CS 43: Computer Networks Security

Kevin Webb Swarthmore College December 7, 2017

Topics

• Spying on network traffic

• Classic problem: buffer overflow attack

• Monetizing botnets

Once upon a time...

• The Internet was "a group of mutually trusting users attached to a transparent network."

- Result: Not much built-in security
 - Email headers
 - IP address spoofing
 - IP prefix hijacking

- ...

Once upon a time...

• Trust is gone, is the network still transparent?

• Switches help, cables can still be tapped...

• What about wireless?

• Wireless example/demo

Encryption

- Multiple options:
 - End to end (SSL, TLS): Browsers use this.
 - Link layer (WEP, WPA): Access point uses this.
- Facebook: enabled E2E encryption?
 July 2013

Cryptography

• Dates back 1000's of years

- Simple substitution cipher (Caesar cipher)

 Shift each letter by three (a -> d, b -> e, ...)
 "Hello world" becomes "khoor zruog"
- Many other, significantly better ciphers since...
 De facto standard today: AES

(Symmetric) Cryptography

• Problem: Encrypting with a cipher requires shared "key" information. (prior to 1970's)

 Sophisticated cipher doesn't help if we have to communicate the secret key!

- You want to ship a package to someone.
- You trust it won't be stolen, but might be read.







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(Asymmetric) Public Key Crypto

- Analogy: locking and unlocking are asymmetric
 - Anybody can lock
 - Very difficult to unlock without the key
- Let's apply this to data.
 - We need a function that's easy to apply in one direction and difficult in the other.

Factoring

- Multiplication is easy.
- Factoring (primes) is hard.

617077493 is the product of two primes.
 What are they?

• Rivest, Shamir, Adleman (RSA)

- Everyone computes two items:
 - Public key (kind of like a pad lock, ok if seen)
 - Private key (keep this to yourself)
- Receiver distributes public key, sender uses it to craft message that only receiver can read.

- Choose two prime numbers of similar length.
 - -P = 41 and Q = 29
- Compute N = P * Q- N = 41 * 29 = 1189
- Compute phi(N) = (P − 1) * (Q − 1)
 − phi(N) = (41 − 1) * (29 − 1) = 40 * 28 = 1120
- Choose a value e, 1 < e < 1120

– e must not divide 1120, we'll pick 13

- P = 41 and Q = 29 N = 1189 phi(N) = 1120 e = 13
- Compute "modular multiplicative inverse" of e – Need: (d * e) % phi(N) = 1 – d = 517

P = 41 and Q = 29

N = 1189, phi(N) = 1120, e = 13, d = 517

• Public key is (n, e), private key is (n, d)

- To encrypt message m = 1000:
 Take 1000¹³ % 1189 = 611
- To decrypt message 611:
 611⁵¹⁷ % 1189 = 1000

In Practice...

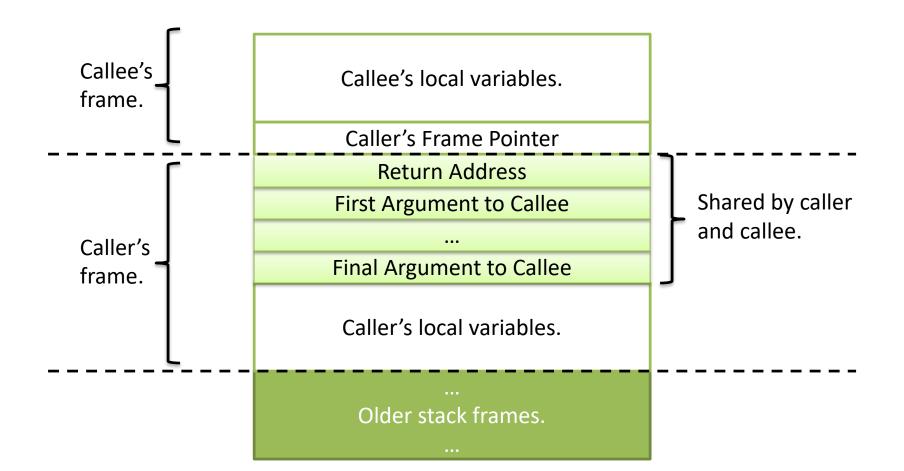
- Result: we can exchange secure messages with parties we've never talked to before!
 – (e.g., your bank)
- Exchange a secure message containing shared secret via RSA (asymmetric crypto)
- Subsequently use shared secret for conventional symmetric crypto (e.g., AES)

Classic Attack: Buffer Overflow

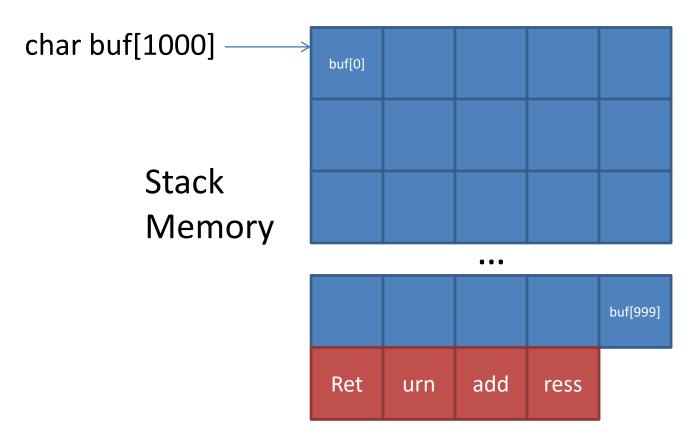
• Encryption ruining your (evil) day?

• Let's try taking control instead!

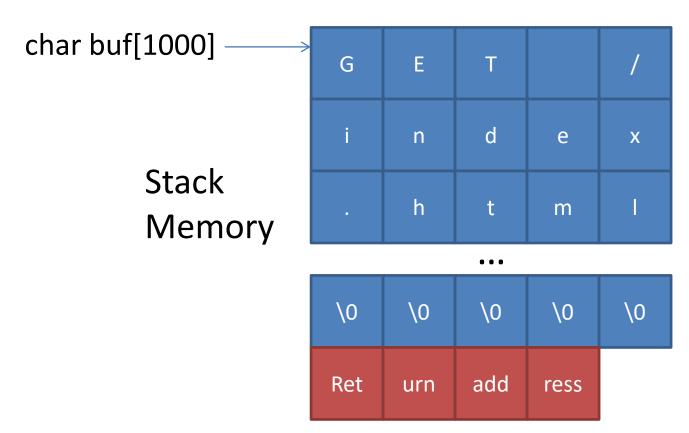
Recall: The Stack



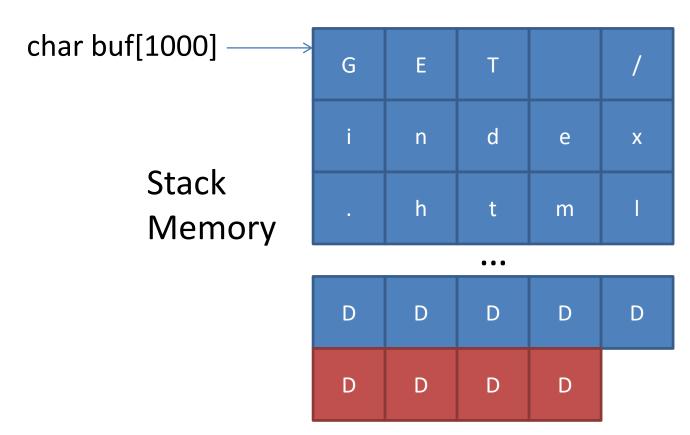
 Suppose we have a protocol that does recv() until it finds \r\n\r\n.



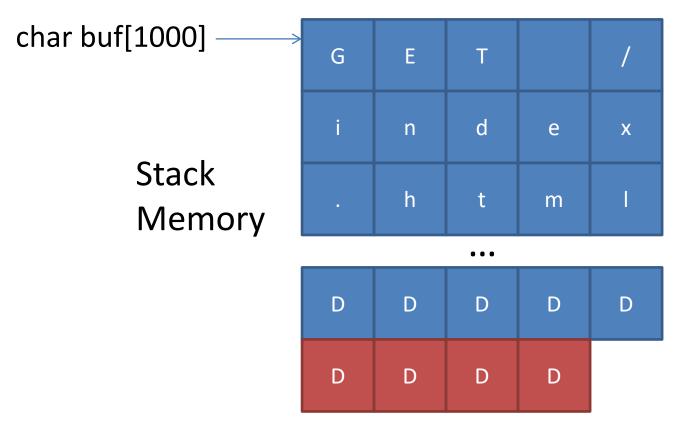
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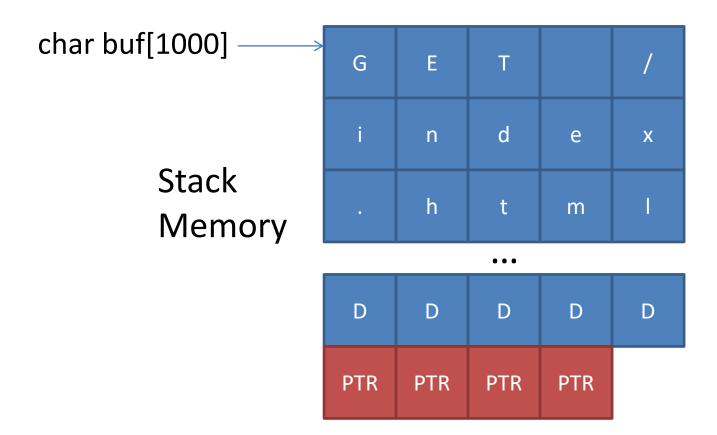
 What happens if we're sent more than 1000 bytes before we see \r\n\r\n? Keep writing...



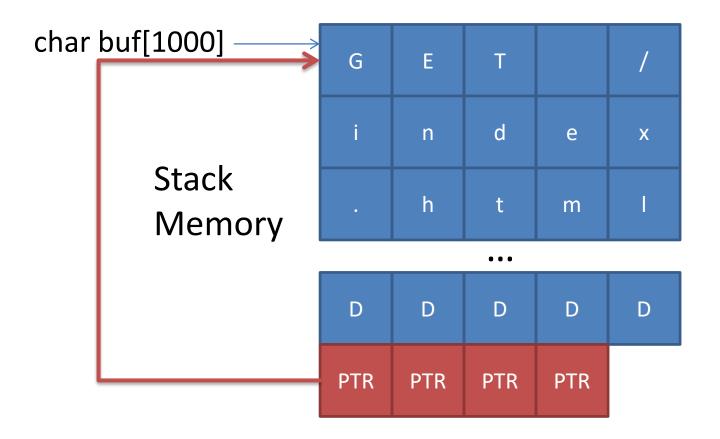
- Uh, if we can overwrite the return address...
- We can control execution on return.



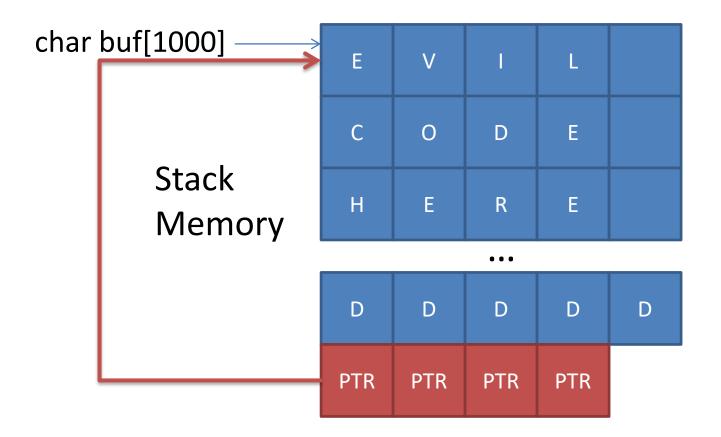
• Let's send malicious data that contains a ptr.



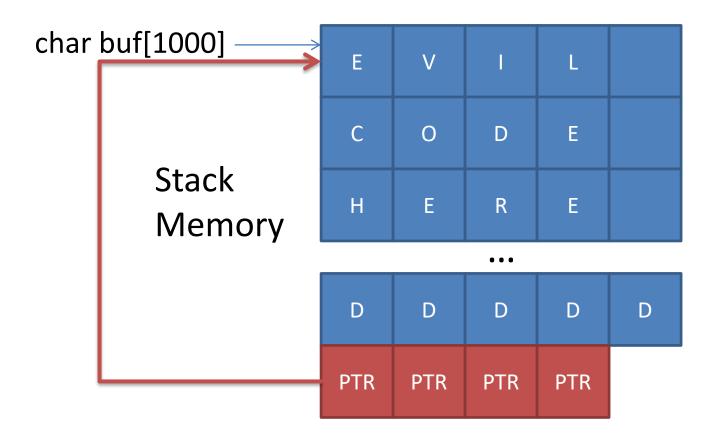
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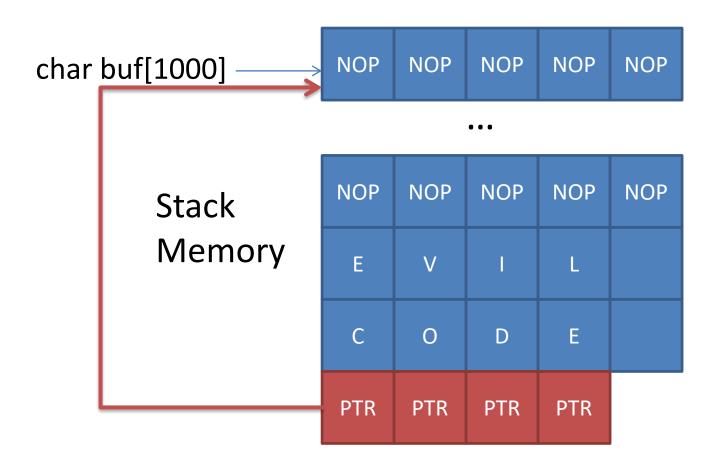
• Oh, and also some commands up here...



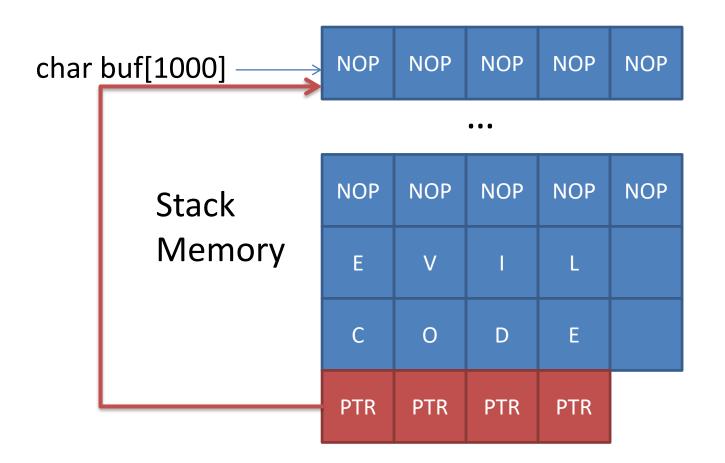
• Function returns, executes evil code.



Improve chances: "NO OP sled"



• See: "Smashing the Stack for Fun and Profit"



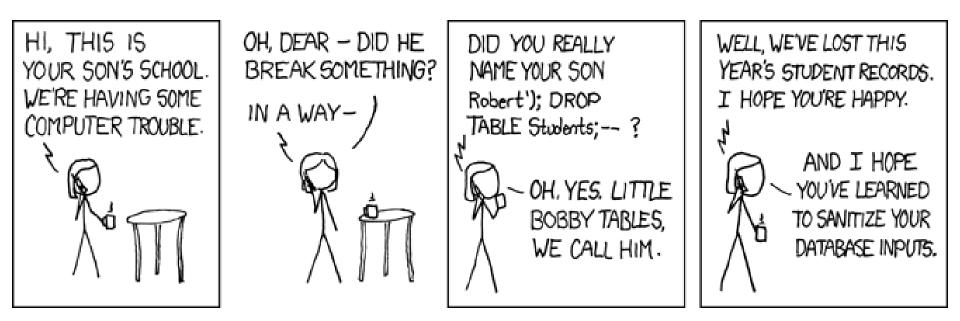
Input from Network

 Programs that receive user input are susceptible to buffer overflow (& more) attacks.

• Potentially much more problematic to receive input from the Internet!

• If attackers can take over program's control flow, they can execute *anything*.

Relevant XKCD #327



Alt text: Her daughter is named Help I'm trapped in a driver's license factory.

Bottom line: be careful about what you're accepting from the network! Make sure the memory you're using is bounded and that the data is valid!

1988: The Morris Worm

- Cornell student Robert Morris
- Exploited buffer overflow in fingerd

 It had a 512-byte buffer, he exploited it to execute /bin/sh, giving him shell access
- Told compromised host to download his worm code, it self-replicated by exploiting others
- Claimed "wanted to gauge size of Internet"

1988: The Morris Worm

- Worm did a check to see if it needed to replicate itself
 - If machine already compromised (process running) don't infect again.
- Worried about admins putting up fake process
 Replicate anyway, at random, 1/7 times.
- This effectively shut down LOTS of machines.

1988: The Morris Worm

- Robert Morris:
 - First person convicted under Computer Fraud and Abuse Act
 - Sentenced to three years probation, 400 hours community service, \$10,000
- Where is he now?

Exploits Today

- Worms
- Trojans (trick user)
- Browser exploits (drive-by downloads)

• Often used in BotNets

BotNets

 Having access to 1000's of machines is lucrative!

- Send Spam.
- Flood target with traffic (DDoS).
- Steal data (CC #'s, state secrets, etc.).
- Mine bitcoins.

Questions?

Final Exam

• Friday, December 15, at 9:00

• LOCATION: SCI 199