CS 43: Computer Networks

09: Bit-Torrent and P2P October 1, 2019



Reading Quiz

Last class

• Email!

– Application Layer Protocol

If SMTP only allows 7-bit ASCII, how do we send pictures/videos/files via email?

- A. We encode these objects as 7-bit ASCII
- B. We use a different protocol instead of SMTP
- C. We're really sending links to the objects, rather than the objects themselves

Base 64

- Designed to be an efficient way to send binary data as a string
- Uses A-Z, a-z, 0-9, "+" and "/" as digits
- A number with digits $d_n d_{n-1} \dots d_1 d_0 = 64^n d_n + 64^{n-1} d_{n-1} + \dots + 64^* d_1 + d_0$
- Recall from CS 31: Other non-base-10 number systems (binary, octal, hex).

Multipurpose Internet Mail Extensions (MIME)

- Special formatting instructions
- Indicated in the header portion of message (not SMTP)
 - SMTP does not care, just looks like message data
- Supports
 - Text in character sets other than ASCII
 - Non-text attachments
 - Message bodies with multiple parts
 - Header information in non-ASCII character sets

MIME

- Adds optional headers
 - Designed to be compatible with non-MIME email clients
 - Both clients must understand it to make sense of it
- Specifies content type, other necessary information
- Designates a boundary between email text and attachments

Mail access protocols



- SMTP: delivery/storage to receiver's server
- mail access protocol: retrieval from server
 - POP: Post Office Protocol: authorization, download
 - IMAP: Internet Mail Access Protocol: more features, including manipulation of stored messages on server
 - HTTP: gmail, Hotmail, Yahoo! Mail, etc.

POP3 protocol



More about POP3

- Previous example uses "download and delete" mode
 Bob cannot re-read e-mail if he changes client
- POP3 "download-and-keep": copies of messages on different clients
- POP3 is stateless across sessions
- Limitations:
 - Can't retrieve just the headers
 - Can't impose structure on messages

IMAP

- Keeps all messages in one place: at server
- Allows user to organize messages in folders
- Keeps user state across sessions:
 - names of folders and mappings between message IDs and folder name
- Can request pieces of a message (e.g., text parts without large attachments)

Webmail

- Uses a web browser
- Sends emails using HTTP rather than POP3 or IMAP
- Mail is stored on the 3rd party webmail company's servers

Summary

- Three main parts to email:
 - Mail User Agent (mail client): read / write for humans
 - Mail Transfer Agent: server that accepts / sends messages
 - SMTP protocol used to negotiate transfers
- No SMTP support for fraud detection
- Extensions (MIME) and encodings (Base64) for sending non-text data

Today

- P2P applications
- BitTorrent
 - Cooperative file transfers
- Briefly: Distributed Hash Tables

- Finding things without central authority

Where we are

Application: the application (So far: HTTP, Email, DNS) Today: BitTorrent, Skype, P2P systems

Transport: end-to-end connections, reliability

Network: routing

Link (data-link): framing, error detection

Physical: 1's and 0's/bits across a medium (copper, the air, fiber)

Designating roles to an endpoint

Client-server architecture

Peer-to-peer architecture





File Transfer Problem

• You want to distribute a file to a large number of people as quickly as possible.

Traditional Client/Server

- Many clients, 1 (or more) server(s)
- Web servers, DNS, file downloads, video streaming

Traditional Client/Server



What is the biggest problem you run into with the traditional C/S model?



- A. Scalability (how many end-hosts can you support?)
- B. Reliability (what happens on failure?)
- C. Efficiency (fast response time)

Traditional Client/Server



P2P Solution



Client-server vs. P2P: example



Let F = file size, client UL rate = u, server rate = u_s , d = client DL rate Assumptions: F/u = 1 hour, $u_s = 10u$, $d_{min} \ge u_s$

P2P Solution



Do we need a centralized server at all? Would you use one for something?



D. Something else.

Napster Architecture



File Search via Flooding in Gnutella



What are the cons of the Gnutella approach?

- A. Access to rare files
- B. Traffic overhead
- C. Redundancy
- D. Something else.

Peer Lifetimes: Highly available?

- Study of host uptime and application uptime (MMCN 2002)
 - Sessions are short (median is 60 minutes)
 - Hosts are frequently offline



Slide 33

Resilience to Failures and Attacks

- Previous studies (Barabasi) show interesting dichotomy of resilience for "scale-free networks"
 - Resilient to random failures, but not attacks
- Here's what it looks like for Gnutella



1771 Peers in Feb, 2001

After top 4% of peers are removed After random 30% of peers removed

P2P file distribution: BitTorrent

- File divided into chunks (commonly 256 KB) \bullet
- Peers in torrent send/receive file chunks



.torrent files

- Contains address of tracker for the file
 - Where can I find other peers?
- Contain a list of file chunks and their cryptographic hashes
 - This ensures pieces are not modified

BitTorrent : Peer Joining

- has no chunks, but will accumulate them over time from other peers
- registers with tracker to get list of peers, connects to subset of peers ("neighbors")



P2P file distribution: BitTorrent

- While downloading, peer uploads chunks to other peers
- Churn: peers may come and go
 - Peer may change peers with whom it exchanges chunks



Requesting Chunks

- At any given time, peers have different subsets of file chunks.
- Periodically, each asks peers for list of chunks that they have.
- Once peer has entire file, it may (selfishly) leave or (altruistically) remain in torrent

Sharing Pieces



If you're trying to receive a file, which chunk should you request next?



Requesting Chunks

Bootstrap: random selection

- Initially, you have no pieces to trade
- Essentially, beg for free pieces at random
- Steady-state: rarest piece first
 - Ensures that common pieces are saved for last
- Endgame
 - Simultaneously request final pieces from multiple peers
 - Cancel connections to slow peers
 - Ensures that final pieces arrive quickly

100%

Sending Chunks

- A node sends chunks to those four peers currently sending it chunks at highest rate
 - other peers are choked (do not receive chunks)
 - re-evaluate top 4 every ~10 secs
- Every 30 seconds: randomly select another peer, start sending chunks
 - "optimistically unchoke" this peer
 - newly chosen peer may join top 4

Academic Interest in BitTorrent

- BitTorrent was enormously successful
 - Large user base
 - Lots of aggregate traffic
 - Invented relatively recently
- Research
 - Modifications to improve performance
 - Modeling peer communications (auctions)
 - Gaming the system (BitTyrant)

Incentives to Upload

- Every round, a BitTorrent client calculates the number of pieces received from each peer
 - The peers who gave the most will receive pieces in the next round
 - These decisions are made by the unchoker
- Assumption
 - Peers will give as many pieces as possible each round
 - Based on bandwidth constraints, etc.
- Can an attacker abuse this assumption?

Unchoker Example



Abusing the Unchocker

• What if you really want to download from someone?



Sybil Attack



BitTyrant

- Piatek et al. 2007
 - Implements the "come in last strategy"
 - Essentially, an unfair unchoker
 - Faster than stock BitTorrent (For the Tyrant user)

Hierarchical P2P Networks

• FastTrack network (Kazaa, Grokster, Morpheus, Gnutella++)



Skype: P2P VoIP



- P2P client supporting VoIP, video, and text based conversation, buddy lists, etc.
 - Overlay P2P network consisting of ordinary and Super Nodes (SN)
- Each user registers with a central server
 - User information propagated in a decentralized fashion

Do the benefits of hierarchical P2P networks out-weight the cons?

- A. Pros: Scalability
- B. Pros: Limits flooding
- C. Cons: No guarantees of performance
- D. Cons: Failure?