

Merge Sort

Split array in half

merge sort each half

merge the halves \rightarrow linear $O(n)$

4 5 1 9 7 2 3 7

4 5 1 9 7 2 3 7

4 5 1 9 7 2 3 7

4 5 1 9 7 2 3 7

[4 5 1 9 2 7 3 7

1 4 5 9 2 3 7 7]

1 2 3 4 5 7 7 9

4 calls
2 elements

↓
1 call
8 elements

2 calls
4 elements

example 3 recursions deep

$$\lceil \log_2 n \rceil$$

$$(c \cdot n)(\log_2 n)$$

$$O(n \log n)$$

vs

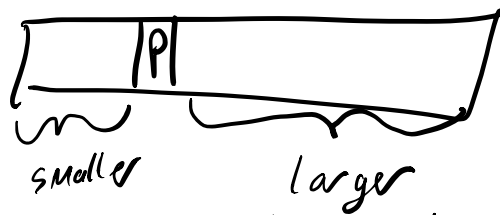
$$O(n^2)$$

Quick Sort

pick some pivot element

two sub-arrays

- everything less than pivot
- everything larger than pivot



quick sort sub-arrays

Function QuickSort (A, start, end)

if start == end: // base case
return

p = partition (A, start, end)

QuickSort (A, start, p-1)

QuickSort (A, p+1, end)

end Function

Function Partition(A, start, end):

$s = \cancel{0}$ start

$l = \cancel{0}$ start

pivot = A[end]

while $l < end$:

if $A[l] < pivot$:

swap $A[l], A[s]$

$s++$;

end if

$l++$;

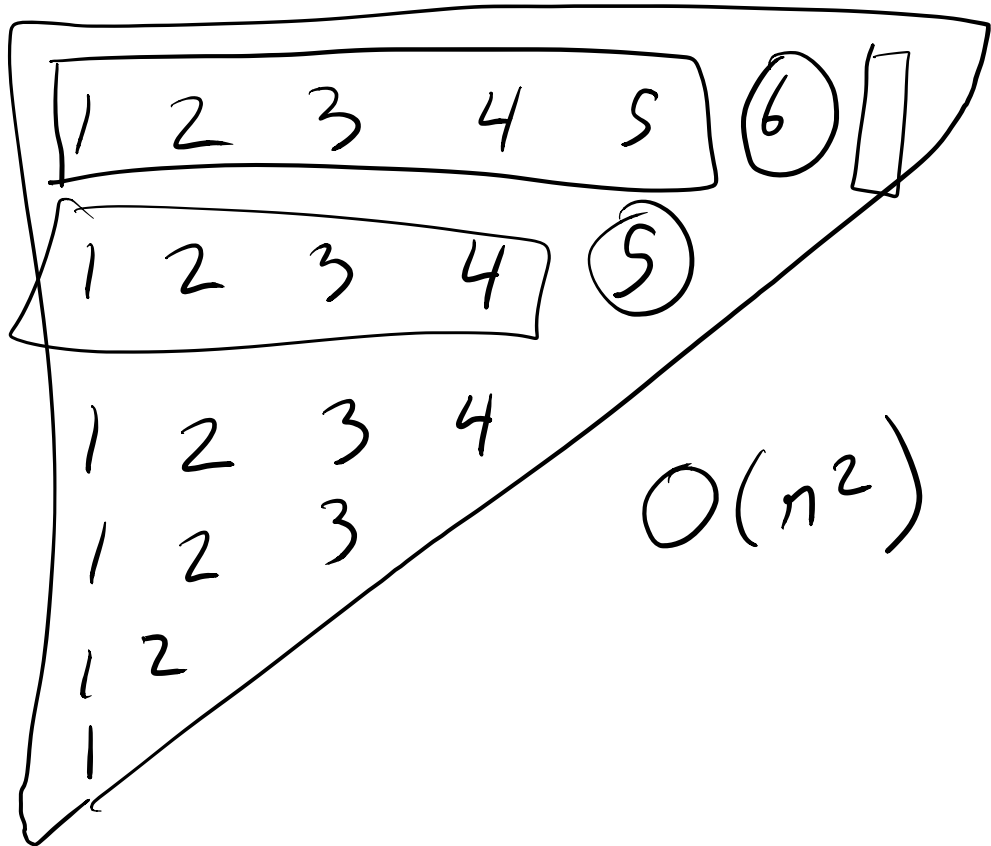
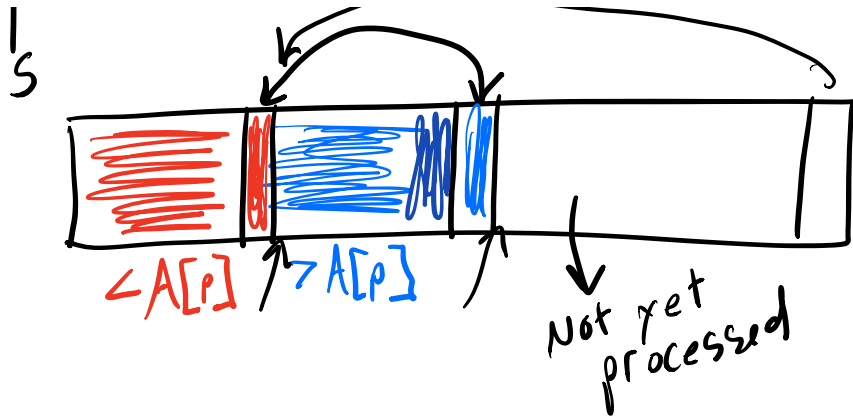
end while

swap $A[s], A[end]$

return s

end Function





Randomize!

Gambler's sort:

flip a coin

if heads : magically sort
in $O(1)$ time

if tails : flip again

$$\frac{1}{2} \cdot 1 + \frac{1}{4} \cdot 2 + \frac{1}{8} \cdot 3 + \dots$$

$$\sum_{i=1}^{\infty} \left(\frac{1}{2}\right)^i \cdot i = 2$$