Algorithm 3 analysis

ALGORITHM3()

- $1 \ k = 1.$
- 2 while you haven't found your friend:
- 3 walk k miles east
- 4 return to start
- 5 walk k miles west
- 6 return to start
- 7 k = 2 * k.

Now let's analyze this algorithm.

<u>Correctness</u>: Clearly you do eventually reach your friend, and then you stop, so the algorithm is correct.

In one complete iteration of the algorithm, you walk 4k miles. How many complete iterations do you do?

Let's say that we stop on the t^{th} iteration, so $2^{t-1} < m \leq 2^t$. (The first inequality is because we finished iteration t-1 and still had not finished; the second inequality is because we find her on this iteration.)

- <u>Case 1:</u> Your friend is m miles east of you. The total distance walked

$$= 4 \cdot 1 + 4 \cdot 2 + 4 \cdot 4 + 4 \cdot 8 + \dots + 4 \cdot 2^{t-1} + m \quad \text{(the last, partial iteration adds the } +m \text{ term})$$

$$= \sum_{k=0}^{t-1} (4 \cdot 2^k) + m$$

$$= 4 \cdot \sum_{k=0}^{t-1} (2^k) + m$$

$$= 4 \cdot (2^t - 1) + m \quad \text{(by the fact that } \sum_{k=0}^{t-1} 2^k = 2^t - 1)$$

$$< 4 \cdot (2m - 1) + m \quad \text{(because } 2^{t-1} < m \text{ so } 2^t < m)$$

$$= 8m - 4 + m$$

$$= O(m)$$

- <u>Case 2</u>: Your friend is *m* miles west of you.

Total distance walked

$$= 4 \cdot 1 + 4 \cdot 2 + \dots + 4 \cdot 2^{t-1} + (2^{t} + 2^{t} + m)$$
 (don't forget the partial last iteration!)

$$= \sum_{k=0}^{t-1} (4 \cdot 2^{k}) + 2^{t+1} + m$$

$$= 4 \cdot \sum_{k=0}^{t-1} (2^{k}) + 2^{t+1} + m$$

$$= 4 (2^{t} - 1) + 2^{t+1} + m$$
 (by the fact that $\sum_{k=0}^{t-1} 2^{k} = 2^{t} - 1$)

$$< 4 (2m - 1) + 4m + m$$
 (because $2^{t-1} < m$ so $2^{t} < 2m$ and $2^{t+1} < 4m$)

$$= 8m - 4 + 4m + m$$

$$= O(m)$$

Note: For this class, you are not expected to already know facts like $\sum_{k=0}^{t-1} 2^k = 2^t - 1$. If you run across calculations like this in homework, you should feel free to use outside resources (math books, Wolfram Alpha, etc.) to help you solve the problem.