CS35X: Competitive Programming

Lecture 8: Priority Queues, Dijkstra's Algorithm

Warmup Kattis Problem: vemvinner

Problem debrief: Button Bashing

Graphs revisited

- A graph G = (V,E) is a set of vertices V along with a set of edges E.
- Edges can be directed or undirected
 - Directed graph: edges u->v are asymmetric
 - *Undirected* graph: edges *u-v* are symmetric
- Edges can be weighted or unweighted

Which type of edges you have depends on application.

Motivation: Emergency Room Triage

Priority Queue ADT

- Maintain collection of (priority, value) pairs.
- Support the following operations:
 - Initialize an empty priority queue.
 - Insert a new (priority, value) pair.
 - Get the highest priority pair.
 - Remove the highest priority pair.
 - Check to see if a PQ is empty.
- Note: "high priority" can be minimum PQ or maximum PQ

Example Syntax

```
• #include <queue>
priority queue<int> pq;
pq.push(5);
                               // insert 5
                               // insert 3
pq.push(3);
cout << pq.top() << endl; // 5</li>
                               // remove 5
pq.pop();

    cout << pq.top() << endl; // 3</li>

                               // insert 6
pq.push(6);
cout << pq.top() << endl;</li>
```

C++ priority_queue details

- Stores just priorities
- To support (priority, value) pairs: use pair class.
 - priority_queue<pair<int,string>> myPQ;
- priority_queue is maximum priority by default.
- Make a minimum PQ by changing the comparison operator:
 - priority_queue<int, vector<int>, std::greater> minPQ;
- Sometimes it is useful to create your own comparison operator:
 - bool operator<(const Edge& rhs) const {
 return weight > rhs.weight;
 }

Exercise: Fancy Dog Show winners! week8/dogshow

Application: Dijkstra's algorithm

- Find shortest path in weighted graph from source to all other vertices.
- Idea: maintain dictionary of distance of shortest path to u we've seen
- Needs:
 - dist: map of current best distances from source to each vertex
 - minPQ: minimum PQ of candidate distances
- At each step, lock in distance from source to one node

```
map<int,int> dist;
priority_queue<pair<int,int>, vector<pair<int,int>>,
                   greater<pair<int,int>>> minPQ;
dist[source]=0;
minPQ.push(pair(0,source))
while(!minPQ.empty()) {
  int u = minPQ.top().second;
  int d = minPQ.top().first;
 minPQ.pop();
  if(d > dist[u]) continue;  // distance more than current best
  for (int j=0; j < g[u]. size(); j++) { // for each neighbor of u
    Edge e = g[u][j];
    v = e.dest;
    newcost = dist[u]+e.weight;
    if(!dist.count(v) || newcost < dist[v]) {</pre>
      dist[v] = newcost;
     minPQ.push(pair(newcost,v));
```

Dijkstra's Algorithm implementation details

 Two ways to handle Infinity: • const int INF = 1e9; // define INF to be a ridiculously large number • vertex not in dictionary == distance is infinite Defining an Edge class can be useful, e.g.: class Edge{ Public: int src, dest, weight; Edge(int src, int dest, int weight) { this->src = src; this->dest = dest; this->weight = weight; Create your own comparison operator: • bool operator<(const Edge& rhs) const { return weight > rhs.weight;

Kattis Problem: bumped