

FOR CLARITY,
Consider a language L .

① Is L decidable?

② How would you design a TM to decide L ?
Design it. Give the algorithm.

These questions are different!

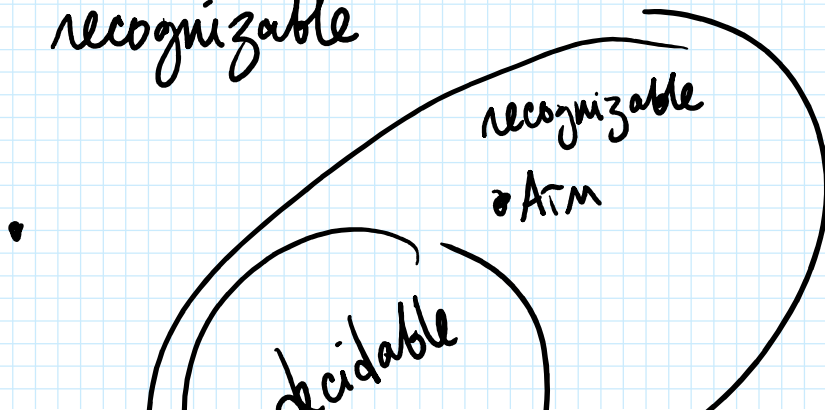
If you can give a TM in Q2, then Q1 is YES.

Sometimes we can prove Q1 is YES
even though we don't have enough info to answer Q2.

$$\overline{A_{TM}} = \{ \langle M, w \rangle \mid \begin{array}{l} M \text{ is a TM} \\ \text{which accepts string } w \end{array} \}$$

A_{TM} is not decidable

A_{TM} is recognizable



(decidable)

If M is a TM, define $L(M) = \{w \mid M \text{ accepts } w\}$.
This is the lang recognized by M .

<u>list of all TMs</u>	<u>all Turing-recognizable langs.</u>	<u>all co-recognizable langs</u>
$\langle M_1 \rangle$	$L(M_1)$	$\overline{L(M_1)}$
$\langle M_2 \rangle$	$L(M_2)$	
$\langle M_3 \rangle$	$L(M_3)$	$\overline{L(M_3)}$
\vdots		

countable ✓

countable
(might contain
duplicates)

also countable ✓

Claim: $\overline{A_{TM}}$ is not decidable.

Proof: (by contradiction)

Suppose M decides $\overline{A_{TM}}$.

$N =$ "On input $\langle R, w \rangle$:

1. Run M on $\langle R, w \rangle$.
2. Do the opposite."

N is a decider for A_{TM} .

$\Rightarrow \Leftarrow$ We know A_{TM} is undecidable!