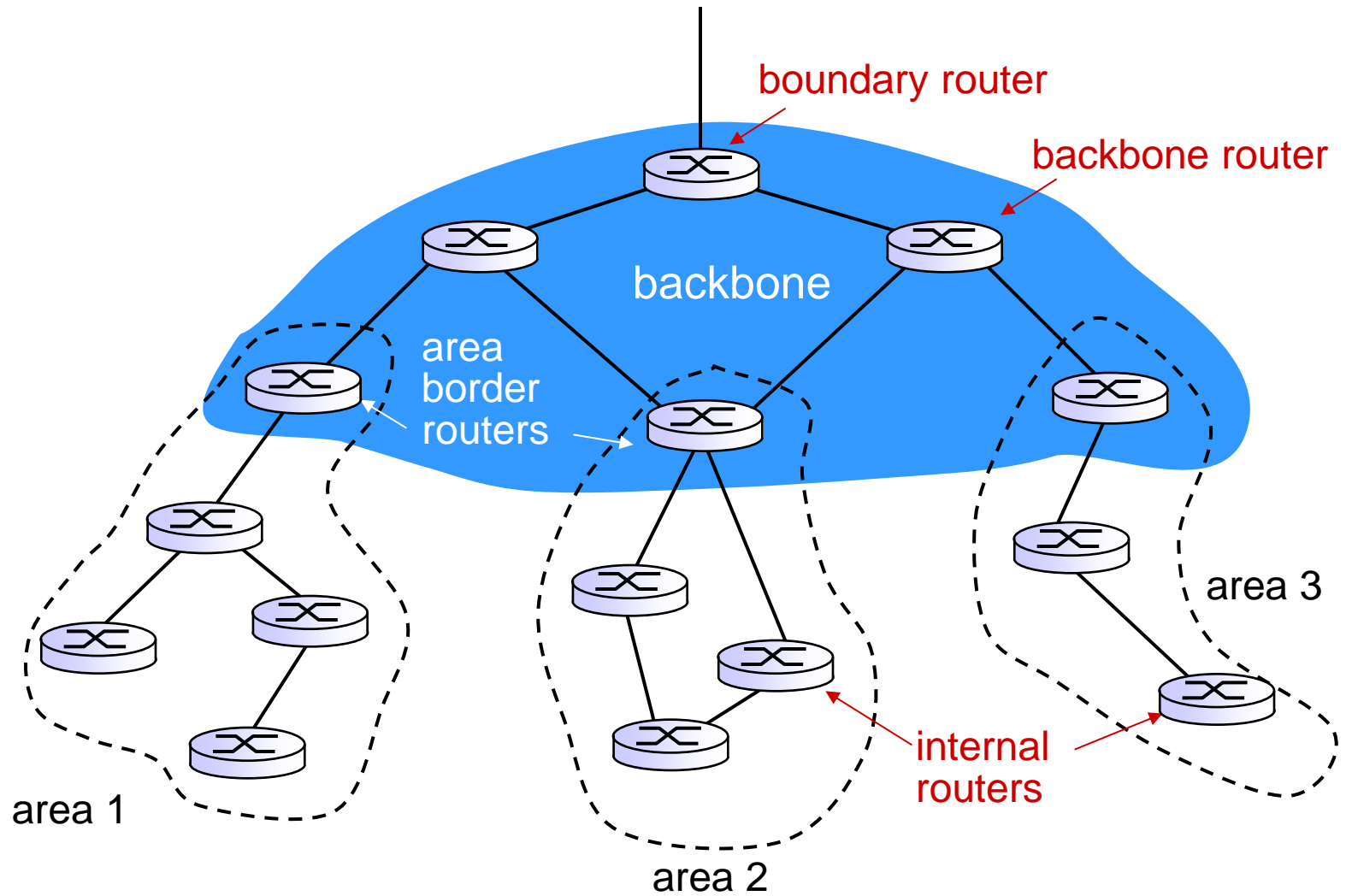
# OSPF (Open Shortest Path First)

- Link state protocol (reliable flooding of LSAs)
- “Open”: standardized, publicly available implementations
- Multiple equal-cost paths allowed (load balancing)
- Additional features:
  - OSPF messages authenticated (to prevent malicious intrusion)
  - Hierarchical OSPF for large autonomous systems.

# Hierarchical OSPF

- *Two-level hierarchy*: local area, backbone.
  - link-state advertisements only in area
  - each nodes has detailed area topology; only know direction (shortest path) to nets in other areas.
- *Area border routers*: “summarize” distances to nets in own area, advertise to other Area Border routers.
- *Backbone routers*: route between local areas
- *Boundary routers*: connect to other AS's.

# Hierarchical OSPF



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# Internet inter-AS routing: BGP

Goal:

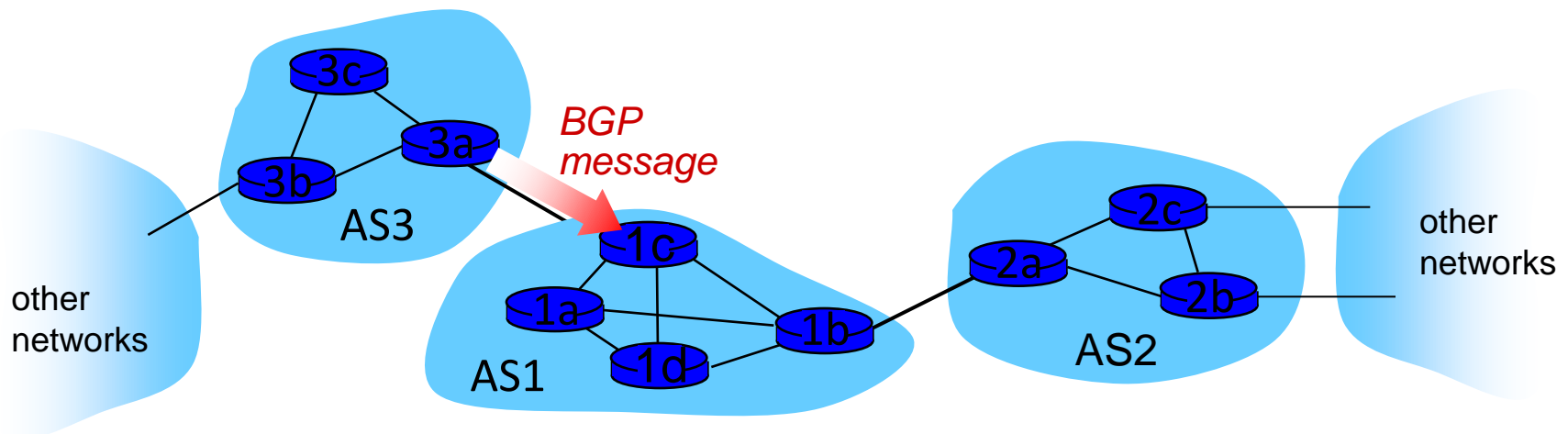
Get traffic from one AS to another.

# Internet inter-AS routing: BGP

- **BGP (Border Gateway Protocol):**  
*The de facto inter-domain routing protocol*
- BGP provides each AS a means to:
  - **external BGP:** obtain subnet reachability information from neighboring ASs.
  - **internal BGP:** propagate reachability information to all AS-internal routers.
  - determine “good” routes to other networks based on reachability information and policy.
- Allows a subnet to advertise its prefix to the rest of the Internet

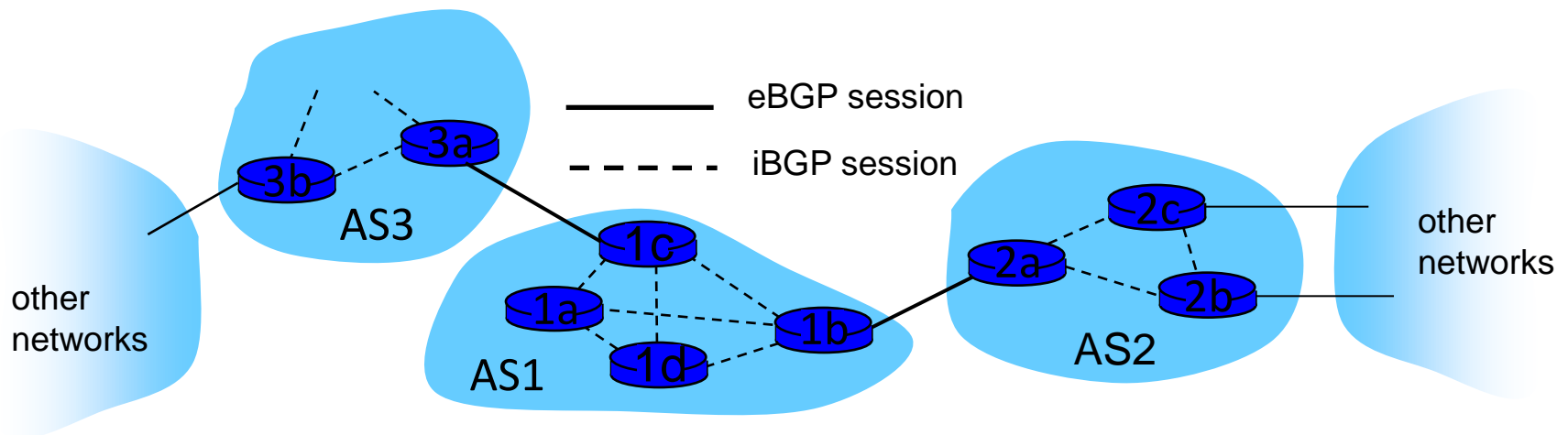
# BGP

- **BGP session:** two BGP routers (“peers”) exchange BGP messages:
  - Advertising *paths* to different destination network prefixes (“path vector”)
  - Exchanged over long-term TCP connections
- When AS3 advertises a prefix to AS1:
  - AS3 *promises* it will forward datagrams towards that prefix
  - AS3 can aggregate prefixes in its advertisement



# BGP: Distributing Path Information

- Using eBGP session between 3a and 1c, AS3 sends prefix reachability info to AS1.
  - 1c can then use iBGP to distribute new prefix info to all routers in AS1
  - 1b can then re-advertise new reachability info to AS2 over 1b-to-2a eBGP session
- When a router learns of a new prefix, it creates an entry for the prefix in its forwarding table.





# Path attributes and BGP routes

- An advertised prefix includes BGP attributes
  - prefix + attributes = “route”
- Two important attributes:
  - **AS-PATH**: contains list of ASs through which prefix advertisement has passed:
    - If AS2 advertises a prefix to AS1, AS1 will advertise path: AS1 AS2...
    - Ignore routes that include yourself in them!
  - **NEXT-HOP**: indicates specific internal-AS router to next-hop AS. (may be multiple links from current AS to next-hop-AS)
- Gateway router receiving route advertisement uses **import policy** to accept/decline
  - e.g., never route through AS x
  - *policy-based* routing

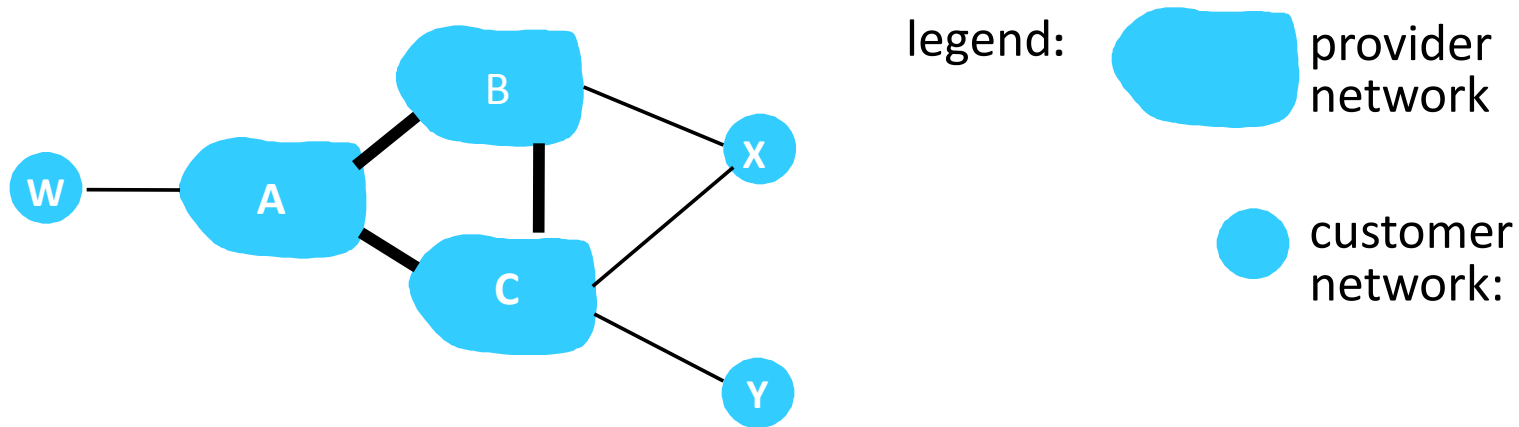
# BGP Route Selection

- Router may learn about more than one route to destination AS, selects route based on:
  - local preference value attribute: administrative policy
  - shortest AS-PATH
  - closest NEXT-HOP router: hot potato routing
  - additional criteria

Which routes a BGP router advertises will depend on...

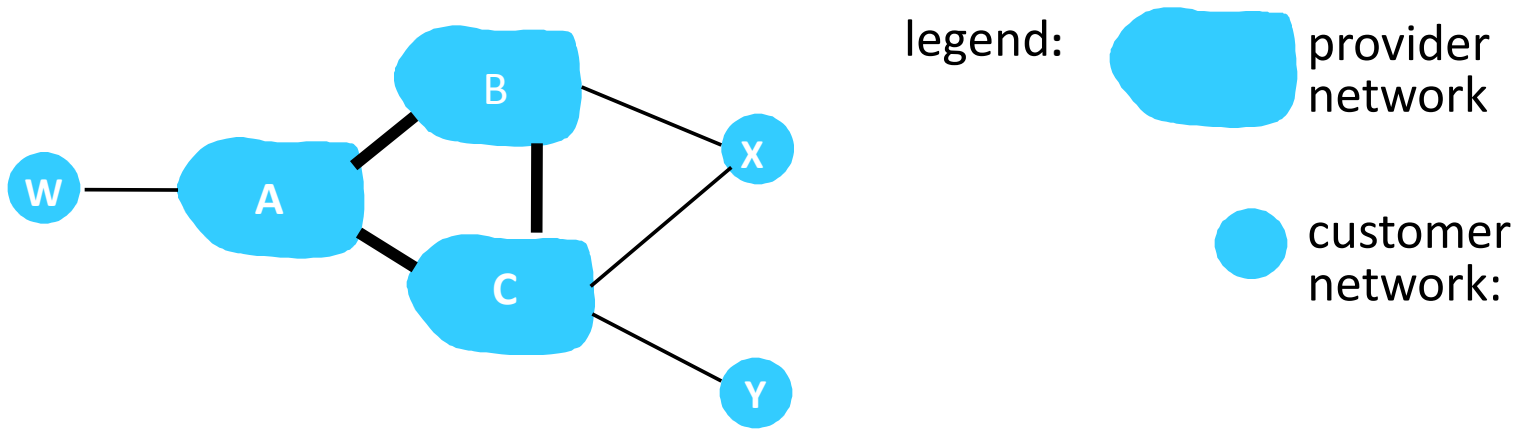
- A. which ISPs have contractual agreements.
- B. the shortest path to a subnet/prefix.
- C. which subnets are customers of an ISP.
- D. More than one of the above. (which?)

# BGP routing policy



- A,B,C are *provider networks*
- X,W,Y are customer (of provider networks)
- X is *dual-homed*: attached to two networks
  - X does not want to route from B via X to C
  - .. so X will not advertise to B a route to C

## BGP routing policy (2)



- A advertises path AW to B
- B advertises path BAW to X
- Should B advertise path BAW to C?
  - B gets no “revenue” for routing CBAW since neither W nor C are B’s customers
  - B wants to force C to route to w via A
  - B wants to route *only* to/from its customers!

# Summary

- As we've seen before (DNS), a hierarchy can help manage state storage constraints.
  - intra-AS routing: lots of info about local routes
  - inter-AS routing: less info about far away routes
- BGP: the inter-AS routing protocol for the Internet
  - Decisions often contractual
- BGP advertises AS prefixes, including:
  - entire path of ASes along the way
  - which border router heard the advertisement (Next Hop)